

AD2426/7/8 Automotive Audio Bus (A²B) Transceiver Technical Reference

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1 Preface

Thank you for purchasing and developing systems using an Automotive Audio Bus A²B[®] Transceiver from Analog Devices.

Purpose of This Manual

The *AD2426/7/8 Automotive Audio Bus A²B Transceiver Technical Reference* provides information about the transceivers, including register and bit descriptions. For timing, electrical, and package specifications, see the *AD2426/7/8 Automotive Audio Bus A²B Transceiver Data Sheet*.

Intended Audience

This manual is intended for system designers and programmers who want to develop systems using the A²B transceiver.

Manual Contents

This manual consists of the following chapters:

- *A²B Overview* - Provides a basic description and the features supported.
- *A²B Operation and Configuration* - Provides information on bringing up the main node and discovery of the subordinate nodes. Provides a simple System Discovery Example.
- *A²B Event Control* - Provides information on system interrupts and their use.
- *A²B System Debug* - Provides information that allows you to perform system diagnostics in order to isolate and correct faults. Additionally, a loop back test mode provides easy validation of I²S/TDM connectivity in main and subordinate nodes.
- *Register Summary* - Provides the register map and bit definitions for the integrated transceiver.
- *Register Descriptions* - Provides the detailed descriptions of the registers and bits.
- *Appendix A: Additional Discovery Flow Examples*

- *Appendix B: Response Cycle Formula*
- *Appendix C: Module ID and Module Configuration Memory*
- *Appendix D: Interrupt Processing*

What's New in This Manual

This is revision (1.3). The following changes were made to content in this revision:

- Document title change to reflect updated part numbering
- Updated the *Simple Discovery Flow* figure in the A²B Operation and Configuration chapter.
- Updated topics:
 - Subordinate Node Interrupt Handling
 - Diagnostics Software Flow
- Updated Register Summary section to correct the reset values for the NODE, MBOX0STAT, MBOX1STAT, and MOBIX1CTL registers.

Register Documentation Conventions

The register sections and diagrams use the following conventions:

- Registers are presented in address order.
- The reset value appears in binary in the individual bits and in hexadecimal to the left of the register.
- Shaded bits are reserved.

NOTE: To ensure upward compatibility with future implementations, write back the value that is read for reserved bits in a register, unless otherwise specified.

Register description tables use the following conventions:

- Each bit's or bit field's access type appears beneath the bit number in the table in the form (read-access/write-access). The access types include:
 - R= read, RC= read clear, RS= read set, R0= read zero, R1= read one, Rx= read undefined
 - W= write, NW= no write, W1C= write one to clear, W1S= write one to set, W0C= write zero to clear, W0S= write zero to set, WS= write to set, WC = write to clear, W1A= write one action, XCVRA/B= transceiver (A-port /B-port)
- Many bit and bit field descriptions include enumerations, identifying bit values and related functionality. Unless otherwise indicated (with a prefix), these enumerations are decimal values.

2 A²B Overview

The Automotive Audio Bus (A²B[®]) connects multichannel I²S synchronous PCM data over a distance of up to 15m between nodes. It also extends the synchronous, time-division multiplexed (TDM) nature of I²S to a system that connects multiple nodes, where each node can consume data, provide data, or both.

The transceivers support these A²B functions with a direct interface to general-purpose digital signal processors (DSPs), field programmable gate arrays (FPGAs), application specific integrated circuits (ASICs), microphones, analog-to-digital converters (ADCs), digital-to-analog converters (DACs), and codecs through a multichannel I²S/TDM interface. They also provide a PDM interface for direct connection of up to four PDM digital microphones. Enabling the A²B bus-powering feature supplies voltage and current to the subordinate nodes over the same, daisy-chained, twisted pair wire cable as used for the communication link. The transceiver also fully supports I²C communication over the A²B link.

The transceivers have the following features.

- Line topology
 - Single main node, multiple subordinate nodes
 - Unshielded, single twisted pair wire (UTP) cable link between nodes (cable length is specified in the product data sheet)
- Communication over distance
 - Synchronous data
 - Multichannel I²S/TDM to I²S/TDM interface
 - Synchronous, phase-aligned clock in all nodes
 - Low-latency subordinate node-to-subordinate node communication
 - I²C to I²C control and status information
 - GPIO over distance
- Configurable with SigmaStudio[®] graphical development tool
- Qualified for automotive applications

- Configurable as A²B bus main or subordinate
- I²C interface
- 8-bit to 32-bit multichannel I²S/TDM interface
 - I²S/TDM/PDM programmable data rate
 - Up to 32 channels (1 x TDM32 or 2 x TDM16), mapped to up to 32 upstream and 32 downstream A²B bus slots
- PDM inputs supporting up to 4 high-dynamic-range microphones
- Unique ID register for each transceiver
- Support for crossover or straight-through cabling
- Programmable settings to optimize EMC performance

A²B Terminology

To make the best use of the A²B system, it is helpful to understand the following terms.

A-Side or A-Port

A²B transceiver interface that faces toward the main (toward the immediately upstream node).

B-Side or B-Port

A²B transceiver interface that faces toward the last-in-line subordinate node (toward the immediately downstream next-in-line subordinate node).

Bus Link

The A²B bus can consist of multiple daisy-chained subordinate nodes connected to a single main node. The physical connection between a main and sub node 0, as well as all physical A²B connections between subordinates, are called bus links. An unshielded twisted wire pair is typically used for each bus link.

Bus Power

Subordinate nodes can tap into the bias voltage on the A²B bus link and use it as the sole power supply.

Data Channel

A data channel carries the synchronous I²S/TDM data for a single sensor/actuator (for example, an ADC, a microphone, or a speaker). The I²S/TDM interface uses equally sized data channels, where the width of the data word is

often smaller than the width of the I²S/TDM data channel. The I²S/TDM interface of the transceiver supports programmable data channel lengths of 16 or 32 bits.

Data Slot

A synchronous data word of a single sensor/actuator (for example, an ADC, a microphone, or a speaker), as mapped onto the A²B bus.

Downstream

Communication flow from the main node toward the subordinate nodes, terminating at the last-in-line subordinate.

Host

Processor that programs the main transceiver. The host is also the source for the synchronous clock on the A²B bus. The clock signal (BCLK) is part of the I²S/TDM interface between the host and main.

I²S/TDM

The inter IC sound (I²S) bus carries pulse code modulated (PCM) information between audio chips on a PCB. The I²S/TDM interface extends the I²S stereo (2-channel) content to multiple channels using time-division multiplexing (TDM).

Local Power Subordinate Node (LPS)

Subordinate nodes that do not operate on A²B bus power use local power, which is sourced by extra wires.

LVDS

Low voltage differential signaling.

Main Node

Originator of the clock (derived from the I²S input), downstream data, network control, and power. The main node is comprised of the host processor and an A²B main transceiver, which receives payloads from the host and sends payloads to the host.

PDM

Pulse Density Modulation (PDM) is used in sigma delta converters. PDM format represents an over-sampled 1-bit sigma delta ADC signal before decimation and is often used as the output format in digital microphones.

PRBS

Pseudo random binary sequence.

Preamble

Synchronization bits to signal the start of a control or response frame. The downstream control frame preamble is sent by the main node for every superframe. Subordinate transceivers synchronize to the downstream control preamble and generate a local, phase-aligned main clock from it.

Response Time

Specifies the time a last node waits after the start of a superframe before the node responds with the Synchronization Response Frame (SRF). Response time is programmed in the main node and all subordinate nodes closest to the main so that these nodes know when to expect the direction to switch from downstream to upstream.

Subordinate Node

Addressable network connection point. Subordinate nodes can be the source and/or destination of both downstream and upstream data slots. Every A²B subordinate node has an A²B subordinate transceiver.

Synchronization Control and Response Frames (SCF/SRF)

Control frame for nodes (control header) and response frame from nodes (response header). Headers include a preamble for synchronization and enable read and write access to all nodes.

Synchronous Data

Data streamed continuously (for example, audio signals) with a fixed time interval (selectable between 44.1 kHz or 48 kHz) between two successive transmissions to and from the same node.

Superframe

The overall frame structure for A²B. It starts with an SCF, includes optional data slots, and concludes with an SRF. Superframes repeat every 1024 bus clock cycles.

Upstream

Communication flow the last-in-line subordinate node to the main node.

A²B Bus Details

The *Communication System Block Diagram* shows an A²B communications system, which is a single-main, multiple-subordinate system where the main transceiver is controlled by the host. The host generates a periodic

synchronization signal (SYNC) on the I²S/TDM interface at a fixed frequency (selectable between 44.1 kHz and 48 kHz), to which all A²B nodes synchronize. Communication over the A²B bus occurs in periodic superframes at this rate. Data is transferred at the A²B system bit clock (SYSBCLK) rate, which is 1024 times faster than the superframe rate (49.152 MHz for a frame rate of 48 kHz, 45.158 MHz for a frame rate of 44.1 kHz). Each superframe is divided into periods of downstream transmission, upstream transmission, and no transmission (where the bus is not driven).

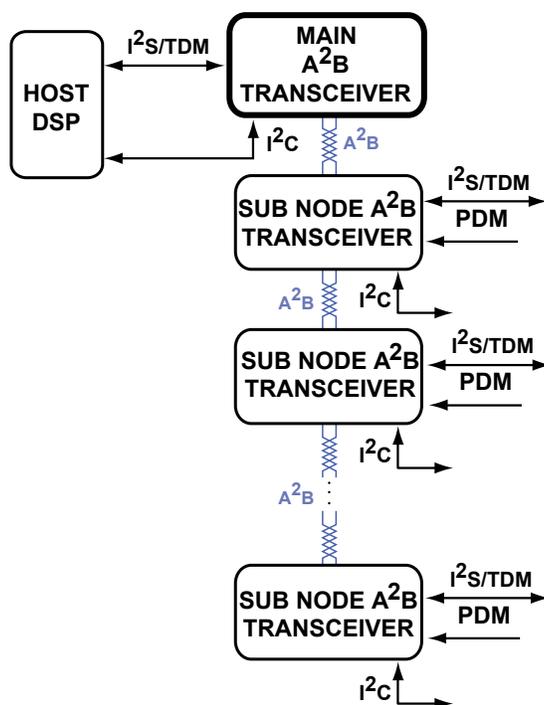


Figure 2-1: Communication System Block Diagram

The *A²B Superframe* figure shows a superframe with an initial period of downstream transmission and a later period of upstream transmission.

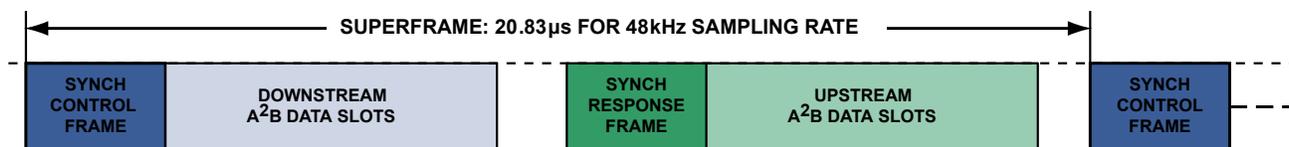


Figure 2-2: A²B Superframe

All signals on the A²B bus are line-coded, and the main node forwards the synchronization signal downstream to the last subordinate node in the form of a synchronization preamble. This preamble is followed by the control frame (SCF). Downstream, TDM synchronous data is added directly after the control frame. Every subordinate node can use or consume some of the downstream data and add data for downstream nodes. The last subordinate node responds after the response time with a response frame (SRF). Upstream synchronous data is added by each node directly after the response frame. Each node can also use or consume upstream data. All synchronous data is organized

into data slots of equal width, though the upstream and downstream slot widths can be different. For more details, see [A²B Slot Format](#).

The embedded control and response frames allow the host to individually address each subordinate node over the A²B bus. In a similar fashion, the host can also access remote peripheral devices that are connected to any discovered subordinate transceivers using I²C- to-I²C communication over distance.

All nodes in an A²B system are sampled synchronously in the same A²B superframe. Synchronous I²S/TDM downstream data from the main arrives at all subordinates in the same A²B superframe, and each node's upstream audio data arrives synchronously in the same I²S/TDM frame at the main. The remaining audio phase differences between subordinates can be compensated for by register-programmable fine adjustment of the SYNC pin signal delay using the `A2B_SYNCOFFSET` register.

Because data is received and transmitted over the I²S/TDM port every sample period, there is a delay incurred for data moving between the A²B bus and the I²S/TDM interfaces. The timing relationship between samples over the A²B bus is shown in the *A²B Bus Synchronous Data Exchange* figure.

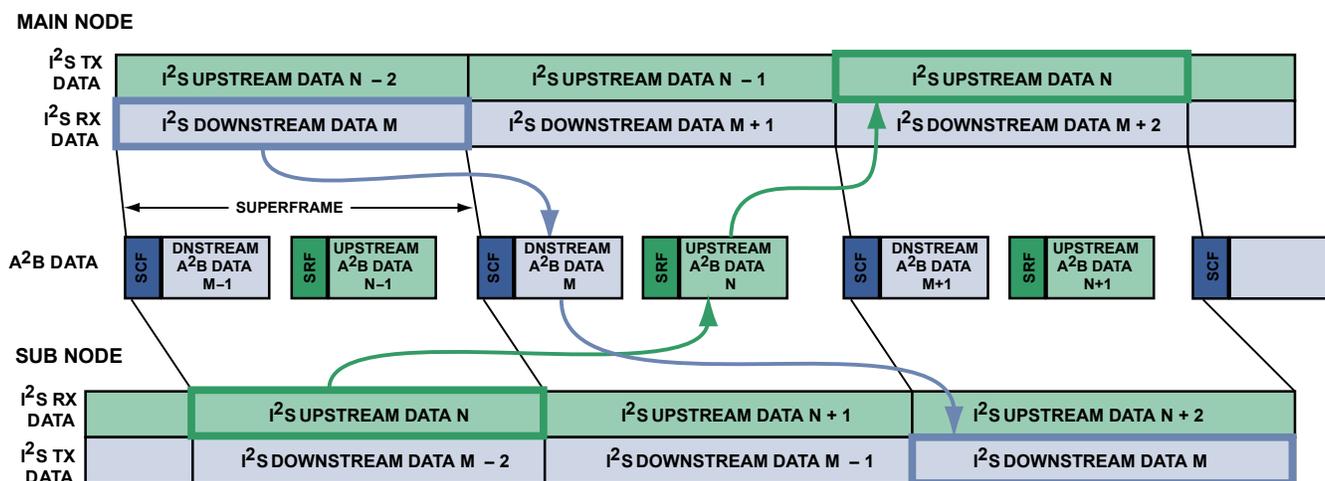


Figure 2-3: A²B Bus Synchronous Data Exchange

Note in the *A²B Bus Synchronous Data Exchange* figure, both downstream and upstream samples are named for the superframe where they enter the A²B system, as follows:

- Data transmitted by the main node transceiver in superframe M creates downstream data M
- Data transmitted by the subordinate node transceivers in superframe N creates upstream data N
- Data received over the I²S/TDM interface by the A²B transceiver chip is transmitted over the A²B bus in the following superframe
- Data on the A²B bus is transmitted over the I²S/TDM interface of an A²B chip transceiver in the following superframe

- Data transmitted across the A²B bus (main node to subordinate node or subordinate to main) has two superframes of latency, plus any internal delay that has accumulated in the transceiver chips, as well as delays due to wire length. Therefore, overall latency is slightly over two superframes from the I²S/TDM interface in one A²B transceiver chip to the I²S/TDM interface of another A²B transceiver chip.

Functional Description

The A²B transceiver connects multichannel I²S (inter-IC sound) synchronous, pulse-code modulated (PCM) data over a distance between nodes (the cable length is specified in the product data sheet). It also extends the synchronous, time-division multiplexed nature of I²S to a system that connects multiple nodes, where each node can consume data, provide data, or both.

The A²B transceiver supports these A²B functions with a direct interface to general-purpose DSPs, FPGAs, ASICs, microphones, ADCs, DACs, and codecs through a multichannel I²S/TDM interface. The data over the A²B bus link is manchester encoded. The transceiver also fully supports I²C communication over the A²B link. The A²B transceiver can be used in either a subordinate node or in a main node. By default, the transceiver starts up as a subordinate transceiver but can be configured as a main transceiver if the host sets the `A2B_CONTROL.MSTR` bit.

The *Simplified A²B System with Four Nodes* figure shows a simple A²B system example. The host programs registers in each of the nodes to control the data traffic on the A²B bus. Microphone data from subordinate nodes 0 and 2 is delivered to the host, and speaker data for subordinate nodes 1 and 2 is delivered from the host to the DACs.

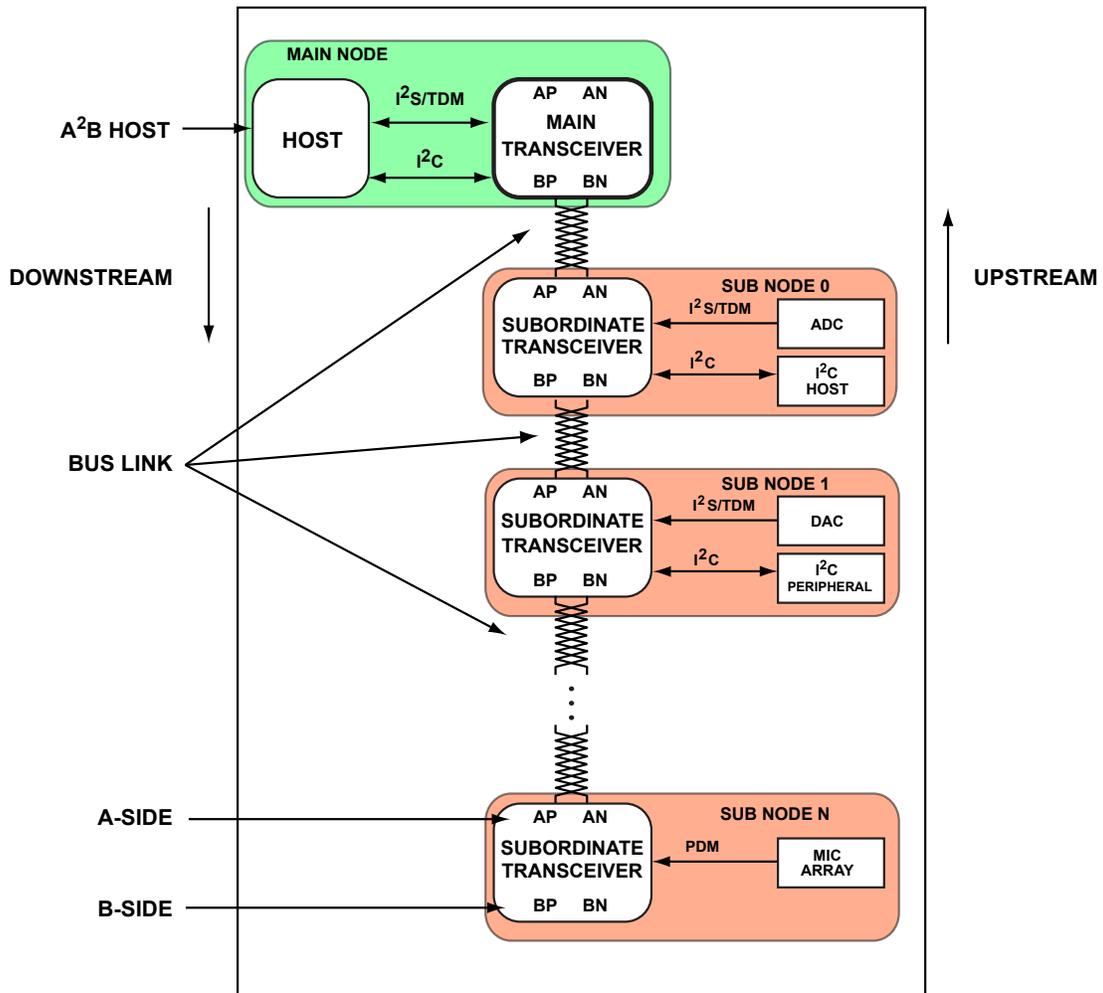


Figure 2-4: Simplified A²B System with Four Nodes

Architectural Concepts

The following sections provide information that describes the hardware blocks, interfaces and interconnections.

I²C Interface

The I²C interface is used to directly access the transceiver register space from a locally connected host and to remotely exchange I²C data over the A²B bus between the main transceiver and any discovered subordinate node in the system. This protocol is referred to as *I²C over distance*, where the exchanged I²C data is embedded within the synchronization control frame (downstream, from the I²C controller to the target) and the synchronization response frame (upstream, from the I²C target to the controller).

The I²C interface in the transceiver is compatible with up to 5 V logic levels and has the following features:

- Target only operation in an A²B main node

- Controller, multi-controller, or target operation in an A²B sub node
- Operations at 100k or 400k bits/s rate, as configured by the A2B_I2CCFG.DATARATE bit
- 7-bit addressing
- Clock stretching

NOTE: The A²B host on the main node must support I²C clock stretching in order to interface to the main transceiver.

A transceiver that is configured as a main recognizes two I²C device addresses:

- BASE_ADDR - for direct accesses via the I²C port to its register space
- BUS_ADDR - for remote access to sub node registers and target node I²C peripherals over the A²B bus using the *I²C over distance* protocol

The I²C BASE_ADDR is set by the logic levels on the ADR2/IO2 and ADR1/IO1 pins at power-on reset, thus providing support for up to four controller devices connecting to the same I²C bus. The LSB of the 7-bit device address determines whether an I²C data exchange uses the BASE_ADDR (bit 1 = 0) to access the transceiver or BUS_ADDR (bit 1 = 1) to access a bus node through a controller-enabled transceiver, as described in the *I²C Address* table.

Table 2-1: I²C Device Address

| ADR2/IO2 Setting | ADR1/IO1 Setting | BASE_ADDR | BUS_ADDR |
|------------------|------------------|-----------|----------|
| 0 | 0 | 0x68 | 0x69 |
| 0 | 1 | 0x6A | 0x6B |
| 1 | 0 | 0x6C | 0x6D |
| 1 | 1 | 0x6E | 0x6F |

A transceiver that is configured as a subordinate does not recognize BUS_ADDR. On subordinate transceivers, the I²C interface allows for both I²C controller and target behavior. It is the I²C controller when the transceiver receives a remote I²C peripheral access request from the host through the A²B bus. The main transceiver functioning as the I²C controller then forwards the I²C transaction to the I²C target address programmed in its A2B_CHIP register. It is the I²C target when the transceiver registers (BASE_ADDR) are accessed by a local external controller through the I²C port.

NOTE: While a local external controller can program the register space of a subordinate transceiver, the A2B_SWCTL, A2B_RESPCYCS, A2B_SLOTFMT, A2B_DATCTL, A2B_RAISE, and A2B_GENERR registers must be written over the A²B bus by the remote host. A write to any of these registers from the local I²C port has no effect on the register.

The I²C interface on the transceiver allows register programming before PLL lock. Write 1 for action (W1A) bits (for example, `A2B_CONTROL.ENDDSC` and `A2B_CONTROL.NEWSTRCT`) have no effect prior to PLL lock since the protocol engine is still in reset.

NOTE: The `A2B_SWCTL`, `A2B_SLOTFMT`, `A2B_DATCTL`, and `A2B_DISCVRY` registers cannot be written in a main transceiver prior to PLL lock. Writes to these registers before PLL lock is established have no effect.

CAUTION: System software must be designed to avoid simultaneous writes to the same subordinate register from both the A²B host (through the A²B bus) and the local processor (through the I²C port). When write contention occurs, both writes complete, but the order in which they complete is unpredictable.

I²C Clock Stretching

The transceiver uses the I²C clock stretching feature to ensure that the I²C accesses have enough time to be processed. It is applied mainly for host I²C accesses to sub node transceivers and target node I²C peripherals over the A²B bus. Clock stretching is initiated by the main transceiver in response to host I²C accesses at the following times:

- During write accesses – before the acknowledge bit after each data byte
- During read accesses – before the acknowledge bit following the read request
- During burst read/write accesses of more than one byte – before the first bit of subsequent data bytes

Pulling the SCL signal low indicates to the host that the transceiver needs more time to process the request. Once the transceiver is ready to acknowledge the request, it lets the SCL signal go high so the host can gain back control of the SCL and proceed with the acknowledge (ACK) and the next byte.

IMPORTANT: It is mandatory that the host (I²C controller) supports I²C clock stretching in an A²B system design.

When a peripheral in a sub node stretches the I²C clock, the SCL signal is also stretched between the main transceiver and the host. If the SCL signal is not released by the peripheral within the time of 32 superframes, the main transceiver registers a timeout (`A2B_INTPND2.I2CERR = 1`), releases the SCL, and ceases stretching of the host clock. This timeout ensures a subordinate peripheral cannot bring the I²C interface of the host to a permanent halt.

Transceiver I²C Accesses

The LSB of the 7-bit device address determines whether an I²C data exchange uses the `BASE_ADDR` (bit 1 = 0) to access the transceiver or `BUS_ADDR` (bit 1 = 1) to access a bus node through a main configured transceiver, as shown in the following table.

Table 2-2: I²C Device Addresses

| Bit Number | Bit 7 | Bit 6 | Bit 5 | Bit 4 | Bit 3 | Bit 2 | Bit 1 (LSB) | Bit 0 (R/ \overline{W}) |
|------------|-------|-------|-------|-------|----------|----------|---------------------|----------------------------|
| Start bit | 1 | 1 | 0 | 1 | ADR2/IO2 | ADR1/IO1 | 0 = BASE 1 = BUS | 0 = write 1 = read |

The A²B transceiver supports the following read and write operations:

- Single-word write operation – the A²B main transceiver (I²C target) issues an acknowledge by pulling SDA low during the ninth clock pulse, thus completing the access.
- Burst mode write sequence – the transceiver automatically increments the register address pointer after each data byte, so sequential data registers can be written without reprogramming the address.
- Single-word read operation – the first read/write (R/ \overline{W}) bit is 0, indicating a write operation. This is because the register address must still be written to set up the internal address. After the I²C target acknowledges the receipt of the register address, the I²C controller must issue a repeated start command, followed by the chip address byte with the R/ \overline{W} bit set to 1 (read). This causes the I²C data line SDA to reverse direction and begin driving data back to the I²C controller. The I²C controller then responds every ninth pulse with an acknowledge pulse to the target.
- Burst mode read sequence – the transceiver automatically increments the register address pointer after every read of a data byte, so sequential data registers can be read without reprogramming the address.

Data transfers over the I²C interface require the following steps:

1. A data transfer is initiated by a microcontroller that is connected to an A²B transceiver.
2. The microcontroller establishes a start condition (a high to low transition on SDA while SCL remains high), which indicates that an address/data stream follows.
3. In the next eight SCL cycles, the A²B transceiver receives a 7-bit address and the R/ \overline{W} bit from the host (MSB first).
4. The A²B transceiver recognizes the transmitted address and responds by pulling the data line low during the ninth clock pulse (acknowledge bit).

The R/ \overline{W} bit determines the direction of the data. When the LSB of the first byte is cleared (=0), the host writes information to the main transceiver. When the LSB of the first byte is set (=1), the host reads information from the main transceiver. Data transfers take place until a stop condition (when SDA transitions from low to high while SCL is held high) is encountered. The register address pointer auto increments to support burst mode I²C writes and burst mode I²C reads for both mains and subordinates.

The *I²C Formats* figure shows the format of the following I²C operations:

- Writes to BASE_ADDR/BUS_ADDR can contain one or more bytes of data. The first byte after the device address sets the register address in the device. The subsequent byte is written to the addressed register. Since the address pointer increments after each write, sequential registers can be written in a single transaction.
- Reads from BASE_ADDR/BUS_ADDR can contain one or more bytes of data. The device address with write indication is followed by the register address in the device and a repeated device address with a read access indication.

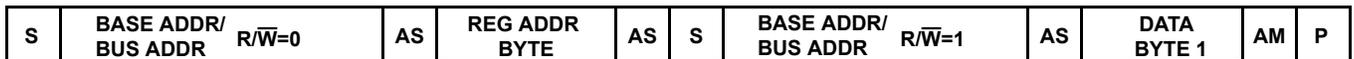
SINGLE WORD I²C WRITE FORMAT



BURST MODE I²C WRITE FORMAT



SINGLE WORD I²C READ FORMAT



BURST MODE I²C READ FORMAT



S = START BIT
P = STOP BIT
AM = ACKNOWLEDGE BY I²C CONTROLLER
AS = ACKNOWLEDGE BY I²C TARGET

Figure 2-5: I²C Formats

The first byte after the repeated device address contains the value of the register addressed. The first byte after the device address sets the register address in the device. It is followed by the repeated device address, but with read access indication. The subsequent bytes contain the values of the automatically incremented register addresses.

The *I²C Write Timing* figure shows I²C write timing.

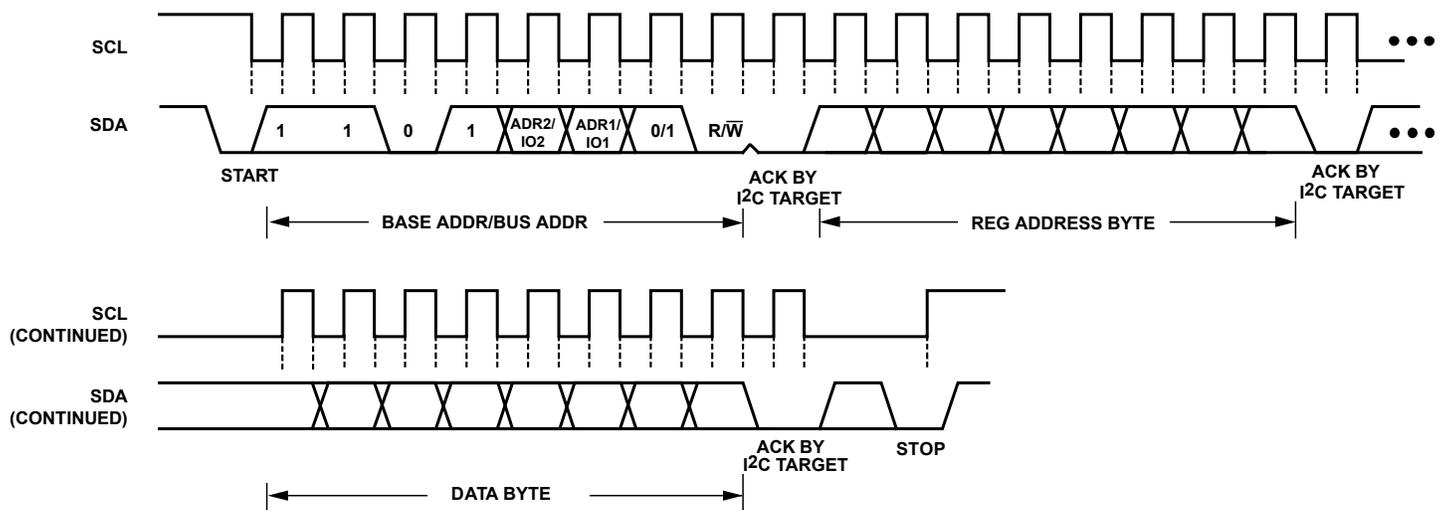
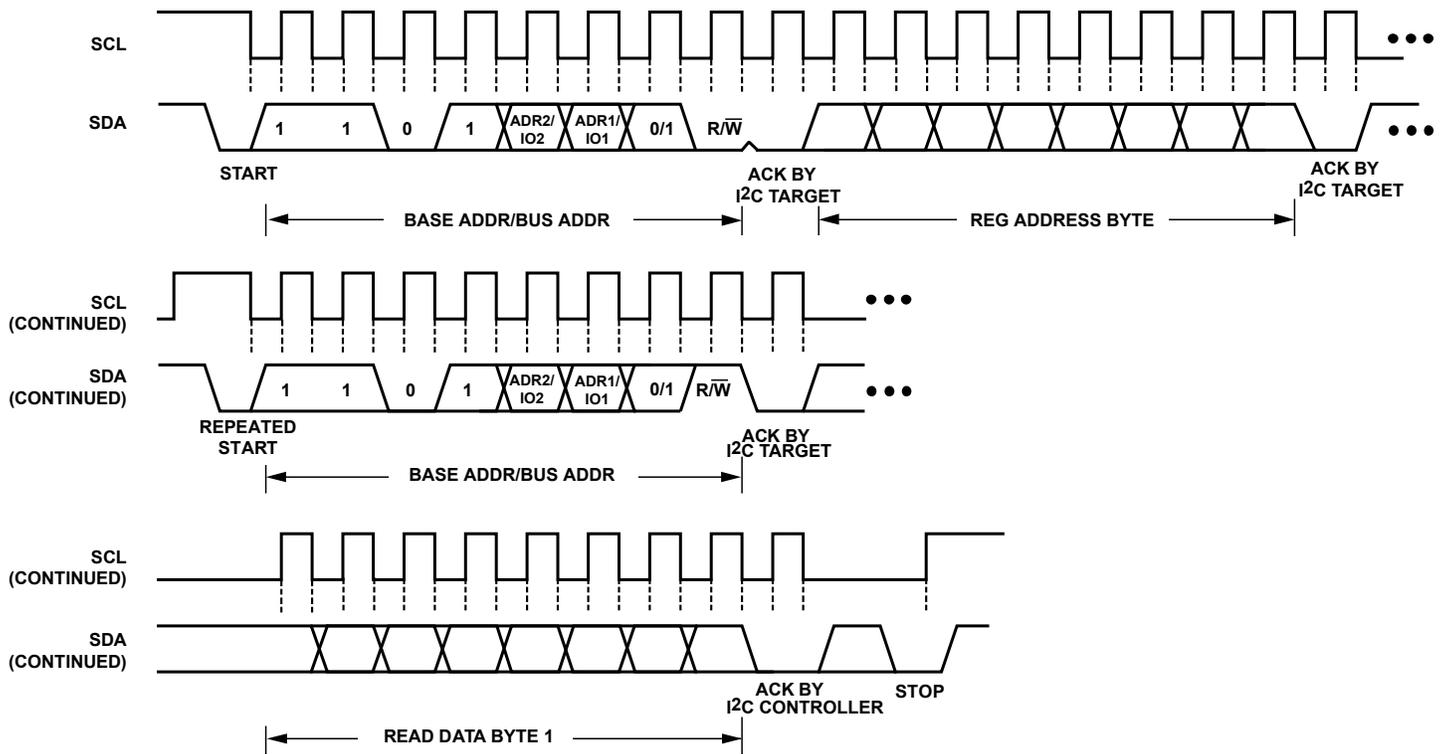


Figure 2-6: I²C Write Timing

The *I²C Read Timing* figure shows I²C read timing.

Figure 2-7: I²C Read Timing

Transceiver I²C Access Latencies

When an I²C access is made over distance to a remote transceiver via the A²B bus, there are latencies incurred. A²B bus latencies for different types of I²C accesses are provided in the *Bus Latencies for I²C Accesses* table.

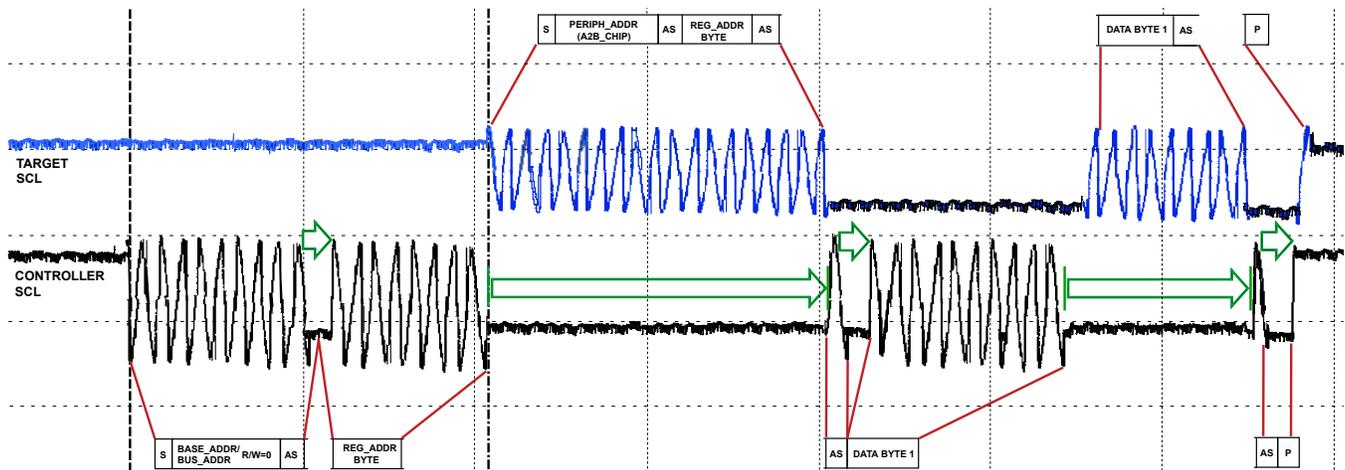
Table 2-3: Bus Latencies for I²C Accesses (48 kHz Superframe Rate)

| I ² C Access Type (Conditions) | Estimated A ² B Bus Latency (μs) |
|--|---|
| I ² C write of N data bytes to subordinate transceiver registers (clock stretching enabled via main A2B_I2CCFG.EACK = 0) | $N \times 22$ |
| I ² C read of N data bytes from subordinate transceiver registers (clock stretching enabled via main A2B_I2CCFG.EACK = 0) | $N \times 22$ |
| I ² C write of N >1 data byte to subordinate transceiver registers (clock stretching disabled via main A2B_I2CCFG.EACK = 1, host I ² C using 400 kHz data rate) | 2 |
| I ² C write of N data bytes to subordinate transceiver registers (clock stretching enabled via main A2B_I2CCFG.EACK= 1, host I ² C using 100 kHz data rate) | 0 |

Table 2-3: Bus Latencies for I²C Accesses (48 kHz Superframe Rate) (Continued)

| I ² C Access Type (Conditions) | Estimated A ² B Bus Latency (μs) |
|--|---|
| I ² C write of N data bytes to remote I ² C peripheral (target A2B_I2CCFG.DATARATE = 0 = 100 kHz) | $((N - 1) \times 113) + 213$ |
| I ² C write of N data bytes to remote I ² C peripheral (target A2B_I2CCFG.DATARATE = 1 = 400 kHz) | $((N - 1) \times 45) + 70$ |

For example, consider a case where a remote peripheral (connected to a subordinate node) register is being written. In the *I²C Access Latency* figure, the I²C access latency is marked with green arrows.

Figure 2-8: I²C Access Latency

NOTE: The latencies described in the *Bus Latencies for I²C Accesses* table are for accesses with no conflicts. If an I²C message doesn't get immediately acknowledged or is otherwise held off due to higher-priority events such as a GPIO interrupt, a line fault interrupt, an I²C issue (NACK), etc., the delay incurred before attempting to execute the message exchange is not included in the values provided in the table.

Pulse-Density Modulation Interface (PDM)

Pulse-density modulation is used in sigma delta converters. The PDM format represents an over-sampled 1-bit sigma delta ADC signal before decimation. It is often used as the output format in digital microphones.

The PDM block supports high dynamic range microphones with a high signal-to-noise ratio (SNR) and an extended maximum sound pressure level (SPL).

The enhanced PDM block of the transceiver supports a lower noise floor than the AD241x transceiver. This provides for an SNR greater than 120 dB. The PDM block on the transceiver supports 24 kHz and 12 kHz sample rates in addition to a 48 kHz sample rate with the same PDM clock rate (3.072 MHz at a 48 kHz frame rate). The cutoff frequency of the high pass filter in the PDM block on the transceiver is fixed to 1 Hz and is not programmable. The highpass filter is a first order IIR filter.

The transceiver is programmable for 1x, 1/2x, or 1/4x PDM sampling (48 kHz, 24 kHz, or 12 kHz typical) relative to the superframe rate (48 kHz typical). For 1/2x or 1/4x PDM sampling, synchronous data in an A²B slot is duplicated in order to match the superframe rate. Even lower PDM sampling rates are possible when the reduced rate feature of the transceiver is used in combination with this (for example, down to 375 Hz).

The PDM bit clock output frequency from the transceiver is 64x faster than the PDM audio sampling rate (typically, 3.072 MHz for 48 kHz PDM audio sampling).

Each PDM-enabled receive pin can receive up to two channels of audio data (stereo). One of the channels is associated with the rising edge of the clock and the other with the falling edge of the clock.

The PDM block is configured using the PDM control (`A2B_PDMCTL`) register:

- When `A2B_PDMCTL.PDM0EN = 1`, the DRX0/IO5 pin is enabled to receive PDM data, and the BCLK pin is an output, typically producing a 3.072 MHz clock for the TDM2 setting. In this mode, the DRX0/IO5 pin data is not passed to the I²S/TDM port. Similarly, the `A2B_PDMCTL.PDM1EN` bit controls PDM data reception on the DRX1/IO6 pin.
- The `A2B_PDMCTL.PDMxSLOTS` bits select whether the PDM signals on the DRX pins use one (mono) or two (stereo) channels.

PDM Sampling Edge of a Connected Microphone

The pulse-density modulation (PDM) interface allows PDM input from two microphones to be time-multiplexed on a single data line using a single clock.

A PDM microphone encodes data such that the left channel is valid on the falling edge of the clock (CLK) signal and the right channel is valid on the rising edge of the CLK signal. After the DATA signal is driven during the appropriate half phase of the CLK signal, the microphone output is tristated. As such, two microphones (one set to the left channel and the other set to the right channel) can share a single DATA line (see the *Stereo PDM Format* figure).



Figure 2-9: Stereo PDM Format

In the transceiver, the PDM block samples the microphone data on all 64 clock edges. The transceiver must be programmed to a TDM mode that produces 64 BCLKs per frame (either the default TDM2/32 or TDM4/16 mode). The TDM settings do not affect the PDM block.

In the transceiver, the data sampled on the rising edge of BCLK is always the first channel. If `A2B_PDMCTL.PDM0SLOTS = 1` or `A2B_PDMCTL.PDM1SLOTS = 1`, the first slot is associated with the rising edges of BCLK, and the second slot is associated with the falling edges of BCLK.

For example, two microphones are connected to each of the DRX0/IO5 and DRX1/IO6 pins of a sub node with the PDM0 and PDM1 slots configured as 2-slot. In this case, the PDM block samples 64-bit data each frame, converts it to 24-bit PCM data, and drives the converted output as follows:

- Right microphone data is sampled on the DRX0 pin on rising clock edges and driven in the first* transmit slot on the A²B bus.
- Left microphone data is sampled on the DRX0 pin on falling clock edges and driven in the second* transmit slot on the A²B bus.
- Right microphone data is sampled on the DRX1 pin on rising clock edges and driven in the third* transmit slot on the A²B bus.
- Left microphone data is sampled on the DRX1 pin on falling clock edges and driven in the fourth* transmit slot on the A²B bus.

Note that * is the actual slot number, based on the system slot configuration.

NOTE: When using the default A2B_PDMCTL2 settings, PDM pins are always sampled with rising edge data first; therefore, the A2B_I2SCFG.RXBCLKINV and A2B_I2SCFG.TXBCLKINV clock inversion settings are ignored when the transceiver is configured in PDM mode.

If using the default A2B_PDMCTL2 settings and A2B_PDMCTL.PDM0SLOTS = 0 or A2B_PDMCTL.PDM1SLOTS = 0, only the right channel data is sampled on the PDM pin. If sampling only left channel data is desired, this can be supported by setting A2B_PDMCTL.PDM0EN = A2B_PDMCTL.PDM0SLOTS = A2B_UPOFFSET = 1.

PDM Enhancements

The default PDM functionality is fully backward-compatible with previous transceiver generations; however, there are several additional features which make the PDM interface more flexible.

PDM Clocking Options

The DRX0 and DRX1 input pins can be configured individually as PDM inputs. When the PDM interface is enabled on an A²B sub node on one or both of the DRX pins, a PDMCLK signal running at $64 \times f_{\text{SYNCM}}$ (3.072 MHz at 48 kHz f_{SYNCM}) is required to clock the PDM device. The transceivers allow either the PDMCLK/IO7 or BCLK pin to produce the required PDMCLK. PDMCLK on IO7 can be enabled by setting the A2B_PDMCTL2.PDMALTCLK bit.

If PDMCLK/IO7 is used instead of BCLK, the restriction limiting operating to TDM2/32 or TDM4/16 is removed. The BCLK frequency can be set to a different frequency using the I²S/TDM registers. In this case, PDMCLK/IO7 is used to capture PDM input on DRX0/DRX1.

BCLK and PDMCLK/IO7 can also be used concurrently to clock the PDM microphones at the same frequency and phase alignment, but with opposite polarity. This is accomplished by setting the A2B_PDMCTL2.PDMALTCLK bit. Additionally, a register controls whether the rising edge data or falling edge data is sampled first:

- When `A2B_PDMCTL2.PDM0FFRST = 0` (default), the PDM0 data on DRX0 is sampled rising edge first. When `A2B_PDMCTL2.PDM0FFRST = 1`, it is sampled falling edge first.
- When `A2B_PDMCTL2.PDM1FFRST = 0` (default), the PDM1 data on DRX1 is sampled rising edge first. When `A2B_PDMCTL2.PDM1FFRST = 1`, it is sampled falling edge first.

NOTE: In a main node, BCLK is always an input; therefore, the clock output to PDM microphones connected to a main transceiver typically comes from PDMCLK/IO7.

PDM Data Routing Options

The PDM interface can be used on main or subordinate transceivers. The PDM data received by the transceiver can then be sent to any node on the A²B bus, sent out to the local I²S port, or both. This option is configured using the `A2B_PDMCTL2.PDMDEST` field.

Full-Duplex I²S With Four PDM Microphones

If both pins (DRX0 and DRX1) are used to receive PDM data, it is possible to change the function of DTX1 so that it acts as the alternate DRX1, enabling concurrent use of up to four PDM microphones and full-duplex I²S communications. This is accomplished by setting the `A2B_I2SGCFG.RXONDTX1` bit.

I²S/TDM Interface

The I²S/TDM serial port operates in full-duplex mode, where both the transmitter and receiver operate simultaneously using the same critical timing bit clock (BCLK) and frame synchronization (SYNC) signals. A²B subordinate transceivers generate the timing signals on the BCLK and SYNC output pins with frequencies based on the settings in the I²S global configuration register (`A2B_I2SGCFG`), the I²S rate register (`A2B_I2SRATE`), and the I²S reduced rate register (`A2B_I2SRRATE`). A²B main transceivers use the same BCLK and SYNC pins as inputs. The host drives the pins which provides the time base for the full A²B bus topology.

Time Division Multiplexing (TDM) Protocol

TDM mode extends an I²S interface to more than a stereo 2-channel (TDM2) signal. When the transceiver is programmed in the `A2B_I2SCFG` register to support a certain number of TDM channels, this number of TDM channels is available on each enabled I²S/TDM data pin (DTX0 and DTX1 or DRX0 and DRX1). TDM2, TDM4, TDM8, TDM12, TDM16, TDM20, TDM24, and TDM32 modes are supported.

For example, if TDM4 is selected and one transmit pin (DTX0) is enabled, there are four transmit data channels. If TDM4 is selected and both transmit pins (DTX1 and DTX0) are enabled, there are eight transmit data channels, shown in the *Data Channel Structure for TDM4 Setting* figure.

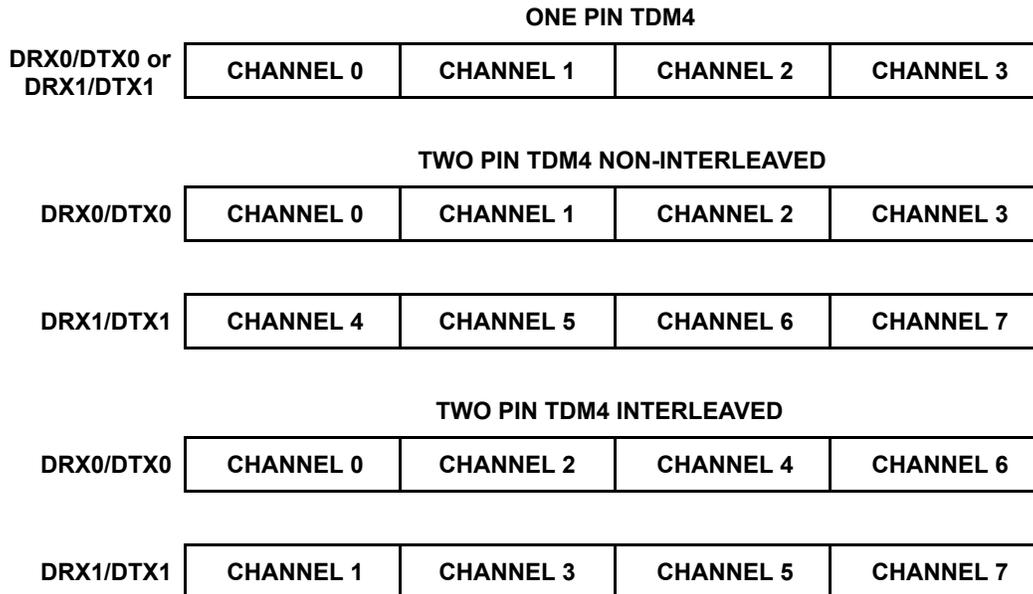


Figure 2-10: Data Channel Structure for TDM4 Setting (TDMMODE == 001)

The I²S/TDM serial port supports data channel widths of 16 bits or 32 bits to carry signals of varying word length. Data words are always represented in MSB first format. The BCLK signal frequencies for different TDM modes are shown in the *I²S/TDM Clock Frequency Settings for 48 kHz Superframe Rates* table.

Table 2-4: I²S/TDM Clock Frequency Settings for 48 kHz Superframe Rates

| TDM Mode | 16-bit TDM Channel Size | | 32-bit TDM Channel Size | |
|----------|-------------------------|---------------------|-------------------------|---------------------|
| | Frequency (MHz) | Comments | Frequency (MHz) | Comments |
| TDM2 | 1.536 | | 3.072 | |
| TDM4 | 3.072 | | 6.144 | |
| TDM8 | 6.144 | | 12.288 | |
| TDM12 | 9.216 | No sub node support | 18.432 | No sub node support |
| TDM16 | 12.288 | | 24.576 | |
| TDM20 | 15.36 | No sub node support | 30.72 | No sub node support |
| TDM24 | 18.432 | No sub node support | 36.864 | No sub node support |
| TDM32 | 24.576 | | 49.152 | |

The DRX0 and DRX1 input pins can be configured individually as PDM inputs. When PDM is enabled on an A²B sub node on one or both of the DRX pins, a PDM clock running at $64 \times f_{\text{SYNCM}}$ (3.072 MHz at 48 kHz f_{SYNCM}) is required to clock the PDM device. Either the PDMCLK/IO7 pin or the BCLK pin can produce the required PDM clock. The transceiver can simultaneously transmit TDM data over the DTX0 or DTX1 pin while receiving PDM streams. However, when BCLK is used as the PDM clock, only I²S/TDM2 and 32-bit channel widths or TDM4 with 16-bit channel widths are supported. Using PDMCLK/IO7 instead of BCLK to clock PDM devices

allows BCLK to be used for a variety of TDM modes. If both DRX0 and DRX1 are used to receive PDM data, it is possible to change the function of DTX1 such that it acts as an alternate DRX1. This enables the concurrent use of up to 4 PDM microphones and full duplex I²S communication.

If using only one pin (DRX0 or DRX1) for PDM, the other pin is available simultaneously for I²S/TDM transfers.

Mailboxes

There are two virtual mailboxes, MBOX0 and MBOX1, that allow for inter-processor communication between the host and a subordinate node control processor.

NOTE: Throughout this section, all specific references to MBOX0 also apply to the MBOX1 instance.

The processor in a subordinate node can send a message over I²C to registers in the A²B subordinate transceiver. In the main node, the host processor is informed about the new message by an interrupt on the main transceiver's IRQ/IO0 pin and can read out the message from A²B subordinate transceiver registers over I²C using the BUS_ADDR. If a mailbox message exchange is from the A²B main node to the A²B subordinate node, the host places a message in A²B subordinate transceiver registers over I²C using the BUS_ADDR. In the subordinate node, the processor is informed of this new message by an interrupt on the subordinate transceiver's IRQ/IO0 pin and can directly read out the message over I²C from A²B subordinate transceiver registers after checking the [A2B_LINTTYPE](#) register.

Mailbox Programming and Operation

The [A2B_MBOX0CTL](#) register provides bit fields to enable the mailbox and control direction, message length, and interrupt capabilities.

By default, mailbox 0 is configured as a receive mailbox (written by the host, read by the subordinate node processor), and mailbox 1 is configured as a transmit mailbox (written by the subordinate node processor, read by the host). Manipulating the [A2B_MBOX0CTL.MB0DIR](#) bit controls the direction of the mailbox.

Each mailbox can hold either 8-, 16-, 24-, or 32-bit messages, as configured in the [A2B_MBOX0CTL.MB0LEN](#) field. The value in this field determines which of the four byte-wide [A2B_MBOX0B0](#) through [A2B_MBOX0B3](#) registers to use for the data, where the first byte is always in the [A2B_MBOX0B0](#) register, and the final byte is in the highest data register required to accommodate the programmed data length, as shown in the following table.

| MBxLEN Field | Final Byte in Register |
|--------------|-----------------------------|
| 0b00 | A2B_MBOX0B0 |
| 0b01 | A2B_MBOX0B1 |
| 0b10 | A2B_MBOX0B2 |
| 0b11 | A2B_MBOX0B3 |

For an enabled receive mailbox ([A2B_MBOX0CTL.MB0EN](#) = 1 and [A2B_MBOX0CTL.MB0DIR](#) = 0), if the [A2B_MBOX0CTL.MB0FIEN](#) bit is set, an interrupt to the subordinate node occurs after the final byte of the mailbox is written by the host and received by the A²B subordinate transceiver. If the [A2B_MBOX0CTL.MB0EIEN](#) bit

is set, an interrupt is propagated back upstream over the A²B bus to the host after the final byte of the mailbox is read by the local processor in the subordinate node.

For an enabled transmit mailbox (`A2B_MBOX0CTL.MB0EN = 1` and `A2B_MBOX0CTL.MB0DIR = 1`), if the `A2B_MBOX0CTL.MB0FIEN` bit is set, an interrupt to the host occurs after the final byte of the mailbox is written by the local processor in the subordinate node. If the `A2B_MBOX0CTL.MB0EIEN` bit is set, an interrupt is propagated downstream over the A²B bus to the subordinate node after the final byte of the mailbox is read by the host.

CAUTION: Dynamic reconfiguration of an enabled mailbox (`A2B_MBOX0CTL.MB0EN = 1`) is forbidden. The host must first disable the mailbox (`A2B_MBOX0CTL.MB0EN = 0`) and then re-enable it in two separate accesses if reconfiguration is required.

The `A2B_MBOX0STAT` register provides status information for the mailboxes:

- When a mailbox is filled, the `A2B_MBOX0STAT.MB0FULL` bit is set, and the `A2B_MBOX0STAT.MB0EMPTY` bit is cleared.
- When a mailbox is emptied, the `A2B_MBOX0STAT.MB0EMPTY` bit is set, and the `A2B_MBOX0STAT.MB0FULL` bit is cleared.
- The `A2B_MBOX0STAT.MB0EIRQ` and `A2B_MBOX0STAT.MB0FIRQ` bits are set when the mailbox signals an interrupt to the host or local processor, and the bits are cleared when the interrupt is processed by the host or local processor.

Multiple subordinate nodes can communicate to the main node through their TX mailboxes. In the main node, the `A2B_INTTYPE` register contains information about the pending interrupt generated by any subordinate node, with the subordinate node indicated in the `A2B_INTSRC` register.

When two subordinates write to their mailboxes simultaneously, the main gets the interrupt indication from the subordinate that is closer to the main. Upon detecting the interrupt, the host extracts the interrupt information by reading the A²B main transceiver's interrupt type (`A2B_INTTYPE`) and interrupt source (`A2B_INTSRC`) registers to determine which interrupt occurred and which subordinate node generated it, respectively. Upon reading the `A2B_INTTYPE` register, the interrupt request for that interrupt is cleared in the subordinate node identified by the value in the `A2B_INTSRC` register. The IRQ/IO0 pin toggles to the deasserted state and then immediately back into the asserted state due to the still active interrupt from the other subordinate node, and the host can again read the main transceiver's `A2B_INTTYPE` and `A2B_INTSRC` registers to acknowledge the other subordinate node's mailbox interrupt.

Mailbox Latency

The mailbox transactions are made up of register reads and writes over the I²C bus. The interrupt request from a subordinate to the main transceiver is part of the SRF packet, so the latency on the subordinate to main transceiver mailbox can include an extra superframe waiting for this time.

The following figures show the system timing for the mailbox transactions in both directions. The light gray slots indicate the SCF field, and the dark gray slots indicate the SRF field.

As shown in the *Mailbox Latency (from Host to Subordinate)* figure, when the mailbox message is from the host to a subordinate processor, the host processor writes the mailbox data to the A²B subordinate node through the SCF field using a 2-byte burst write access to the main transceiver BUS_ADDR device address. When the writes complete, the subordinate transceiver immediately generates the interrupt to its local node processor. As a result, the subordinate interrupt request (SUBORDINATE IRQ) asserted on IRQ/IO0 aligns with the SCF field. Once this interrupt is asserted, the locally-connected processor can use the subordinate transceiver BASE_ADDR device address to interrogate the A²B_LINTTYPE register to determine that it is the mailbox full interrupt, after which it can then extract the data from the mailbox data registers using a 2-byte burst read. Once those transactions finish, the mailbox empty interrupt is generated at the main node (MAIN IRQ), aligned with the SRF field, and the host proceeds with reading the A²B_INTSRC and A²B_INTTYPE registers of the main transceiver (using the main transceiver BASE_ADDR device address) to determine that it is the mailbox empty interrupt originating with the indicated subordinate.

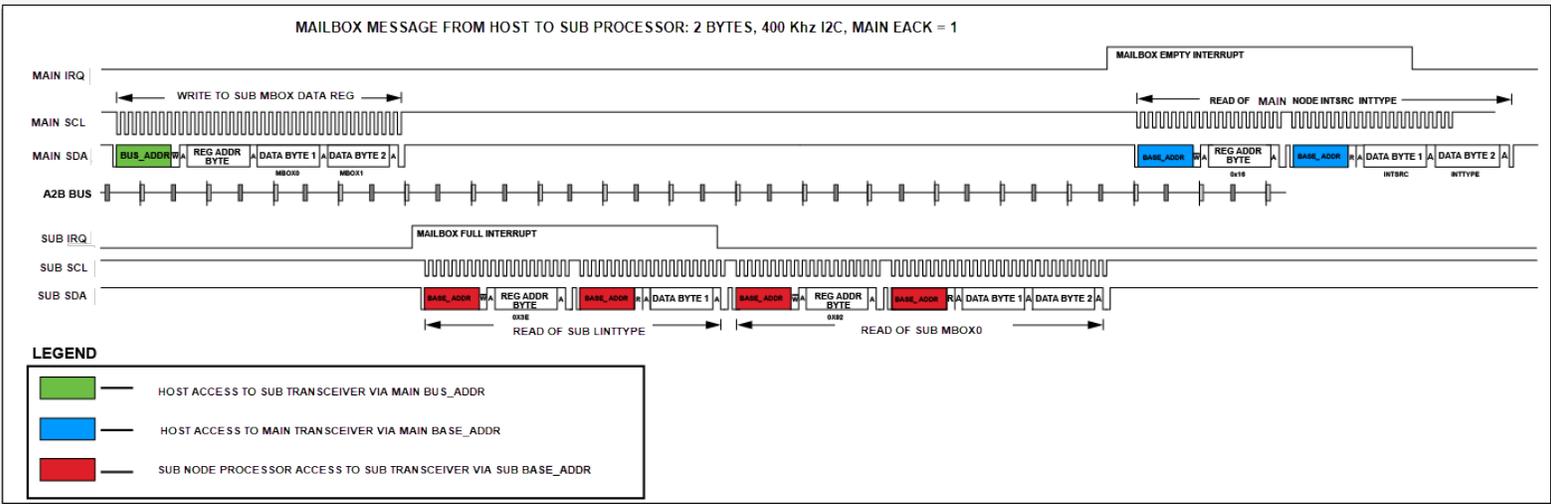


Figure 2-11: Mailbox Latency (from Host to Subordinate)

Similarly, as shown in the *Mailbox Latency (from Subordinate to Host)* figure, when the mailbox message is from a subordinate processor to the host, the subordinate node processor populates the mailbox data registers at any time by issuing writes to the registers using the subordinate transceiver BASE_ADDR device address, and the interrupt indication to the main A²B node goes through the SRF field. As a result, the main mailbox full interrupt request (MAIN IRQ) asserted on IRQ/IO0 aligns with the SRF field. Once this interrupt is asserted, the host (using the main transceiver BASE_ADDR device address) interrogates the A²B_INTSRC and A²B_INTTYPE registers to determine that it is the mailbox full interrupt originating with the indicated subordinate.

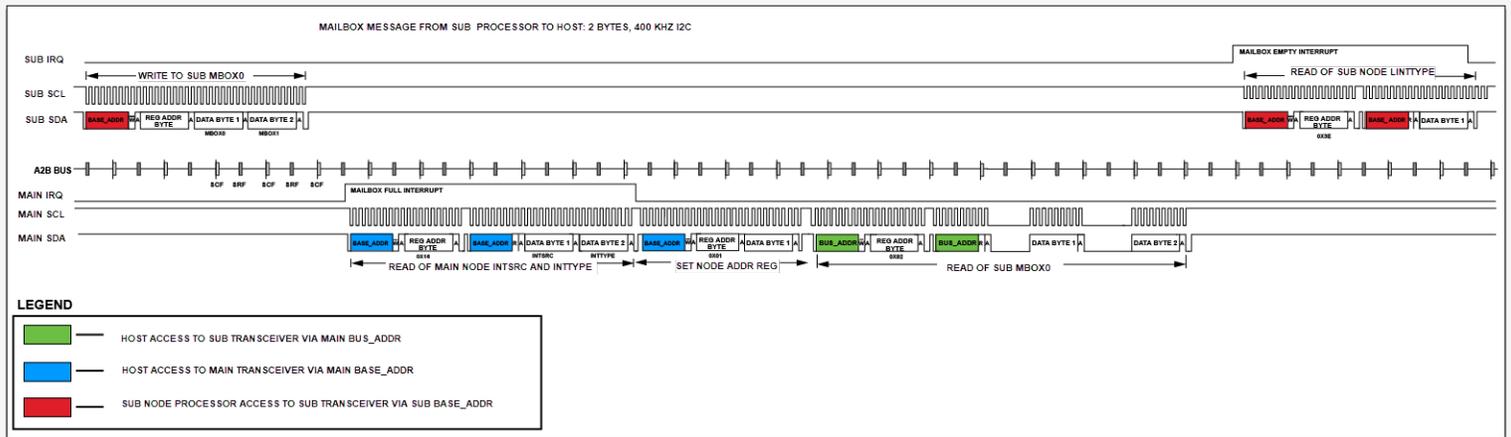


Figure 2-12: Mailbox Latency (from Subordinate to Host)

To subsequently extract the data from the mailbox of the subordinate transceiver, the host must first set the `A2B_NODEADR` register to the subordinate node that generated the interrupt (using a main transceiver `BASE_ADDR` write access), and then issue the `BUS_ADDR` accesses to read the mailbox data byte registers of the subordinate transceiver (note the superframe spacing required for these reads to take place). Once the last byte is read by the host, the mailbox empty interrupt request of the subordinate node (SUBORDINATE IRQ) gets asserted in the next SCF. Then, the subordinate node processor can use a subordinate transceiver `BASE_ADDR` access to read the `A2B_LINTTYPE` register and take action after identifying that it was the mailbox empty interrupt that occurred (for example, load the mailbox data registers again to restart the process).

3 A²B Operation and Configuration

The A²B bus is high-level programmable and can address many use cases. A²B systems are easy to configure, based on knowledge of the system, nodes, and peripherals. The exact system configuration can be gained by collecting information individually from each subordinate. As an example, the same A²B module can be supplied by different vendors, with each of the modules having unique register programming requirements. One module can use TDM4 as an audio interface, while another one uses TDM8. One module can provide two upstream channels, while another can provide three upstream channels, all with the host not having prior knowledge of how many nodes are connected.

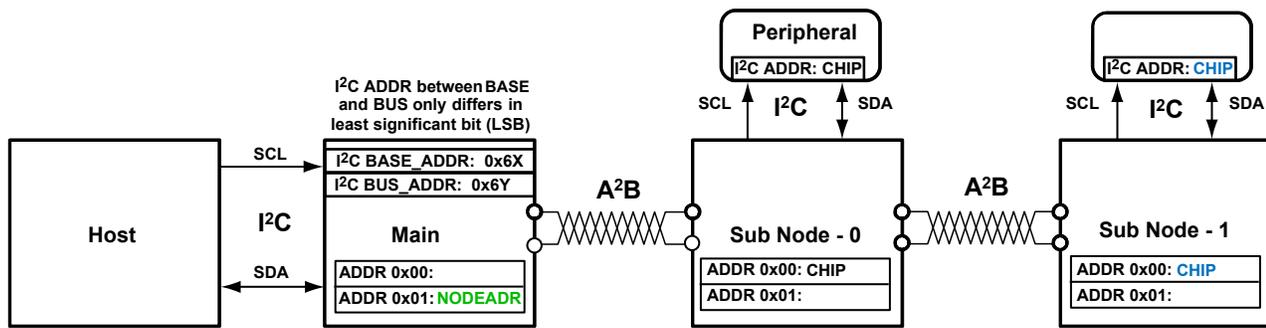
IMPORTANT: Ensure that the register programming results in a valid system configuration.

Analog Devices provides free SigmaStudio™ (<http://www.analog.com/SigmaStudio>) tools featuring an intuitive graphical user interface to architect, configure, and set up the A²B bus. The tools also generate driver code for embedded software.

Linux and QNX software drivers also are available upon request.

I²C Port Programming Concepts

Main transceiver enabled registers are programmed directly by the A²B host via the I²C port using [Direct I²C Register Accesses](#). Transceiver sub node enabled registers can also be programmed in this fashion by an I²C-connected controller on the subordinate node; however, A²B subordinate transceiver registers are typically programmed remotely by the A²B host through the main transceiver over the A²B bus using [Remote Target I²C Register Accesses](#). Further, if a subordinate transceiver is locally connected to an I²C target device on the subordinate node, that connected I²C target device can also be accessed remotely by the A²B host over the A²B bus using [Remote Peripheral I²C Accesses](#). The *Programming Sequence for I²C Accesses* figure is a graphical representation of the programming sequences that are required when programming transceiver registers and accessing target node I²C peripheral devices.



Main I²C Access:

<I2C ADDR : BASE > R/W < ADDR > < R/W Data > // Read and write directly from/to main node

Sub Node I²C Access:

<I2C ADDR : BASE > R/W < ADDR : 0x01 > < PERI=0, NODE=0 > // Set sub node number in main node (Sub Node 0)

<I2C ADDR : BUS > R/W < ADDR > < Data > // Read and write directly from/to sub node

Peripheral of Sub Node I²C Access:

<I2C ADDR : BASE > R/W < ADDR : 0x01 > < PERI=0, NODE=1 > // Set sub node number in main (Sub Node 1)

<I2C ADDR : BUS > R/W < ADDR : 0x00 > < CHIP > // set device address for peripheral (CHIP) in sub node

<I2C ADDR : BASE > R/W < ADDR : 0x01 > < PERI=1, NODE=1 > // Set main node to access a sub node's peripheral (Peripheral of Sub Node 1)

<I2C ADDR : BUS > R/W < ADDR > < Data > // Read and write directly from/to peripheral of sub node

Figure 3-1: Programming Sequence for I²C Accesses

In the *Programming Sequence for I²C Accesses* figure:

- I2C ADDR is the controller transceiver I²C device address:
 - [Direct I²C Register Accesses](#) to the main transceiver use BASE_ADDR (I2C ADDR: BASE).
 - [Remote Target I²C Register Accesses](#) to a subordinate transceiver and [Remote Peripheral I²C Accesses](#) to an I²C-connected peripheral on a subordinate node use BUS_ADDR (I2C ADDR: BUS).

NOTE: See [Transceiver I²C Accesses](#) for more details regarding BASE_ADDR and BUS_ADDR.

- NODEADR is the main transceiver A2B_NODEADR register:
 - NODE is the A2B_NODEADR.NODE field.
 - PERI is the A2B_NODEADR.PERI bit.
- CHIP is the A2B_CHIP register:
 - Black text indicates the A2B_CHIP register itself.
 - Blue text indicates the value of the A2B_CHIP register.

Direct I²C Register Accesses

The I²C port can be used to directly access the transceiver register space, whether the transceiver is configured as a main or as a subordinate:

- On the main node, the A²B host directly accesses the main transceiver register space using this method.
- On a subordinate node, a locally-connected I²C host directly accesses the subordinate transceiver register space using this method.

As shown in the *Main I²C Access* portion of the *Programming Sequence for I²C Accesses* figure, a main transceiver register access requires the I²C transfer from the host to consist of the main transceiver I²C device address (I2C ADDR: BASE = BASE_ADDR), followed by the register address (ADDR), followed finally by the data associated with the main transceiver register (R/W Data). For further details, see [Transceiver I²C Accesses](#).

NOTE: This *Main I²C Access* sequence is identical for an I²C-connected host on the subordinate node directly accessing a subordinate transceiver's register space.

Remote Target I²C Register Accesses

Though a locally-connected I²C host on a subordinate node can directly program subordinate transceiver registers over the I²C port, A²B systems are typically fully configured by the A²B host from the main node. As shown in the *Subordinate 0 I²C Access* portion of the *Programming Sequence for I²C Accesses* figure, the 2-step process consists of the A²B host first directly configuring the main transceiver before using remote I²C accesses to program a specific subordinate transceiver over the A²B bus. The A²B host must use the following programming sequence to access an A²B subordinate transceiver register space remotely over the A²B bus from the main node.

1. Use a [Direct I²C Register Accesses](#) to set the main transceiver A2B_NODEADR.NODE field to the subordinate node ID to be accessed. Be sure the A2B_NODEADR.PERI bit is set to 0 in this write so that subsequent bus accesses target the indicated subordinate transceiver register space rather than an I²C peripheral connected to the indicated subordinate.

ADDITIONAL INFORMATION: Setting the A2B_NODEADR.NODE field to 0 means that subsequent bus accesses will target subordinate node 0. If this field were set to 1, subsequent bus accesses would target subordinate node 1. If the intent is to broadcast the write to all of the discovered nodes (main and subordinates), be sure to also set the broadcast bit (A2B_NODEADR.BRCST) in this write.

2. To access the subordinate transceiver register, the I²C transfer from the host consists of the main transceiver's bus address (I2C ADDR: BUS = BUS_ADDR), followed by the subordinate transceiver register address (ADDR), followed finally by the data associated with the subordinate transceiver register (Data). For more details, see [Transceiver I²C Accesses](#).

Remote Peripheral I²C Accesses

The *Peripheral of Target 1 I²C Access* portion of the *Programming Sequence for I²C Accesses* figure illustrates the sequence required for the A²B host to access a peripheral connected to the I²C port of a subordinate transceiver over

the A²B bus using remote peripheral I²C accesses. The A²B host must follow the below programming sequence to access an I²C peripheral on an A²B subordinate node (for example, a microphone or a DAC) over the A²B bus.

1. Use a [Direct I²C Register Accesses](#) write access to set the main transceiver A2B_NODEADR.NODE field to the subordinate node ID that is connected to the peripheral to be accessed. Be sure the A2B_NODEADR.PERI bit is cleared in this write so that the subsequent bus access is to the targeted subordinate transceiver's register space, not to the subordinate peripheral itself.

ADDITIONAL INFORMATION: The A2B_NODEADR.NODE field is set to 1 in this write so that subsequent bus accesses target subordinate node 1. If the intent is to broadcast the peripheral write to all of the discovered nodes (main and subordinates), be sure to also set the A2B_NODEADR.BRCST bit in this write. If the A2B_CHIP register in the targeted subordinate transceiver is already set to the I²C address of the intended peripheral access, perform this write with the A2B_NODEADR.PERI bit set (rather than cleared) and proceed directly to the final step.

2. Use a [Remote Target I²C Register Accesses](#) write access to program the desired subordinate transceiver's A2B_CHIP register with the I²C device address of the peripheral connected to the subordinate.
3. Use a [Direct I²C Register Accesses](#) write access to set the main transceiver A2B_NODEADR.PERI bit (while maintaining the content of the A2B_NODEADR.NODE field) such that subsequent BUS_ADDR accesses go to the desired subordinate node I²C peripheral.
4. To access the subordinate node peripheral, the I²C transfer from the host must consist of the main transceiver's BUS_ADDR (I2C ADDR: BUS), followed by the address that the subordinate transceiver uses to access the subordinate node I²C peripheral (ADDR), followed finally by the data associated with the address (Data).

System Bring-Up and Discovery

An A²B system is brought up by the A²B host. Once power is properly established, each node in the system must be discovered and configured in order, starting with the main node.

Reset and Operating States

Loss of PLL lock resets all of the register information except A2B_BMMCFG and A2B_CONTROL.MSTR.

The *Transceiver State Diagram* figure shows transceiver state information that is important to understand when bringing up and running a complete A²B system.

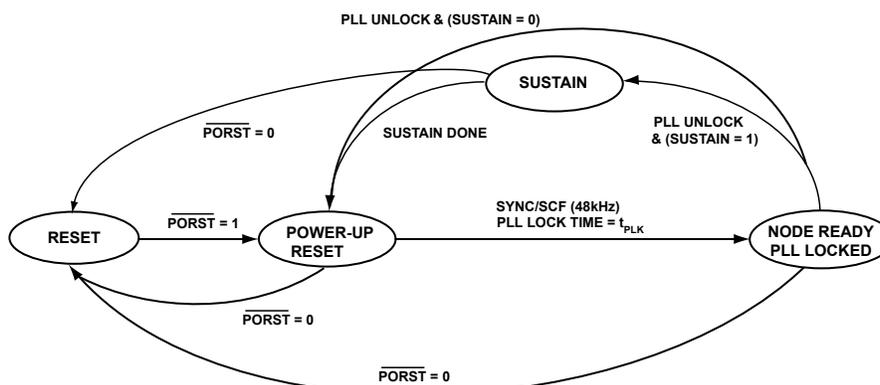


Figure 3-2: Transceiver State Diagram

NOTE: As sustain mode is a subordinate-only feature, a main transceiver never enters the SUSTAIN state. A loss of PLL lock on a main transceiver results in a direct return to the POWER-UP state.

Transceiver Power-On and Reset

When the transceiver is initially being powered on, it is in the RESET state. When in RESET, all A²B system registers are held in reset, and no registers can be programmed until the transceiver advances to the POWER-UP RESET state, which is a function of power (VIN) applied to the transceiver.

An internal power-on reset circuit monitoring the state of the VIN power supply pin holds an internal power-on reset signal ($\overline{\text{PORST}}$) asserted low until the V_{RSTN} specification is met, at which point $\overline{\text{PORST}}$ is deasserted high to indicate that the transceiver is properly powered. The transceiver then transitions to the POWER-UP RESET state. After the transceiver enters POWER-UP RESET, the $\overline{\text{PORST}}$ signal remains deasserted high unless the voltage sensed on the VIN power supply pin drops into its V_{RST} specification range, in which case $\overline{\text{PORST}}$ is asserted low to bring the transceiver back to RESET.

Main Node Bring-Up and Operation

Referring to the *Transceiver State Diagram* figure, the ADR1/IO1 and ADR2/IO2 pins are latched to determine the I²C device address when the transceiver transitions to the POWER-UP RESET state, and the transceiver is I²C-device accessible no more than 2.5 ms later. The A²B host then sets the A2B_CONTROL.MSTR bit in the transceiver before driving the SYNC pin (the input clock to a main transceiver-enabled transceiver's PLL) at the audio sampling rate of the system (selectable between 48 kHz or 44.1 kHz). The main transceiver locks its PLL to the received SYNC signal according to the PLL Lock Time (t_{PLK}) specification.

NOTE: It is recommended that the host set a timeout in excess of the PLL Lock Time (t_{PLK}) specification so that a non-responsive main transceiver can be detected by software.

Upon PLL lock, the main transceiver transitions to the NODE READY PLL LOCKED state, at which point it generates the MSTR_RUNNING (0xFF) interrupt to the host (the IRQ/IO0 pin is driven high), as stored in the

interrupt type register ([A2B_INTTYPE](#)), indicating the main transceiver is ready for programming via the I²C interface.

NOTE: Once the PLL is locked, writing the `A2B_CONTROL.MSTR` bit has no effect.

If the main transceiver PLL becomes unlocked during bus operation, the transceiver goes back to POWER-UP, as the SUSTAIN state is a subordinate-only feature (SUSTAIN = 0). All of the registers return to their reset values except the [A2B_CONTROL](#) register.

Subordinate Node Bring-Up and Operation

Referring to the *Transceiver State Diagram*, the transceiver is in the POWER-UP RESET state once local or A²B bus power is established, and the ADR1/IO1 and ADR2/IO2 pins are latched to determine the I²C device address. The transceiver is by default a subordinate and is ready to be discovered and programmed 2.5 ms after entering POWER-UP RESET.

NOTE: The BCLK and SYNC outputs are three-stated in the POWER-UP RESET state.

While in the POWER-UP RESET state, a subset of the subordinate transceiver register space can be configured through the I²C port by a locally-connected host using [Direct I²C Register Accesses](#). These registers include:

| | | |
|---|---|---|
| A2B_BMMCFG | A2B_CHIP | A2B_BCDNSLOTS |
| A2B_LDNSLOTS | A2B_LUPSLOTS | A2B_DNSLOTS |
| A2B_UPSLOTS | A2B_INTMSK0 | A2B_INTMSK1 |
| A2B_BECCTL | A2B_TESTMODE | A2B_I2CCFG |
| A2B_SYNCOFFSET | A2B_PDMCTL | A2B_ERRMGMT |
| A2B_GPIODAT | A2B_GPIOOEN | A2B_GPIOIEN |
| A2B_PINTEN | A2B_PINTINV | A2B_PINCFG |
| A2B_I2SRATE | A2B_I2SRRCTL | A2B_I2SRRSOFFS |
| A2B_CLK1CFG | A2B_CLK2CFG | A2B_UPMASK0 - A2B_UPMASK3 |
| A2B_UPOFFSET | A2B_DNMASK0 - A2B_DNMASK3 | A2B_DNOFFSET |
| A2B_GPIODEN | A2B_GPIOD0MSK - A2B_GPIOD7MSK | A2B_GPIODINV |
| A2B_MBOX0CTL - A2B_MBOX1CTL | A2B_I2STEST | A2B_I2SRATE |
| A2B_I2SGCFG | A2B_I2SCFG | |

Even though these registers can be written in the POWER-UP RESET state, programmed values do not take effect until the transceiver advances to the NODE READY PLL LOCKED state, with the exception of the slot registers ([A2B_BCDNSLOTS](#), [A2B_LDNSLOTS](#), [A2B_LUPSLOTS](#), [A2B_DNSLOTS](#), [A2B_UPSLOTS](#), [A2B_UPMASK0](#) through [A2B_UPMASK3](#), and [A2B_DNMASK0](#) through [A2B_DNMASK3](#)). Values programmed to the listed slot registers do not take effect until the main transceiver [A2B_DATCTL](#) register is programmed and the new structure is subsequently applied (`A2B_CONTROL.NEWSTRCT = 1`).

In the POWER-UP RESET state, a subordinate transceiver awaits synchronization control frames (SCFs) coming from the A²B main, which is initiated when the host initiates the discovery process for that particular A²B system subordinate by setting the main transceiver's `A2B_DISCVRY` register to the response time for the targeted subordinate. When this write occurs, the main initiates discovery by sending discovery frames with the response time value embedded in the SCF. The subordinate being discovered then extracts the information to set its response time (`A2B_RESPCYCS`). These discovery frames provide the input clock to the subordinate transceiver, which the subordinate transceiver locks its PLL to in accordance with the PLL Lock Time (t_{PLK}) specification. Once the subordinate transceiver PLL is locked, it is in the NODE READY PLL LOCKED state and starts generating synchronization response frames (SRFs) to upstream nodes, which causes the main transceiver to generate the DSCDONE interrupt (`A2B_INTTYPE = 0x18`) indicating that the subordinate transceiver is ready for programming over the A²B bus using [Remote Target I²C Register Accesses](#).

ATTENTION: When simultaneous writes to the same register are attempted from both the A²B bus (using [Remote Target I²C Register Accesses](#)) and the I²C port (using [Direct I²C Register Accesses](#)), the order in which these I²C accesses occur cannot be predicted. Therefore, special care must be taken when I²C transactions are also coming from both sources.

TIP: If local node programming (using [Direct I²C Register Accesses](#)) is desired in the application, this potential contention can be avoided by using a mailbox handshake with the main node such that the host writes one of the subordinate's mailboxes when it is ready to begin making register accesses and then waits for the subordinate to read that mailbox as an indication that its initialization sequence is complete. See [Mailboxes](#) for more information.

In the NODE READY PLL LOCKED state, the BCLK and SYNC output are driven low until any I2S/TDM/PDM port data pin is enabled in the subordinate transceiver's `A2B_PDMCTL` (for PDM mode) or `A2B_I2SCFG` (for I²S/TDM modes) registers.

If the subordinate transceiver PLL becomes unlocked during bus operation, it goes back to the POWER-UP state if the clock sustain feature is disabled (`A2B_SUSCFG.SUSDIS = 1`, denoted in the *Transceiver State Diagram* figure as SUSTAIN = 0). Once back in the RESET state, the A²B main can issue another discovery sequence.

Clock Sustain Functionality

By default (and denoted in the *Transceiver State Diagram* figure as SUSTAIN = 1), the subordinate transceiver has a clock sustain feature to power down subordinate nodes with processors and DACs, where audio signals of locally powered subordinate nodes are gracefully muted. When the bus loses communication and a reliable clock cannot be recovered by the subordinate transceiver (PLL UNLOCK in the *Transceiver State Diagram* figure), the subordinate transceiver enters the SUSTAIN state, provided the clock sustain feature has not been disabled (`A2B_SUSCFG.SUSDIS = 1`). Upon entering the SUSTAIN state, the transceiver:

- Runs at the current clock frequency for 1024 SYNC periods
 - I²S/TDM ports continue running
 - Signals SUSTAIN state on a GPIO, if enabled

- PLL relock is not attempted while in the SUSTAIN state
- Resets and re-enters the POWER-UP RESET state
- Transitions to the NODE READY PLL LOCKED state if stable SCF discovery frames are present

If the sustain GPIO output enable (A2B_SUSCFG.SUSOE) bit is set, the sustain signal from the PLL is driven high on the GPIO pin selected by the A2B_SUSCFG.SUSSEL bit field while the transceiver is in the SUSTAIN state. This feature has a higher priority than other GPIO outputs, but a lower priority than function outputs on the pins. For example, if clock output 1 is enabled (A2B_CLK1CFG.CLK1EN = 1), the ADR1/IO1 pin is driven as a clock output. Setting the A2B_SUSCFG.SUSOE bit and configuring the sustain output on the ADR1/IO1 pin (A2B_SUSCFG.SUSSEL = 1) does not override this behavior.

The sustain signal from the PLL goes high near the beginning of a superframe. Once the sustain signal is high, decaying data values are produced on the DTX0/IO3 and DTX1/IO4 pins, starting on the following I²S/TDM frame.

The data in the TX frame buffer (see [Managing A²B System Data Flow](#)), as received from the A²B bus, contains a 32-bit value that is output to either or both of the I²S DTX0/DTX1 data pins. Negative values gradually attenuate to 0, while positive values gradually attenuate to -109 dB (0x00001F00) on the enabled data pins.

Node Discovery and Initialization

This section provides information regarding simple node discovery and initialization for an A²B bus system. Modified, optimized, and advanced discovery flows are described in [Appendix A: Additional Discovery Flow Examples](#) of this manual. Any of these software flow diagrams can be used as a guideline for discovery and initialization.

Simple Discovery Flow

All subordinate nodes are discovered sequentially from sub 0 to the last available sub in the system with the software flow shown in the *Simple Discovery Flow* figure. In this figure, the stages show commands as issued over the I²C interface between the host and the main-enabled transceiver. Write commands are identified as "wr" and read commands are identified as "rd" along with the REGISTER_NAME being accessed. The "M" indicates an access to the BASE_ADDR, and the "S" indicates an access to the BUS_ADDR.

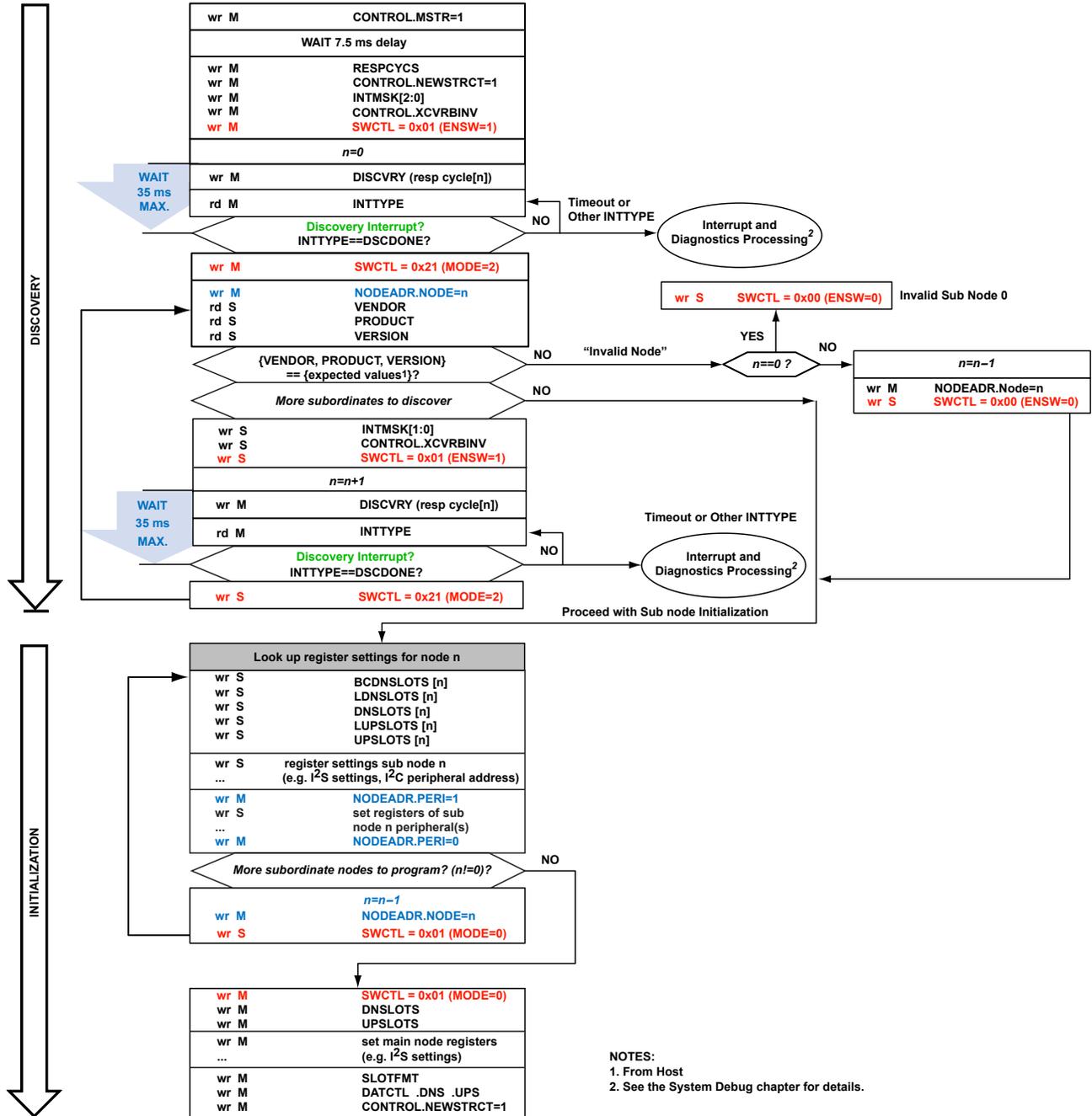


Figure 3-3: Simple Discovery Flow

NOTE: In the *Simple Discovery Flow* figure, setting the `A2B_SWCTL.ENS` bit in the main node or in any sub node causes it to begin sending SCFs downstream to the next connected sub node, thus allowing that next sub transceiver to begin locking its PLL before the main node initiates discovery frames targeting it.

Use the following guidelines for the reverse-wire feature `A2B_CONTROL.XCVRBINV` (Invert Data to/from LVDS XCVR B):

1. In the main node, set the `A2B_CONTROL.XCVRBINV` bit prior to writing to the `A2B_SWCTL.ENS` bit. Be careful to avoid inadvertently clearing the `A2B_CONTROL.XCVRBINV` bit when writing to the `A2B_CONTROL` register for other purposes, such as writing to the `A2B_CONTROL.NEWSTRCT` bit.
2. In any sub node, the `A2B_CONTROL.XCVRBINV` bit must be set before writing to the `A2B_SWCTL.ENS` bit.

Once all of the sub nodes are discovered, initialize the nodes for synchronous data exchange. The example flow diagram starts initialization with the last node and finishes with the main node.

The discovery finishes quickly, providing earlier access to all nodes and their I²C peripherals before the initialization for synchronous audio, which takes extra time to finish.

There is no further need for bus management after all of the nodes are discovered and programmed. Interrupt service routines can be used to react to special interrupt request (IRQ) events (for example, from an IO pin). Alternatively, the `A2B_INTTYPE` register can be polled to monitor interrupt events.

The [Optimized Discovery Flow](#) and [Advanced Discovery Flow](#) sections illustrate how to perform auto-configuration.

Response Cycles

The `A2B_RESPCYCS` register sets the relative time from the start of a synchronization control frame (SCF) to the moment the last subordinate responds with a synchronization response frame (SRF). The register setting indicates to earlier nodes when to expect the response from the last subordinate. If the last node does not respond, the previous node that is next to the presumed last node does respond.

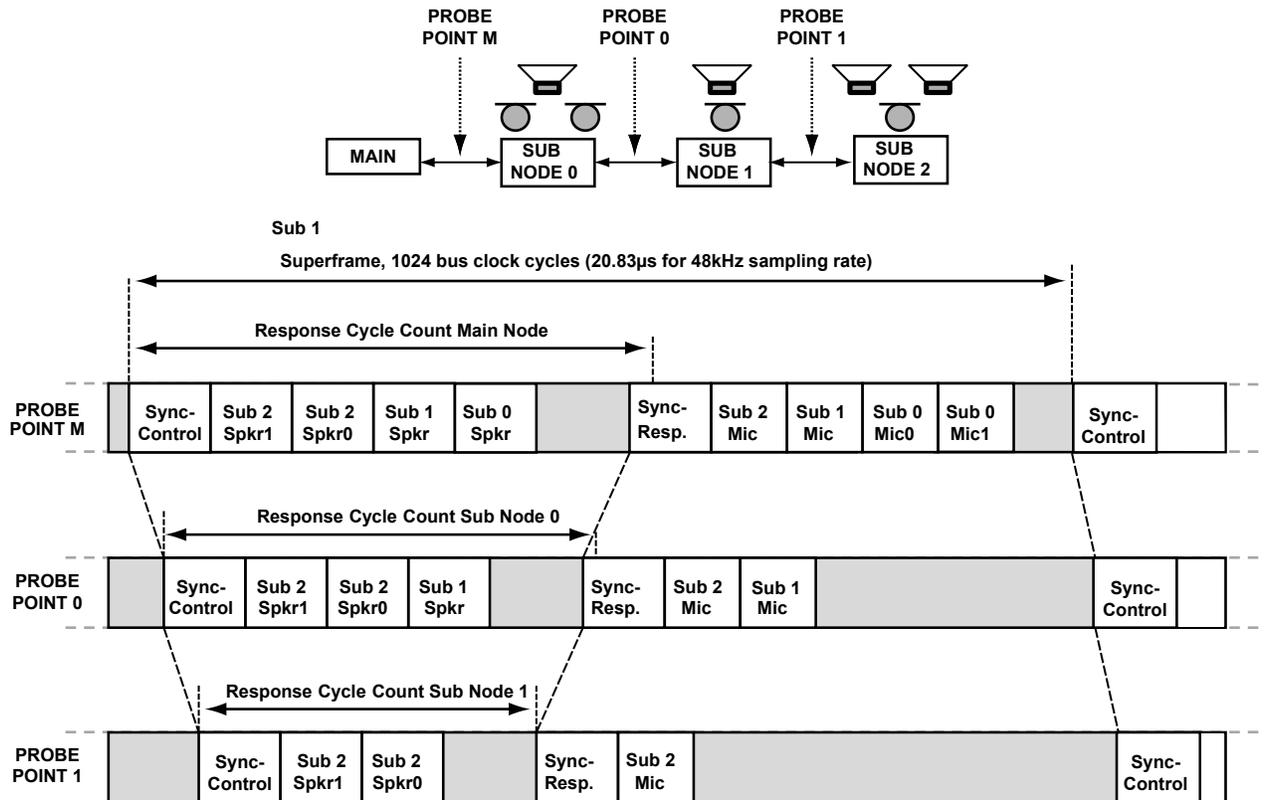


Figure 3-4: Synchronous Data Enabled

The response cycle values for the transceivers are discussed in [Appendix B: Response Cycle Formula](#) as a function of the following parameters:

- Number of subordinate nodes
- Number of downstream slots
- Downstream slot size
- Number of upstream slots
- Upstream slot size
- Main I²S/TDM channel configuration

NOTE: The main transceiver response cycle values are calculated using the above parameters in the response cycle calculator spreadsheet or in SigmaStudio software. For more information, contact your local Analog Devices representative.

Subordinate Node Response Cycles

The *Subordinate Node Response Cycle* figure shows the relative timing between SCFs and SRFs on the A and B XCVR ports of a subordinate node. A subordinate node generates the SRF approximately $((4 * A2B_RESPCYCS)$

+ 7) bits after the SCF starts on the A XCVR. For example, when `A2B_RESPCYCS= 128 (0x80)`, the subordinate node generates the SRF beginning at the 519th $((4 * 128) + 7 = 519)$ bit.

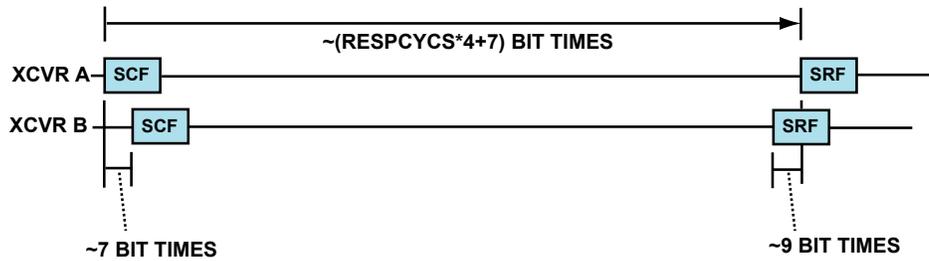


Figure 3-5: Subordinate Node Response Cycle

As shown in the *Subordinate Node Response Cycle* figure, there are transceiver delays (TD) incurred to pass the superframe from one side of the transceiver to the other. For the downstream portion of the superframe, there is a delay (TD_{DOWN}) of seven (± 2) bits incurred when going from the A-side to the B-side of the transceiver. Conversely, there is a delay (TD_{UP}) of nine (± 2) bits going from the B-side to the A-side during the upstream portion of the same superframe. These delays are summarized for the supported frame rates in the *Transceiver Delays* table, as governed by the equation:

$$\text{Delay Range} = \text{Nominal Latency Range} / (\text{SYNC Rate} * 1024)$$

Table 3-1: Transceiver Delays

| Time Delay (Direction) | SYNC Rate (kHz) | Nominal Latency Range (SYSBCLK) | Delay Range (ns) |
|---|-----------------|---------------------------------|------------------|
| TD_{DOWN} (A-Side to B-Side Downstream) | 44.1 | 7 ± 2 | 110.7 - 199.3 |
| TD_{DOWN} (A-Side to B-Side Downstream) | 48.0 | 7 ± 2 | 101.7 - 183.1 |
| TD_{UP} (B-Side to A-Side Upstream) | 44.1 | 9 ± 2 | 155.0 - 243.6 |
| TD_{UP} (B-Side to A-Side Upstream) | 48.0 | 9 ± 2 | 142.4 - 223.8 |

In addition to these transceiver delays, cable delays (CD) between nodes also change the relative timing between when the SCF is received in the downstream portion of the superframe and when the complementary SRF returns to that point during the upstream portion of the same superframe. There is a 5-bit time window (expected bit time \pm

2) in which the SRF is correctly received on the B-side and passed to the A-side of a subordinate node. An SRF outside of this window is still detected, and the expected response time is gradually (and automatically) adjusted by the transceiver during discovery to compensate for mismatches, with an adjustment range of -4 bit times to +15 bit times to span the cable length specifications. As such, the `A2B_RESPCYCS` formula works for all supported cable lengths. If the cable length is known during the system design phase, this recommendation can be applied for all discovery flows. If the cable lengths are unknown, the default response cycles calculation (assuming 4m cable length) is adequate. Although some errors can be observed during discovery (CRCERR, SRFERR, or SRFRCRERR) when longer cables are used, the system runs cleanly after discovery completes due to this automatic adjustment capability.

The automatic response cycle adjustment performed during discovery works as follows:

1. The host programs the main to expect the SRF at the 519th bit of the superframe by setting `A2B_RESPCYCS = 128 (0x80)`, as detailed above ($(4 * 128) + 7 = 519$).
2. The main node initiates discovery of subordinate 0 when the host writes `0x80` to its `A2B_DISCVRY` register. When subordinate 0 starts sending SRFs, the main adjusts its response time to align with subordinate 0.

Short cable lengths (up to 20cm) do not impact the main node's ability to receive the SRF at the 519th bit of the superframe.

Longer cable lengths, however, introduce a physical cable delay (CD) on the order of 5ns/m to the time at which the SRF is captured at the receiving node. For example, a 10m cable between the main node and the subordinate 0 node delays the SRF reception time at the main node by 100ns (50ns downstream CD plus 50ns upstream CD). This 100ns total CD equates to five A²B bits, thus causing the main node in this case to adjust its response cycles to expect the SRF at the 524th (± 2) bit of the superframe.

3. The main node initiates discovery of subordinate 1 when the host writes `0x7C` to the `A2B_DISCVRY` register. When subordinate 1 starts sending SRFs, subordinate 0 adjusts its response time to align with subordinate 1, which causes the SRFs from subordinate 0 to be delayed, thus adding further delay to the time at which the SRF reaches the main node.

The main node receives the SRF as a function of the CD between the main node and subordinate 0 and the CD between subordinate 0 and subordinate 1. Continuing with the above example, a second 10m cable between subordinate 0 and subordinate 1 delays the SRF reception time at the main node by an additional five bits, thus causing the main node to adjust its response cycles to expect the SRF at the 529th (± 2) bit of the superframe.

The *SRF Response* figure illustrates how cable and transceiver delays affect the SRF response. In this case, the SRF miss error is not observed because the response cycles are adjusted during the discovery phase.

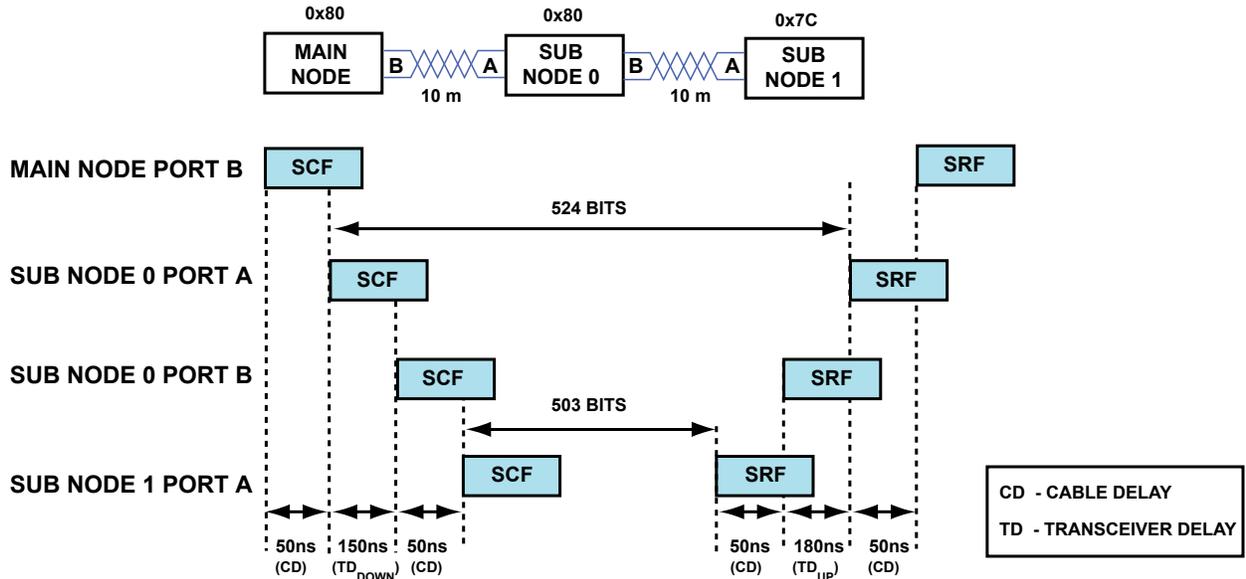


Figure 3-6: SRF Response

In this example:

- Subordinate 1 is the last-in-line subordinate node, which is responsible for initiating the SRF to commence the upstream portion of the superframe. When programmed with `A2B_RESPCYCS = 124 (0x7C)`, subordinate 1 is configured to generate the SRF at the 503rd bit of the superframe ($(4 * 124) + 7 = 503$).
- From the perspective of the upstream subordinate 0, the total delay between the SCF arriving to the subordinate 0 A-side transceiver during the downstream portion of the superframe and the corresponding SRF appearing there during the upstream portion of the same superframe is 430ns (21 bits), comprised of:
 - the downstream transceiver delay of subordinate 0 ($TD_{DOWN} = 150ns$),
 - the downstream cable delay between subordinate 0 and subordinate 1 ($CD = 5ns/m \times 10m = 50ns$),
 - the upstream cable delay between subordinate 1 and subordinate 0 ($CD = 5ns/m \times 10m = 50ns$), and
 - the upstream transceiver delay of subordinate 0 ($TD_{UP} = 180ns$)

Therefore, the number of bits between SCF arrival to the subordinate 0 A-side transceiver and the corresponding SRF being generated there is calculated to be $503 + 21 = 524$ bits for a 10m cable length between the subordinate 0 node and the subordinate 1 node.

- From the perspective of the main node, the total delay between generating the SCF and the corresponding SRF appearing during the upstream portion of the same superframe is 100ns (5 bits), which is comprised of:
 - the downstream cable delay between the main and subordinate 0 ($CD = 5ns/m \times 10m = 50ns$) and
 - the upstream cable delay between subordinate 0 and the main ($CD = 5ns/m \times 10m = 50ns$)

Therefore, the number of bits between SCF field generation and the corresponding SRF being received is calculated to be $524 + 5 = 529$ bits.

Managing A²B System Data Flow

Each main and subordinate transceiver in a full A²B system must be properly configured for the desired slot management scheme and format for both upstream and downstream traffic on the A²B bus between any two transceivers.

Each transceiver features two internal frame buffers:

- TX frame buffer - populated by the A²B bus and outputs to the DTX0 and/or DTX1 pins
- RX frame buffer - populated by the DRX0 and/or DRX1 input pins and outputs to the A²B bus

These frame buffers are populated and drained during each superframe, with the downstream slot content occupying the lower-order locations of the buffer and the upstream slot content in the higher-order locations. The frame buffers are 32 locations deep and 32-bits wide because any given transceiver can occupy up to 32 slots on the A²B bus and supports up to 32-bit data.

ATTENTION: If a transceiver is configured to receive more than 32 combined downstream and upstream slots from the A²B bus, the extra data associated with the upstream slots that cannot be accommodated by the frame buffer is dropped.

The TX frame buffer is populated by the A²B bus. The number of enabled downstream slots and specific slot masks determine which downstream slots are stored to the TX frame buffer during the downstream portion of the superframe. Similarly, the number of upstream data slots and specific slot masks determine which upstream slots are subsequently stored to the TX frame buffer after the downstream data. This combined buffer of data is then presented to the enabled DTXn data pins as a function of the number of transmit data pins enabled and whether or not interleaving is enabled.

The RX frame buffer is populated by the I²S/TDM port over the enabled DRXn data pins. The number of enabled receive pins and whether or not interleaving is turned on determine how data is placed into the RX frame buffer. Once the buffer is populated, the number of enabled downstream slots and specific slot masks determine which downstream slots are populated by the RX frame buffer during the downstream portion of the superframe. Similarly, the number of upstream data slots and specific slot masks determine which upstream slots are subsequently populated by the RX frame buffer after the downstream data has been sent.

Definition of dnmaskrx and upmaskrx

The dnmaskrx value is determined from the value of the [A2B_DNMASK0](#) through [A2B_DNMASK3](#) registers.

```
if (DNMASK3.RXDNSLOT31==1) dnmaskrx = 32;
else if (DNMASK3.RXDNSLOT30==1) dnmaskrx = 31;
else if (DNMASK3.RXDNSLOT29==1) dnmaskrx = 30;
. . .
else if (DNMASK0.RXDNSLOT02==1) dnmaskrx = 3;
```

```

else if (DNMASK0.RXDNSLOT01==1) dnmaskrx = 2;
else if (DNMASK0.RXDNSLOT00==1) dnmaskrx = 1;
else dnmaskrx = 0;
    
```

The upmaskrx value is determined from the value of the [A2B_UPMASK0](#) through [A2B_UPMASK3](#) registers

```

if (UPMASK3.RXUPSLOT31==1) upmaskrx = 32;
else if (UPMASK3.RXUPSLOT30==1) upmaskrx = 31;
else if (UPMASK3.RXUPSLOT29==1) upmaskrx = 30;
. . .
else if (UPMASK0.RXUPSLOT02==1) upmaskrx = 3;
else if (UPMASK0.RXUPSLOT01==1) upmaskrx = 2;
else if (UPMASK0.RXUPSLOT00==1) upmaskrx = 1;
else upmaskrx = 0;
    
```

A²B Slot Format

The normal (default) format of both upstream and downstream data slots is the data followed by a single parity bit. However, alternate formats supporting floating-point compression or ECC protection are also available. Both the size and the format of upstream and downstream data slots are configured using the [A2B_SLOTFMT](#) register. The *Slot Format* table summarizes the possible data formats configured by the [A2B_SLOTFMT.DNFMT](#), [A2B_SLOTFMT.DNSIZE](#), [A2B_SLOTFMT.UPFMT](#), and [A2B_SLOTFMT.UPSIZE](#) bits. In the *Slot Format* table, the FMT column is the [A2B_SLOTFMT.DNFMT](#) bit or the [A2B_SLOTFMT.UPFMT](#) bit, and the SIZE column is the 3-bit [A2B_SLOTFMT.DNSIZE](#) or [A2B_SLOTFMT.UPSIZE](#) field, depending on whether it is the downstream or upstream (respectively) slot format that is being configured.

Table 3-2: Slot Format

| FMT | SIZE | A ² B Slot Size | Compression | Protection | Data Width | A ² B Bus Bits |
|-----|-------|----------------------------|-------------|------------|------------|---------------------------|
| 0 | 0b000 | 8-bit | None | Parity | 8-bit | 9 |
| 0 | 0b001 | 12-bit | None | Parity | 12-bit | 13 |
| 0 | 0b010 | 16-bit | None | Parity | 16-bit | 17 |
| 0 | 0b011 | 20-bit | None | Parity | 20-bit | 21 |
| 0 | 0b100 | 24-bit | None | Parity | 24-bit | 25 |
| 0 | 0b101 | 28-bit | None | Parity | 28-bit | 29 |
| 0 | 0b110 | 32-bit | None | Parity | 32-bit | 33 |
| 0 | 0b111 | RESERVED | | | | |
| 1 | 0b000 | RESERVED | | | | |
| 1 | 0b001 | 12-bit | FP | Parity | 16-bit | 13 |
| 1 | 0b010 | 16-bit | FP | Parity | 20-bit | 17 |
| 1 | 0b011 | 20-bit | FP | Parity | 24-bit | 21 |
| 1 | 0b100 | 24-bit | None | ECC | 24-bit | 30 |

Table 3-2: Slot Format (Continued)

| FMT | SIZE | A ² B Slot Size | Compression | Protection | Data Width | A ² B Bus Bits |
|-----|-------|----------------------------|-------------|------------|------------|---------------------------|
| 1 | 0b101 | RESERVED | | | | |
| 1 | 0b110 | 32-bit | None | ECC | 32-bit | 39 |
| 1 | 0b111 | RESERVED | | | | |

NOTE: In the *Slot Format* table, the I²S/TDM Data Width column indicates the width of the actual data being exchanged over the I2S/TDM/PDM port in MSB-first format. Use cases for data widths in this column from 8 to 16 bits can optionally set the A2B_I2SGCFG.TDMSS bit to utilize 16-bit TDM channel data width over the I2S/TDM/PDM port. Data widths from 20-32 bits require the A2B_I2SGCFG.TDMSS bit to be cleared (32-bit TDM channel data width). See [I²S/TDM Port Programming Concepts](#) for more details.

ECC Protection

The transceiver provides support for both 24- and 32-bit data with ECC protection for the A²B bus data slots.

NOTE: As shown in the *Slot Format* table, there are six ECC bits for 24-bit data and seven ECC bits for 32-bit data.

ECC protection is useful in an environment where strong noise interferences (shorter than the superframe) are present, which otherwise can generate bit errors. ECC can be used in addition to the audio data error correction (repeat of last known good data), but it may only be used for non-audio data because it requires extra bus bandwidth.

Floating-Point Data Compression

The A²B protocol engine provides optional floating-point data compression/decompression so that less bandwidth is used on the A²B bus for a given data size (with better quality than the immediately lower data size). This compression can be used for A²B data sizes of 12, 16, and 20 bits, corresponding to the I²S data width. The compression encodes the number of leading sign bits in the source data as a 3-bit field and concatenates the sign bit itself, followed by $N-4$ bits of data (where N is the A²B data size). An example of 16-bit to 12-bit compression is shown in the *16-Bit to 12-Bit Compression Example* table. In the table, s is the sign bit and $\sim s$ is the inverse of the sign bit.

Table 3-3: 16-Bit to 12-Bit Compression Example

| 16-Bit Data | | | | | | | | | | | | | | | --> | 12-Bit FP Data | | | | | | | | | | | | | | |
|-------------|----|----|----|----|---|---|---|---|---|---|---|---|---|---|-----|----------------|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| s | ~s | x | x | x | x | x | x | x | x | x | y | y | y | y | y | --> | 0 | 0 | 0 | s | x | x | x | x | x | x | x | x | x | x |
| s | s | ~s | x | x | x | x | x | x | x | x | y | y | y | y | y | --> | 0 | 0 | 1 | s | x | x | x | x | x | x | x | x | x | x |
| s | s | s | ~s | x | x | x | x | x | x | x | x | y | y | y | y | --> | 0 | 1 | 0 | s | x | x | x | x | x | x | x | x | x | x |
| s | s | s | s | ~s | x | x | x | x | x | x | x | y | y | y | y | --> | 0 | 1 | 1 | s | x | x | x | x | x | x | x | x | x | x |

Table 3-3: 16-Bit to 12-Bit Compression Example (Continued)

| 16-Bit Data | | | | | | | | | | | | | | | --> | 12-Bit FP Data | | | | | | | | | | | | | |
|-------------|---|---|---|---|----|----|----|----|---|---|---|---|---|----------|----------|----------------|---|---|---|---|---|---|---|---|---|---|---|---|---|
| s | s | s | s | s | ~s | x | x | x | x | x | x | x | x | y | y | --> | 1 | 0 | 0 | s | x | x | x | x | x | x | x | x | x |
| s | s | s | s | s | s | ~s | x | x | x | x | x | x | x | x | y | --> | 1 | 0 | 1 | s | x | x | x | x | x | x | x | x | |
| s | s | s | s | s | s | s | ~s | x | x | x | x | x | x | x | x | --> | 1 | 1 | 0 | s | x | x | x | x | x | x | x | x | |
| s | s | s | s | s | s | s | s | ~s | x | x | x | x | x | x | x | --> | 1 | 1 | 1 | s | x | x | x | x | x | x | x | x | |

Data decompression reverses the process. The LSB of the compressed data (**L** in the *12-Bit to 16-Bit Data Decompression Example* table) is used to generate any remaining LSBs of the decompressed data that are not stored in the compressed format.

Table 3-4: Example of Data Decompression: 12 Bit to 16 Bit

| 12-Bit FP Data | | | | | | | | | | | | --> | 16-Bit Decompressed Data | | | | | | | | | | | | | | | | |
|----------------|---|---|---|---|---|---|---|---|---|---|----------|-----|--------------------------|----|----|----|----|----|----|---|---|---|----------|----------|----------|----------|----------|----------|----------|
| 0 | 0 | 0 | s | x | x | x | x | x | x | x | L | --> | s | ~s | x | x | x | x | x | x | x | x | L |
| 0 | 0 | 1 | s | x | x | x | x | x | x | x | L | --> | s | s | ~s | x | x | x | x | x | x | x | L | L | L | L | L | L | |
| 0 | 1 | 0 | s | x | x | x | x | x | x | x | L | --> | s | s | s | ~s | x | x | x | x | x | x | L | L | L | L | L | | |
| 0 | 1 | 1 | s | x | x | x | x | x | x | x | L | --> | s | s | s | s | ~s | x | x | x | x | x | L | L | L | L | | | |
| 0 | 0 | 0 | s | x | x | x | x | x | x | x | L | --> | s | s | s | s | s | ~s | x | x | x | x | L | L | L | | | | |
| 0 | 0 | 1 | s | x | x | x | x | x | x | x | L | --> | s | s | s | s | s | s | ~s | x | x | x | x | L | L | | | | |
| 0 | 1 | 0 | s | x | x | x | x | x | x | x | L | --> | s | s | s | s | s | s | ~s | x | x | x | x | L | | | | | |
| 0 | 1 | 1 | s | x | x | x | x | x | x | x | L | --> | s | s | s | s | s | s | s | x | x | x | x | L | | | | | |

Selecting FP compression is a good method to reduce the data slot size. It is beneficial in systems that requires multiple data channels. Sometimes it is beneficial to have enough or not enough data slots available. Reducing the slot size also reduces the current draw, which can be important in bus powered nodes.

The full dynamic range (24 bit = 144.49 dB) of the audio signal is preserved when data compression is enabled. The human ear can listen to sounds near the noise level in a quiet environment, but the human ear masks very quiet audio content in the presence of very loud audio content. The floating-point compression (to 20 bit) takes advantage of this psychoacoustic effect and removes low-level content in the presence of high-level audio content. The floating-point compression preserves all low-level content (here, 16-bits = 96.33 dB for 20-bit data slots) when there is no high-level audio content and supports the full dynamic range for strong audio signals (up to 144.49 dB for 20 bit data slots), always with 16 bit = 96.33 dB resolution.

Downstream Data Slots

Subordinate nodes can selectively receive downstream bus slots for output onto the DTXn pins. A programmable number of I²S/TDM data channels on the DRXn pins (`A2B_DNOFFSET`) can be skipped before the next in line channels are presented as downstream data slots to the A²B bus. As a result, this mode allows subordinate nodes to receive and transmit downstream data.

The `A2B_DNMASK0` through `A2B_DNMASK3` registers provide one bit for each possible downstream data slot. These downstream mask bits select which downstream slots are consumed by the transceiver and placed in its TX frame buffer for output over the I²S/TDM port, as governed by the downstream mask enable (`A2B_LDNSLOTS.DNMaskEN`) bit.

When `A2B_LDNSLOTS.DNMaskEN= 0`, the `A2B_DNSLOTS` register defines the number of downstream data slots, starting immediately after the SCF, which are passed downstream through the subordinate node, and the `A2B_LDNSLOTS` register defines the number of downstream data slots which are captured by the transceiver during the downstream portion of the superframe. The transceiver consumes and does not pass these data slots downstream to the next node. As such, a subordinate transceiver receives "`A2B_BCDNSLOTS + A2B_DNSLOTS + A2B_LDNSLOTS`" downstream data slots on the A-side transceiver and transmits "`A2B_BCDNSLOTS + A2B_DNSLOTS`" downstream data slots on the B-side transceiver.

NOTE: When the `A2B_LDNSLOTS.DNMaskEN` bit is cleared in a subordinate transceiver, the `A2B_DNMASK0` through `A2B_DNMASK3` registers are ignored.

When `A2B_LDNSLOTS.DNMaskEN= 1`, the `A2B_LDNSLOTS` register defines the number of data slots which the local node adds during the downstream portion of the superframe. These data slots are passed downstream through the local node after the `A2B_DNSLOTS` data slots. The most significant bit that is set in the `A2B_DNMASK0` through `A2B_DNMASK3` registers determines the number of slots that must be received by the transceiver (`dnmaskrx`) for it to then identify which individual slots are placed in its RX frame buffer for output over the I²S/TDM port. To that, a subordinate node receives $\text{MAX}(\text{A2B_DNSLOTS}, \text{dnmaskrx})$ downstream data slots on the A-side transceiver and transmits "`A2B_DNSLOTS + A2B_LDNSLOTS`" downstream data slots on the B-side transceiver.

NOTE: When the `A2B_LDNSLOTS.DNMaskEN` bit is set in a subordinate transceiver, the `A2B_BCDNSLOTS` register is ignored.

The value of the `A2B_DNOFFSET` register is meaningful only when the subordinate transceiver is configured to transmit downstream data (`A2B_LDNSLOTS.DNMaskEN= 1`, and the `A2B_LDNSLOTS` register is non-zero). Data is placed in the enabled downstream slots starting with the beginning of the RX frame buffer unless the `A2B_DNOFFSET` register has been programmed to apply an offset into the RX frame buffer from which it begins populating the enabled downstream slots.

The *Subordinate Node Using the A2B_DNMASKn and A2B_DNOFFSET Registers* figure is an example of how downstream data slots are used in a subordinate transceiver after programming the `A2B_DNMASK0`, `A2B_DNMASK1`, and `A2B_DNOFFSET` registers.

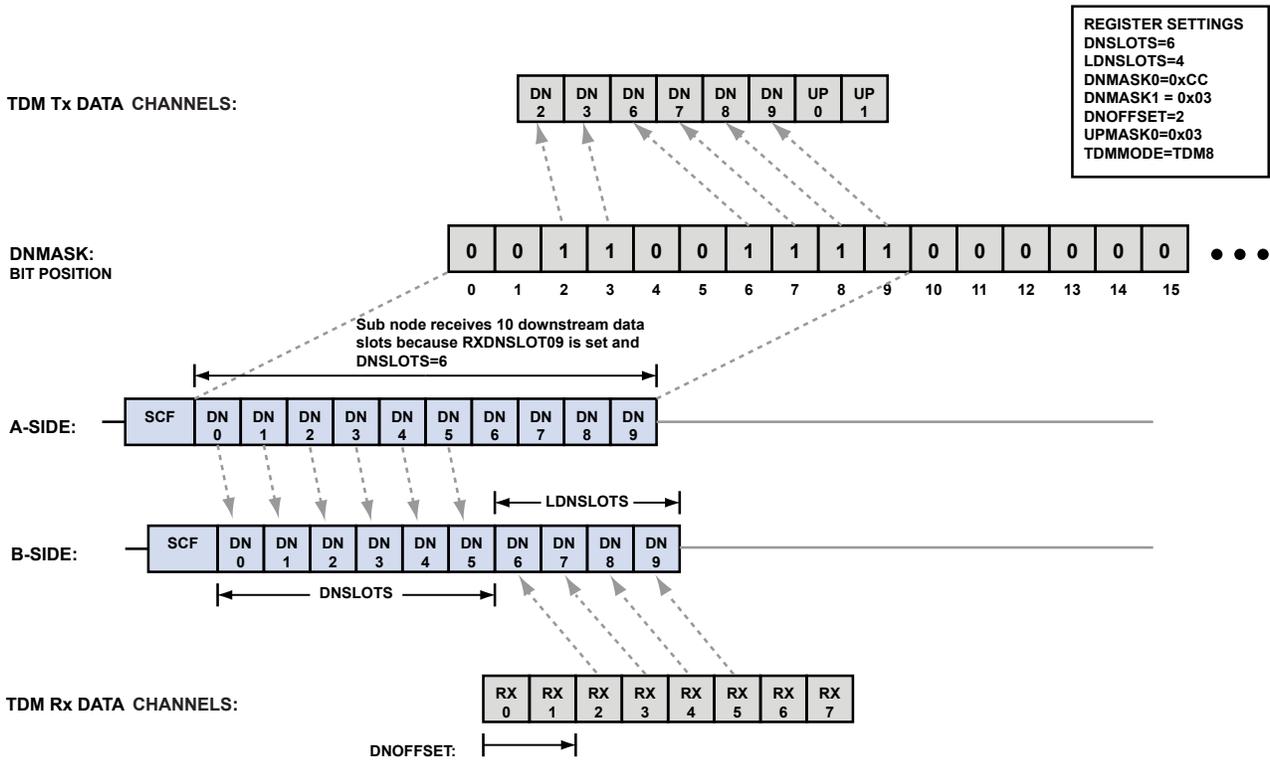


Figure 3-7: Subordinate Node Using the A2B_DNMAKn and A2B_DNOFFSET Registers

Upstream Data Slots

The [A2B_UPSLOTS](#) register defines the number of upstream data slots. For a main transceiver, this register defines the number of data slots that will come upstream to the main from the first-in-line subordinate transceiver. For a subordinate transceiver, this register defines the number of upstream data slots, starting immediately after the SRF with slot 0, that are passed upstream through the transceiver, whether or not that subordinate transceiver uses the information contained in those slots.

The [A2B_LUPSLOTS](#) register defines the number of data slots that the subordinate transceiver appends to the upstream portion of the superframe after the data slots being passed upstream by the subordinate, as defined in the [A2B_UPSLOTS](#) register. The data placed in the upstream data slots comes from the transceiver's internal RX frame buffer, as populated by its I²S/TDM/PDM port.

A subordinate transceiver selectively receives upstream A²B bus data slots into its TX frame buffer for output onto its DTXn pin(s) for use in the subordinate node. In subordinate transceivers, the [A2B_UPMASK0](#) through [A2B_UPMASK3](#) registers provide one bit for each possible upstream data slot. When a bit is set in any of these registers, the subordinate transceiver takes the upstream data from the corresponding slot and places it in its TX frame buffer after any received downstream data slots, which will then be output to the appropriate DTXn pin(s) via the I²S/TDM port.

The most significant bit set in the [A2B_UPMASK0](#) through [A2B_UPMASK3](#) registers defines the number of slots (upmaskrx) that the transceiver must receive in order to then appropriately place enabled slots into the TX frame

buffer for output to the I²S/TDM port. To that, a subordinate transceiver receives MAX (A2B_UPSLOTS, up-maskrx) upstream data slots on the B-side transceiver. It then transmits "A2B_UPSLOTS + A2B_LUPSLOTS" upstream data slots on the A-side transceiver.

A programmable number of I²S/TDM data channels on the DRXn pins (A2B_UPOFFSET) can be skipped before the next-in-line channels are presented as upstream data slots to the A²B bus. By default, a subordinate node populates the enabled upstream slots with the first entry in its RX frame buffer. The A2B_UPOFFSET register can be written to define an offset into the RX frame buffer from which it begins populating the enabled upstream slots.

The *Subordinate Node Using the A2B_UPMASKn and A2B_UPOFFSET Registers* figure provides an example of how upstream data slots are used in a subordinate transceiver after programming the A2B_UPMASK0, A2B_UPMASK1, and A2B_UPOFFSET registers.

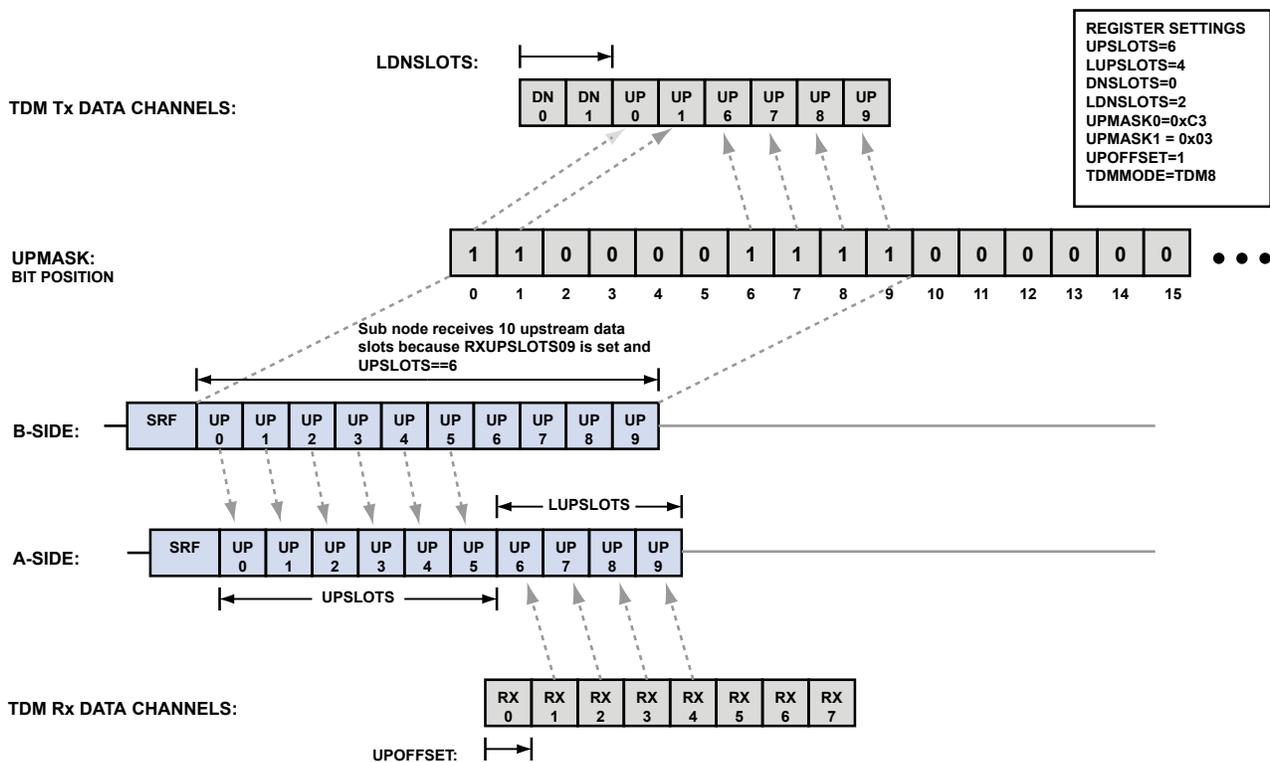


Figure 3-8: Subordinate Node Using the A2B_UPMASKn and A2B_UPOFFSET Registers

A²B Bandwidth

All upstream data is restricted to the same slot size (bits per slot), and all downstream data is restricted to the same slot size (bits per slot), but the downstream slot size can be different from the upstream slot size. A detailed calculation spreadsheet and the SigmaStudio software are available from Analog Devices to calculate bandwidth for all possible cases.

The *Bandwidth Examples* table provides bandwidth examples for 48 kHz sampled, synchronous upstream and downstream data slots. To simplify the table, every node uses the same number of upstream and downstream slots. Use up to 32 slots for upstream data and up to 32 slots for downstream data.

Table 3-5: Bandwidth Examples

| Sub Nodes | Downstream Slots Per Node (Speaker) | Upstream Slots Per Node (Mics) | Slot Size (Bits Per Slot) | Sum of Downstream Slots (Max. 32) | Sum of Upstream Slots (Max. 32) | Sum of Downstream and Upstream Slots |
|-----------|-------------------------------------|--------------------------------|---------------------------|-----------------------------------|---------------------------------|--------------------------------------|
| 9 | 2 | 2 | 16 | 18 | 18 | 36 |
| 8 | 4 | 1 | | 32 | 8 | 40 |
| 7 | 4 | 2 | | 28 | 14 | 42 |
| 6 | 5 | 2 | | 30 | 12 | 42 |
| 5 | 6 | 3 | | 30 | 15 | 45 |
| 4 | 8 | 3 | | 32 | 12 | 44 |
| 3 | 10 | 6 | | 30 | 18 | 48 |
| 2 | 16 | 9 | | 32 | 18 | 50 |
| 1 | 32 | 19 | | 32 | 19 | 51 |
| 9 | 3 | 0 | 24 | 27 | 0 | 27 |
| 8 | 3 | 0 | | 24 | 0 | 24 |
| 7 | 4 | 0 | | 28 | 0 | 28 |
| 6 | 5 | 0 | | 30 | 0 | 30 |
| 5 | 6 | 0 | | 30 | 0 | 30 |
| 4 | 8 | 0 | | 32 | 0 | 32 |
| 3 | 10 | 1 | | 30 | 3 | 33 |
| 2 | 16 | 1 | | 32 | 2 | 34 |
| 1 | 32 | 2 | | 32 | 2 | 34 |

I²S/TDM Port Programming Concepts

Programming the I²S/TDM interface involves selecting the mode of operation for the port, controlling how many data pins are enabled for both transmit and receive operations, and configuring the polarity and timing of the BCLK and SYNC signals relative to data.

The `A2B_I2SGCFG` and `A2B_I2SCFG` registers are used to configure the I²S/TDM port to support these various modes of operation. The *Serial Mode Data and Clock Formats* table provides a summary of the different data and clock formats supported by both main and subordinate transceivers.

Table 3-6: Serial Mode Data and Clock Formats

| Bit Setting | Data and Clock Format |
|------------------------------------|---|
| <code>A2B_I2SGCFG.EARLY = 0</code> | SYNC pin changes in the same cycle as the MSB of Data Channel 0 |

Table 3-6: Serial Mode Data and Clock Formats (Continued)

| Bit Setting | Data and Clock Format |
|-------------------------|--|
| A2B_I2SGCFG.EARLY =1 | SYNC pin changes one cycle before the MSB of Data Channel 0 |
| A2B_I2SGCFG.ALT =0 | SYNC pin is driven high for one BCLK cycle at the start of each sampling period |
| A2B_I2SGCFG.ALT =1 | SYNC pin is driven high at the beginning of each sampling period and low in the middle of each sampling period |
| A2B_I2SGCFG.INV =0 | Rising edge of SYNC references the first channel (Channel 0) |
| A2B_I2SGCFG.INV =1 | Falling edge of SYNC references the first channel (Channel 0) |
| A2B_I2SCFG.RXBCLKINV =0 | DRX0, DRX1, and SYNC pins are sampled on the rising edge of BCLK |
| A2B_I2SCFG.TXBCLKINV =0 | DTX0, DTX1, and SYNC pins change on the rising edge of BCLK |
| A2B_I2SCFG.RXBCLKINV =1 | DRX0, DRX1, and SYNC pins are sampled on the falling edge of BCLK |
| A2B_I2SCFG.TXBCLKINV =1 | DTX0, DTX1, and SYNC pins change on the falling edge of BCLK |

To support more than a stereo two-channel (TDM2) signal, the A2B_I2SGCFG.TDMMODE field must be set to enable any of the supported TDM modes of operation. Once configured, this is the operating mode used for each of the enabled data pins, as controlled by the A2B_I2SCFG.RX0EN, A2B_I2SCFG.RX1EN, A2B_I2SCFG.TX0EN, and A2B_I2SCFG.TX1EN bits.

When both data pins in either direction are enabled, the interleaving feature can be enabled by setting the respective two-pin interleave (A2B_I2SCFG.RX2PINTL and A2B_I2SCFG.TX2PINTL) bit. When set, the even slot data is associated with the DTX0/DRX0 data pin, and the odd slot data is associated with the DTX1/DRX1 data pin. When cleared, the lower half of the enabled slots are associated with the DTX0/DRX0 data pin, and the upper half of the enabled slots are associated with the DTX1/DRX1 data pin. For example, if the data format is set for I²S or TDM2 mode, the *Data Channel Structure for TDM2 Setting* figure summarizes how the data is aligned.

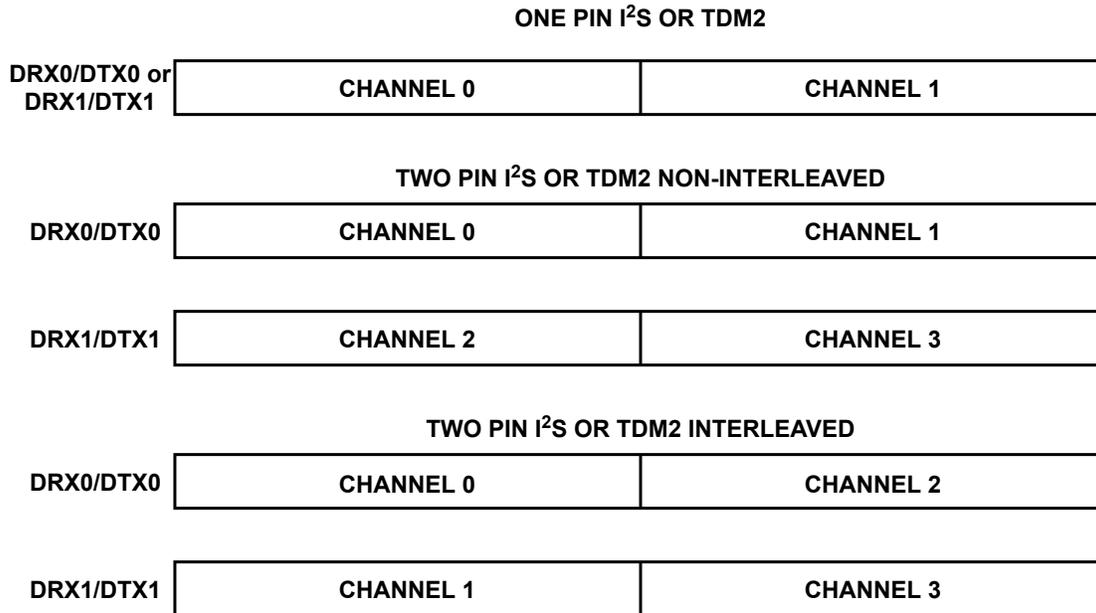


Figure 3-9: Data Channel Structure for TDM2 Setting (TDMMODE == 000)

NOTE: Single-pin transmit can be on either the DTX0 or DTX1 pin, and single-pin receive can be on either the DRX0 or DRX1 pin.

The `A2B_I2SGCFG.TDMSS` bit selects between 16-bit and 32-bit serial data for the I²S/TDM port, and it is the responsibility of the host to ensure that the appropriate timing signals are provided to accommodate the full window of data. For example, if TDM8 mode is selected (`A2B_I2SGCFG.TDMMODE = 0b010`), then the host must provide either 128 (8 x 16-bit, when `A2B_I2SGCFG.TDMSS = 1`) or 256 (8 x 32-bit, when `A2B_I2SGCFG.TDMSS = 0`) BCLK pulses for the data and the appropriate SYNC signal (to be either pulsed or held for a 50% duty cycle, per the setting of the `A2B_I2SGCFG.ALT` bit), as shown in the *I²S/TDM8 Example Timing* figure.

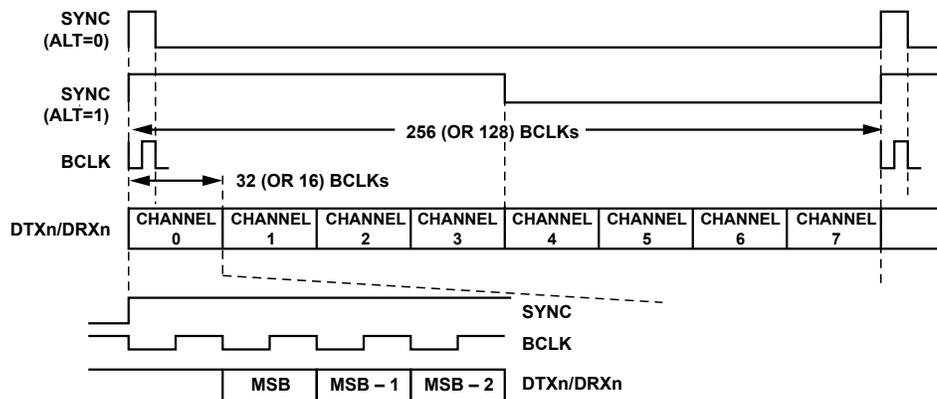


Figure 3-10: I²S/TDM8 Example Timing

As shown in the *I²S/TDM8 Example Timing* figure, the TDM channel data is in MSB-first format. When the data being exchanged over the A²B bus is not exactly 16-bit or 32-bit (as configured by the `A2B_I2SCFG.TDMSS` bit), the transceiver expects the input TDM data channels to arrive in MSB-first format and disregards any lower-order bits. When outputting to the local node, the transceiver presents the received A²B slot data to the I²S/TDM port in MSB-first format with the unused lower-order bits zero-filled. For example, if the A²B slot is configured for 12-bit data (`A2B_SLOTFMT.UPSIZE = 1` for upstream slots or `A2B_SLOTFMT.DNSIZE = 1` for downstream slots), the 12-bit input data must be left-justified in the TDM channel, and output data consists of the 12-bit A²B slot data followed by four zero bits.

If SYNC arrives one bit earlier, it can be rephrased to data arriving one bit later than the relevant edge on the SYNC signal. The *I²S/TDM2 to TDM16 A²B Main or Subordinate* figure shows the typical timing for I²S, as well as the TDM2 to TDM16 interface modes with programmable options. Data is provided on one edge of BCLK and sampled on the opposite edge of BCLK (`A2B_I2SCFG.TXBCLKINV ≠ A2B_I2SCFG.RXBCLKINV`).

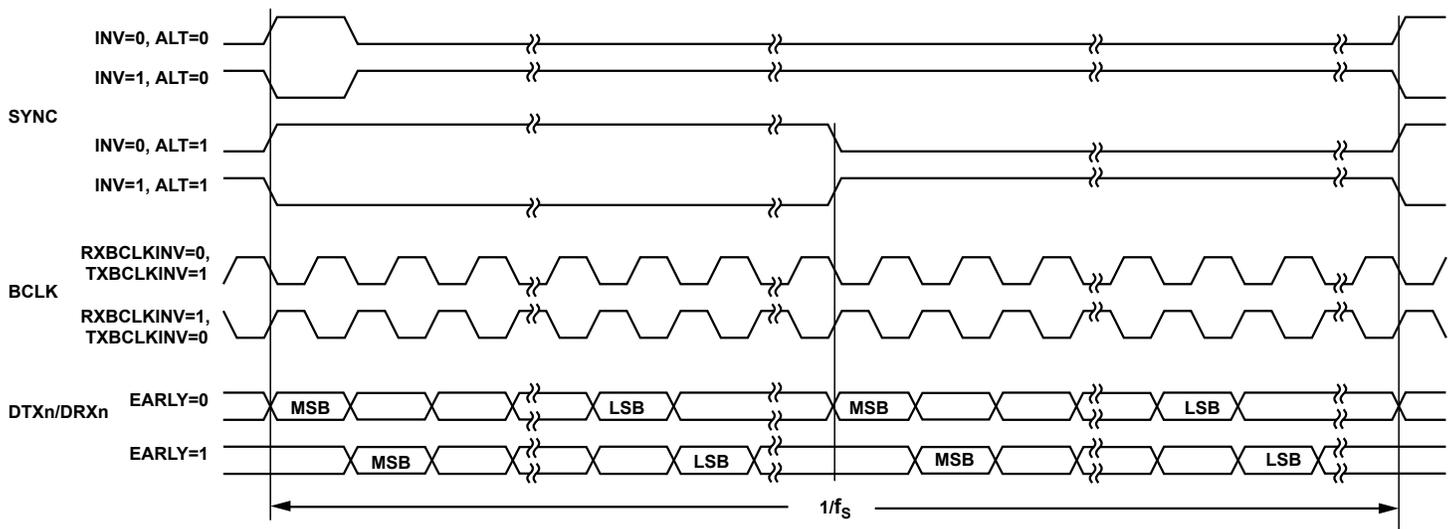


Figure 3-11: I²S/TDM2 to TDM16 A²B Main or Subordinate

The full 32-channel combined bandwidth is available when both data pins are enabled in TDM16 mode.

CAUTION: Be cautious if only one data pin is available for a TDM32 interface, as this increases the BCLK rate to a speed at which race conditions can occur.

The A²B main samples data on a BCLK edge and changes data on the previous, same polarity BCLK edge (`A2B_I2SCFG.TXBCLKINV = A2B_I2SCFG.RXBCLKINV`), as shown in the *TDM32 A²B Main* figure.

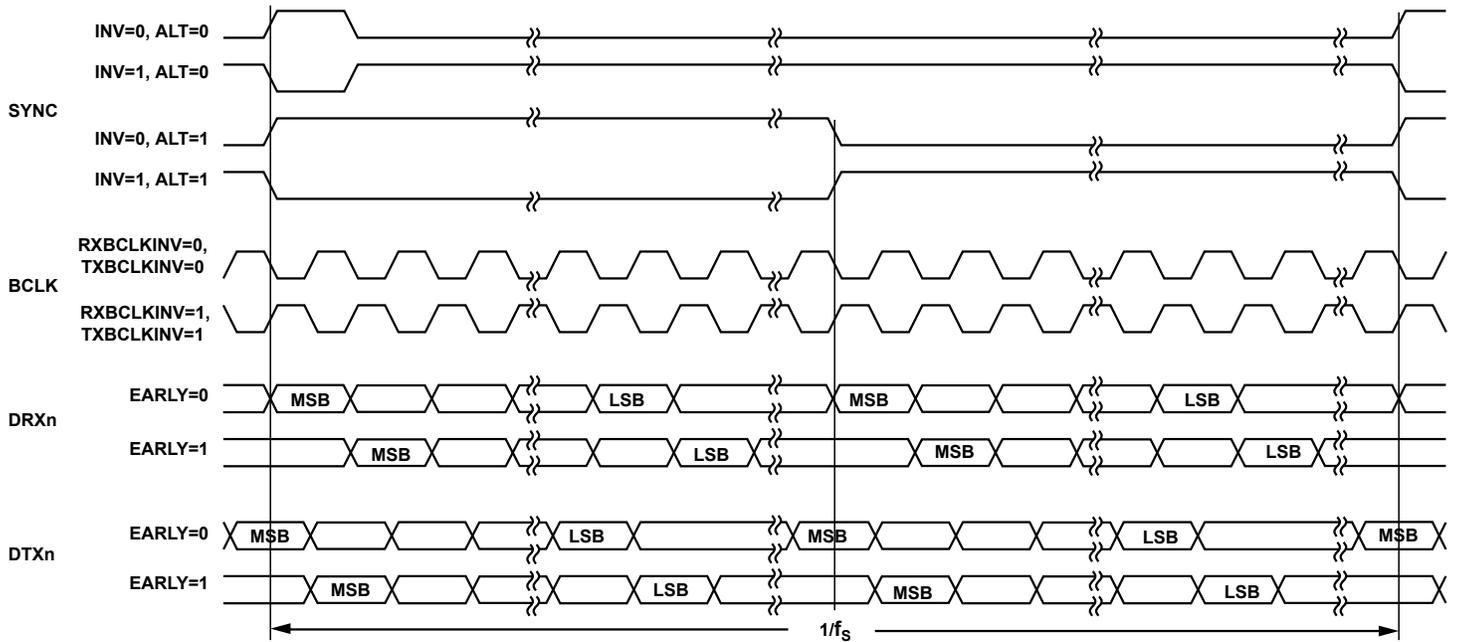


Figure 3-12: TDM32 A²B Main

The A²B subordinate changes data on a BCLK edge and samples data on the next, same polarity BCLK edge ($A2B_I2SCFG.TXBCLKINV = A2B_I2SCFG.RXBCLKINV$), as shown in the *TDM32 A²B Subordinate* figure.

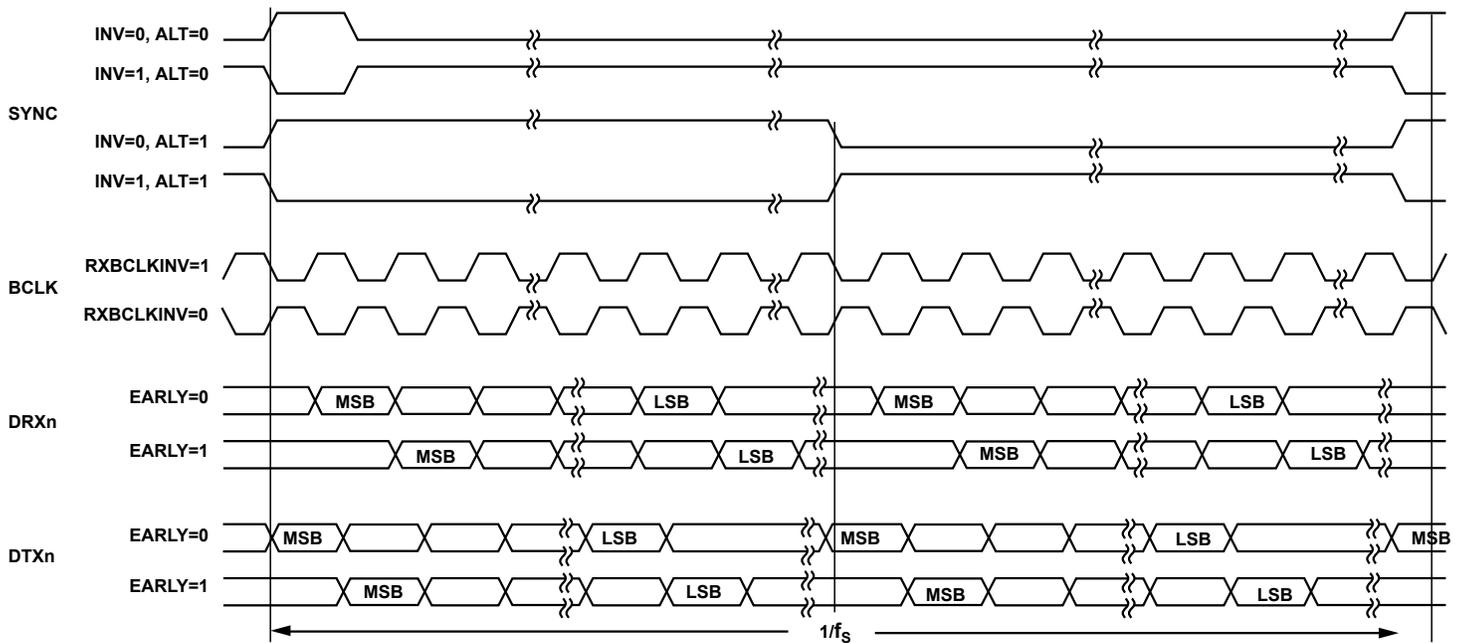


Figure 3-13: TDM32 A²B Subordinate

Synchronizing Subordinate Nodes

A²B subordinate nodes can all sample at exactly the same time by individually compensating for their propagation delay in the `A2B_SYNCOFFSET` register. Writing a non-zero value to this register adjusts the A²B bus clock (f_{SYSBCLK}) cycle on which the SYNC pin indicates the start of an audio frame for that particular subordinate transceiver. As the programmed value is the 8-bit signed two's complement representation of the integer number of SYSBCLK cycles between where the SYNC occurs and where the superframe subsequently begins, only negative values are valid.

The maximum value that can be programmed into the `A2B_SYNCOFFSET` register defines a SYNC signal to occur 104 SYSBCLK cycles before the start of the superframe ($-104 = 0x98$), but this value is only valid for the subordinate node that is the furthest away from the main node in a fully populated A²B network topology (`A2B_NODEADR.NODE = 0x09`). For any subordinate node n that is nearer to the main node, the valid ranges supporting a predictable transfer of I²S/TDM data to A²B slots are a function of the location of subordinate node n in the network, as governed by the formula:

$$(-32 - 8n) \leq \text{A2B_SYNCOFFSET} \leq 0$$

The *Supported SYNC Offset* table summarizes the valid settings for the `A2B_SYNCOFFSET` register for any given subordinate node in SYSBCLK cycles (Offset Range).

Table 3-7: Supported SYNC Offset

| Subordinate Node n | Offset Range | <code>A2B_SYNCOFFSET</code> Range |
|--------------------|--------------|-----------------------------------|
| 0 | -32 to 0 | 0xE0 to 0x00 |
| 1 | -40 to 0 | 0xD8 to 0x00 |
| 2 | -48 to 0 | 0xD0 to 0x00 |
| 3 | -56 to 0 | 0xC8 to 0x00 |
| 4 | -64 to 0 | 0xC0 to 0x00 |
| 5 | -72 to 0 | 0xB8 to 0x00 |
| 6 | -80 to 0 | 0xB0 to 0x00 |
| 7 | -88 to 0 | 0xA8 to 0x00 |
| 8 | -96 to 0 | 0xA0 to 0x00 |
| 9 | -104 to 0 | 0x98 to 0x00 |

I²S Reduced Data Rate

Subordinate nodes can also run the I²S/TDM interface at a reduced rate frequency with respect to the superframe rate (f_{SYNCM}). The reduced-rate frequency is derived by dividing the superframe rate by a programmable set of values. Different subordinate nodes can be configured to run at different reduced I²S/TDM rates.

The `A2B_I2SRATE.I2SRATE` bit field is used to divide the superframe A²B rate down to the reduced I²S rate. It also provides a control bit, `A2B_I2SRRATE.RBUS`, to enable reduced-rate data slots on the bus. The A²B data slots on the bus are transmitted only once every `A2B_I2SRRATE.RRDIV` superframes.

The `A2B_I2SRATE.I2SRATE` bit field can be used to program the division factor to 2, 4, or as configured in the `A2B_I2SRRATE.RRDIV` field. The `A2B_I2SRATE.SHARE` bit enables the shared A²B bus slots in a reduced-rate subordinate node, provided the node has the I²S transmit disabled.

The `A2B_I2SRRCTL` register provides bits to allow a processor to track the full-rate audio frame, which contains new reduced-rate samples. The pin can be used as a strobe by setting the `A2B_I2SRRCTL.ENSTRB` bit, which indicates the audio frame where reduced-rate data is updated. The `A2B_I2SRRCTL.STRBDIR` bit configures the direction of the pin when used as a strobe. The reduced rate strobe output at the main node is based on the `A2B_I2SRRATE.RRDIV` field setting. When the `A2B_I2SRRATE.RRDIV` field is not one, the reduced rate count is maintained in each node, and the strobe output signal is generated accordingly. When the strobe is an input, it is sampled on the active edge of SYNC, and the reduced rate count is synchronized to it. The user must create a strobe signal that matches the `A2B_I2SRRATE.RRDIV` setting.

The `A2B_I2SRRSOFFS` register provides a bit field to move the SYNC edge in a reduced-rate subordinate node in superframe increments.

The *Reduced Data Rate* figure shows how the upstream slots from the transceiver can reduce the superframe rate on the bus, allowing the subordinate nodes to run at a reduced-sample frequency with both sharing disabled (`A2B_I2SRATE.SHARE = 0`) and enabled (`A2B_I2SRATE.SHARE = 1`). This figure is drawn for a system with one main node and one subordinate node.

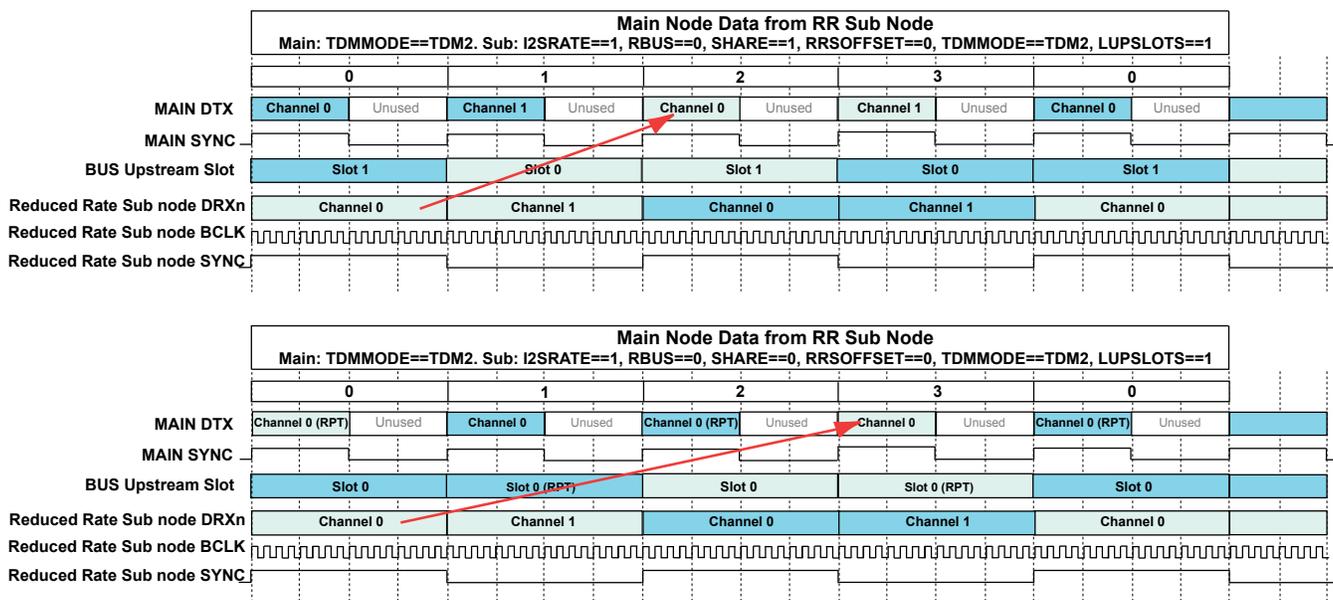


Figure 3-14: Reduced Data Rate

The following table shows the I²S/TDM sampling rates categorized into system modes for the reduced rate.

Table 3-8: I²S/TDM Sampling Rates Categorized into System Modes for Reduced Rate

| Mode | Host I ² S/TDM Rate | Bus Data Slots | Subordinate Rate(s) | Channels |
|------|--------------------------------|----------------------------|--|----------|
| 1 | Set in A2B_I2SRRATE .RRDIV | Set in A2B_I2SRRATE .RRDIV | Set in A2B_I2SRRATE .RRDIV | 1 - 32 |
| 2 | 48 kHz | Set in A2B_I2SRRATE .RRDIV | Set in A2B_I2SRRATE .RRDIV | 1 - 32 |
| 3 | 48 kHz | 48 kHz | Set in A2B_I2SRRATE .RRDIV | 1 - 32 |
| 4 | 48 kHz | 48 kHz | Set in A2B_I2SRRATE .RRDIV | 1 - 128 |
| 5 | 48 kHz | 48 kHz | Set in A2B_I2SRRATE .RRDIV, 1/4x, 1/2x, 1x, 2x, 4x | 1 - 32 |

The reduced rate feature allows system designers to add the following functionality:

1. Subordinate nodes can run the I²S/TDM interface at a reduced rate divided from the superframe rate, as divided down from the superframe rate. For example, reduced rates for a 48 kHz superframe rate are 24 kHz, 12 kHz, 6 kHz, 4 kHz, 3 kHz, 2.4 kHz, 2 kHz, 1.71 kHz, or 1.5 kHz. The I²S/TDM RX data on the subordinate node can be sent either upstream or downstream at the reduced rate.

Different subordinate nodes can run at different reduced I²S/TDM rate.

2. The SYNC signal of the reduced-rate subordinate node can be adjusted in superframe increments to ensure minimum latency on the delivery of reduced-rate data.
3. Control of the BCLK signal generation can minimize a delay by quick sampling at the reduced-rate I²S data (for example, within a 48 kHz I²S/TDM frame) or sampling at the reduced I²S/TDM rate.
4. Options to notify a processor when the reduced-rate I²S/TDM data channels are updated.
5. Option to run the bus data slots at the full, continuous audio rate (nominally 48 kHz) or a reduced rate. The rate can be reduced by:
 - a. Skipping data slots for superframes that do not contain data (for example, only reduced sampling rate microphone nodes on the A²B bus). This approach saves power by reducing the bus activity level but does not increase channel bandwidth on the bus. When the same A²B data slots are shared between multiple I²S/TDM channels in a node, the program cannot skip the A²B data slots.
 - b. Time-dividing bus data slots of a node into multiple I²S/TDM channels and not skipping data slots for superframes. This approach is used if different types of subordinate nodes connecting on the same A²B bus (for example, a multi-axis accelerometer node with a microphone or amp nodes on the same bus). The bus must run at the full-data rate to allow for A²B data slot sharing. This approach provides for increased channel bandwidth on the bus by allowing reduced-rate subordinate nodes to time-multiplex I²S/TDM data words over bus data slots.

- Subordinate nodes running at ½ rate can use 2:1 time multiplexing (two I²S/TDM channels in the same subordinate node alternate on one A²B slot).
- Subordinate nodes running at lower rates can use 4:1 time multiplexing (four I²S/TDM channels in the same subordinate node alternate on one A²B slot).
- Time multiplexing of A²B data slots beyond 4:1 is not supported.
- Time multiplexing of A²B data slots between nodes is not supported.
- The bus must be run with A²B data slots at the full, continuous audio rate for data slots to be shared.
- The I²S/TDM RX reduced rate data can be transmitted upstream or downstream.

I²S Reduced Rate Restrictions

Observe the following general restrictions when using the I²S reduced rate feature.

- Each subordinate node can only run at a single I²S/TDM rate.
- Configure subordinate nodes running at a reduced I²S/TDM rate for the I²S/TDM RX data, not the I²S/TDM TX data. This means that the reduced-rate subordinate nodes must have `A2B_I2SCFG.TX0EN = 0` and `A2B_I2SCFG.TX1EN = 0`.
- If `A2B_I2SRRATE.RBUS` is set and a reduced rate is configured (`A2B_I2SRRATE.RRDIV > 1`), subordinate nodes must have an `A2B_I2SRATE.I2SRATE` value of 0 (SFF x 1) or 3 (SFF / `A2B_I2SRRATE.RRDIV`).

Restrictions on Data Slot Sharing (`A2B_I2SRATE.SHARE = 1`)

Observe the following data slot sharing restrictions when using the I²S reduced rate feature.

- The bus must run at the full-data rate (`A2B_I2SRRATE.RBUS = 0`) to allow for A²B data slot sharing. A²B data slot skipping cannot be used when the same A²B data slots are shared between multiple I²S/TDM channels in a node.
- Data slots on the A²B bus produced by a reduced-rate subordinate node with `A2B_I2SRATE.SHARE = 1` must be received from the A²B bus by full- or increased-rate nodes.
- If the `A2B_I2SRATE.SHARE` bit is set in a reduced-rate subordinate node, the maximum synchronization offset is one superframe (`A2B_I2SRRSOFFS.RRSOFFSET` must be 0 or 1).

If the `A2B_I2SRATE.SHARE` bit is set in a reduced-rate subordinate node and there is no synchronization offset (`A2B_I2SRRSOFFS.RRSOFFSET = 0`), there is a further constraint on the node programming relative to *N* (the number of usable up and down slots). For example, if TDMS is the number of slots per frame on one pin of a reduced-rate subordinate node (which is 2, 4, 8, 16, or 32), *N* is calculated as shown in the following table:

| I ² S/TDM Divide Ratio | Number of Slots (N) |
|-----------------------------------|---|
| 2 | TDMS >> 1 |
| 4 | (TDMS >> 1) + (TDMS >> 2) |
| > 4 | (TDMS >> 1) + (TDMS >> 2) + (TDMS >> 3) |

If the reduced-rate subordinate node has the `A2B_I2SCFG.RX0EN`, `A2B_I2SCFG.RX1EN`, and `A2B_I2SCFG.RX2PINTL` bits all set, "`A2B_LUPSLOTS + A2B_UPOFFSET`" must be $\leq 2N$. Otherwise, "`A2B_LUPSLOTS + A2B_UPOFFSET`" must be $\leq N$.

If the reduced-rate subordinate node is generating downstream data slots (`A2B_LDNSLOTS.DNMASKEN = 1`), the same constraint applies to "`A2B_LDNSLOTS + A2B_DNOFFSET`".

Restrictions on Alternate BCLK Rate (`A2B_I2SRATE.BCLKRATE`)

Observe the following alternate BCLK rate restrictions when using the I²S reduced rate feature.

- In a reduced-rate subordinate node, if the I²S rate setting is $SFF / 2$ (`A2B_I2SRATE.I2SRATE = 1`), do not set the BCLK frequency to $SYNC \times 4096$ (`A2B_I2SRATE.BCLKRATE != 2`).
- If the system-level reduced rate divisor is 1 (`A2B_I2SRRATE.RRDIV = 1`) and the I²S rate setting is " $SFF / A2B_I2SRRATE.RRDIV$ " (`A2B_I2SRATE.I2SRATE = 3`), do not set the BCLK frequency to " $SYNC \times 2048$ " (`A2B_I2SRATE.BCLKRATE = 1`) or " $SYNC \times 4096$ " (`A2B_I2SRATE.BCLKRATE = 2`).
- If the system-level reduced rate divisor is 2 (`A2B_I2SRRATE.RRDIV = 2`) and the I²S rate setting is " $SFF / A2B_I2SRRATE.RRDIV$ " (`A2B_I2SRATE.I2SRATE = 3`), do not set the BCLK frequency to " $SYNC \times 4096$ " (`A2B_I2SRATE.BCLKRATE = 2`).
- If the BCLK frequency is not determined by the value programmed in the `A2B_I2SGCFG` register (`A2B_I2SRATE.BCLKRATE != 0`) in a reduced rate subordinate node, the synchronization offset cannot exceed 1 superframe (`A2B_I2SRRSOFFS.RRSOFFSET < 2`).

I²S Increased Data Rate

The A²B subordinate transceiver supports increased sampling rates at the I²S/TDM interface with respect to the superframe rate (f_{SYNCR}). The local sampling rate of the subordinate node can be programmed to $1 f_{SYNCR}$, $2 f_{SYNCR}$, or $4 f_{SYNCR}$ in the `A2B_I2SRATE` register. For example, given a 48 kHz superframe frequency, the local sampling rate can be set to 48 kHz, 96 kHz, or 192 kHz, respectively. The *Increased Data Rate* figure shows how the downstream and upstream slots from the A²B superframe are distributed on the DTX0/DTX1 and DRX0/DRX1 pins in the subordinate transceiver for different `A2B_I2SRATE` bit settings (with `A2B_I2SRATE.REDUCE = 0`) in a system with one controller and one subordinate node.

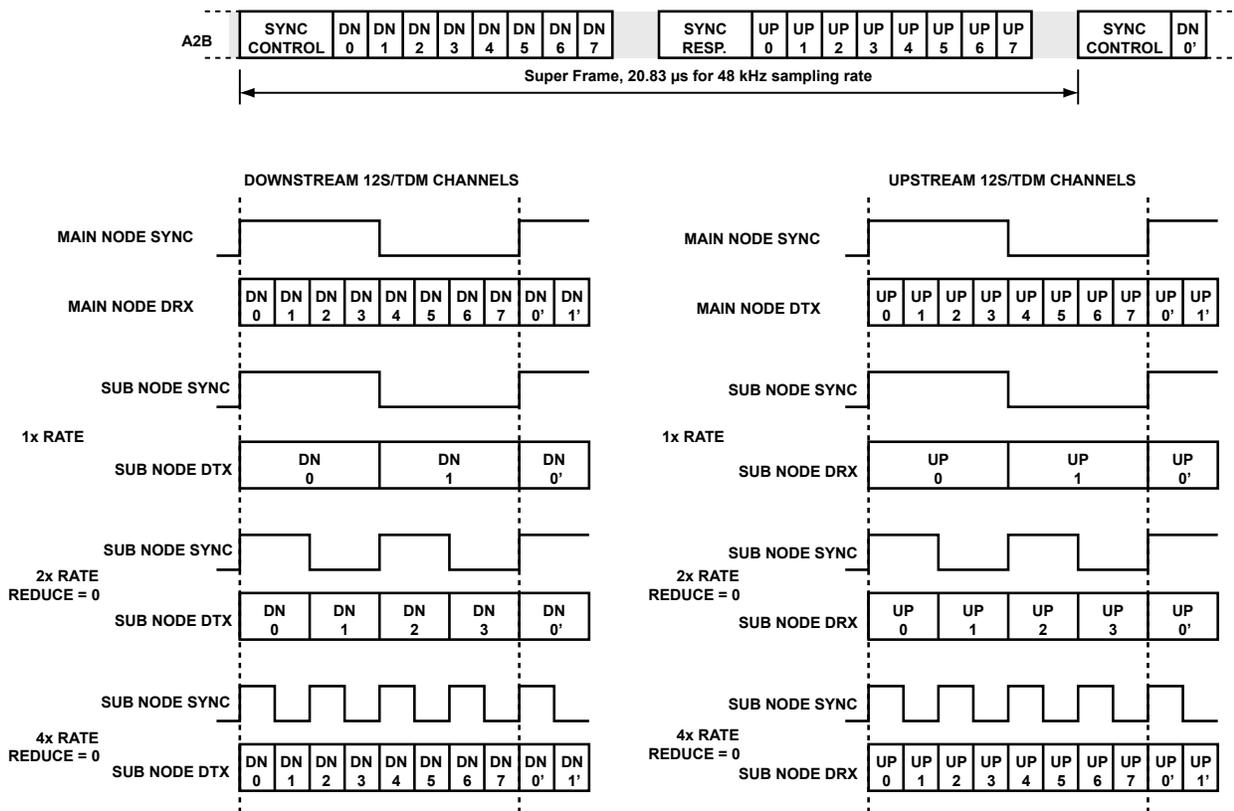


Figure 3-15: Increased Data Rate

The *Increased Data Rate Example* figure further illustrates the behavior of the `A2B_I2SRATE` register settings based on an example system. In the figure, both subordinate transceivers (S1 and S2) are set to $2 f_{\text{SYNCM}}$ rate mode. However, S1 has the `A2B_I2SRATE.REDUCE` bit set to 1. The waveforms in the figure illustrate the effect of the `A2B_I2SRATE.REDUCE` bit for both upstream and downstream slots. When the `A2B_I2SRATE.REDUCE` bit is set, only the first two channels on the DRX0/DRX1 pin are used for the upstream slots, and the other two channels are ignored for $2 f_{\text{SYNCM}}$ rate. For the DTX0/DTX1 transmitter, the two local downstream slots are duplicated on the DTX0/DTX1 pins for a $2 f_{\text{SYNCM}}$ rate when the `A2B_I2SRATE.REDUCE` bit is set.

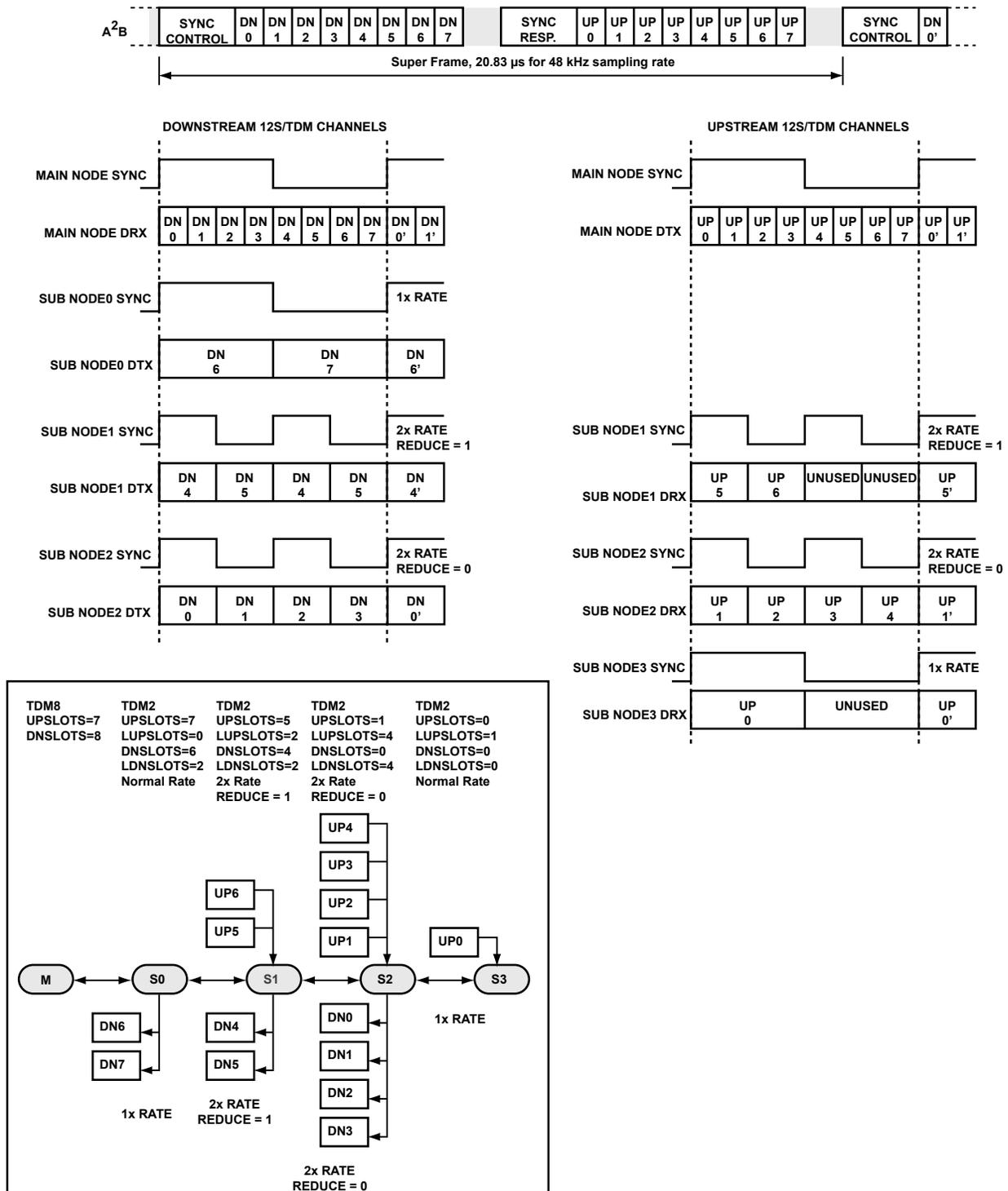


Figure 3-16: Increased Data Rate Example

NOTE: The increased data rate feature is not a sample rate converter. The mapping of frame buffer entries to TDM slots happens sequentially, so increased rate TDM requires enough bus slots to get the data aligned.

Alternatively, the TX crossbar feature can be used produce a expected TX mapping that looks like sample conversion.

GPIO Over Distance

This feature allows GPIO communication to occur over the A²B bus without host intervention after initial programming. The host is only required to initialize the GPIO over distance feature through the use of virtual ports. The GPIO over distance functionality has the following features:

- Eight parallel 1-bit virtual ports, managed by the main node. The main node can read the state of each virtual port can be read in the [A2B_GPIODDAT](#) register.
- Flexible mapping scheme of GPIO pins to virtual ports 0 through 7.
- GPIO pins can be configured as inputs that update the content of the [A2B_GPIODDAT](#) register or as outputs that reflect the content of the [A2B_GPIODDAT](#) register.
- When multiple virtual ports are mapped to one GPIO output pin, the values are OR'ed together.
- When multiple GPIO input pins are mapped to one virtual port, the values are OR'ed together even if they are from multiple nodes.

Configuration

Before attempting to configure the GPIO over distance function on a given pin, first verify that it is available for GPIO, as shown in the *GPIO Pin Configuration* table.

Table 3-9: GPIO Pin Configuration

| IO Bit | Pin Name | Pin Available for GPIO in Main Node | Pin Available for GPIO in Sub Node |
|--------|------------|---|------------------------------------|
| IO0 | IRQ/IO0 | Never | Always |
| IO1 | ADR1/IO1 | If <code>A2B_CLK1CFG.CLK1EN = 0</code> | |
| IO2 | ADR2/IO2 | If <code>A2B_CLK2CFG.CLK2EN = 0</code> | |
| IO3 | DTX0/IO3 | If <code>A2B_I2SCFG.TX0EN = 0</code> | |
| IO4 | DTX1/IO4 | If <code>A2B_I2SCFG.TX1EN = A2B_I2SGCFG.RXONDTX1 = 0</code> | |
| IO5 | DRX0/IO5 | If <code>A2B_I2SCFG.RX0EN = A2B_PDMCTL.PDM0EN = 0</code> | |
| IO6 | DRX1/IO6 | If <code>A2B_I2SCFG.RX1EN = A2B_PDMCTL.PDM1EN = 0</code> | |
| IO7 | PDMCLK/IO7 | If <code>A2B_PDMCTL.PDM0EN = A2B_PDMCTL.PDM1EN = A2B_PDMCTL2.PDMALTCLK = A2B_I2SRRCTL.ENSTRB = 0</code> | |

If the pin is available as GPIO, GPIO over distance is enabled by setting the appropriate enable bit in the GPIO over distance enable ([A2B_GPIODEN](#)) register. When a bit is set, the corresponding GPIO pin can then be mapped to one or more GPIO over distance virtual ports using the GPIO over distance mask registers

([A2B_GPIOD0MSK](#) through [A2B_GPIOD7MSK](#) , corresponding to GPIO-capable pins IO0 through IO7, respectively). Bits 0 through 7 in these registers correspond to virtual ports 0 through 7, respectively. If a bit is set within one of these registers, it maps the GPIO pin associated with the register to the corresponding virtual port.

If GPIO over distance is enabled for a given GPIO-capable pin, the direction of the pin is controlled exclusively via the GPIO output enable register ([A2B_GPIOOEN](#)) rather than a combination of this register and the complementary GPIO input enable register ([A2B_GPIOIEN](#)). When a bit in the [A2B_GPIOOEN](#) register is set, the associated GPIO pin is an output for GPIO over distance. If the bit is cleared in the [A2B_GPIOOEN](#) register, the associated GPIO pin is an input to GPIO over distance. It is not necessary to program the [A2B_GPIOIEN](#) register when using GPIO over distance for the pins of interest.

If the GPIO pin is an input (the associated bit in [A2B_GPIOOEN](#) = 0), the local node updates the virtual port(s) associated with the set bit(s) in the corresponding GPIO over distance mask registers ([A2B_GPIOD0MSK.IOD0MSK](#) through [A2B_GPIOD7MSK.IOD7MSK](#)). The virtual port values can be read in the GPIO over distance data register ([A2B_GPIODDAT](#)).

If the GPIO pin is an output (the associated bit in [A2B_GPIOOEN](#) = 1), the virtual ports that are mapped to that pin, as determined by the set bits in the associated GPIO over distance mask registers (through [A2B_GPIOD7MSK.IOD7MSK](#)) are OR'ed together to produce the GPIO output value (the logic OR of the corresponding bits in the [A2B_GPIODDAT](#) register).

NOTE: The [A2B_GPIODDAT](#) register is read only. It is recommended that the host always read this register from the main node.

The GPIO over distance inversion register ([A2B_GPIODINV](#)) allows for inversion of GPIO pin input or output. When a bit is set in this register, the associated GPIO pin signal is inverted. The inversion is applied on the way in from the pin if the GPIO pin is an input to a virtual port (the associated bit in [A2B_GPIOOEN](#) = 0), and it is applied on the way out to the GPIO pin if the pin is an output from a virtual port (the associated bit in [A2B_GPIOOEN](#) = 1).

If multiple nodes are updating the same virtual port, the [A2B_GPIODINV](#) register settings can be used to change the behavior from wired OR to wired AND. For example, to create a wired AND of multiple, active-high GPIO bits, the GPIO inputs and GPIO outputs must be inverted.

Mapping Multiple GPIO Inputs to One Virtual Port

When more than one node has a GPIO input mapped to the same virtual port, the protocol treats the input pins as a wired OR into the virtual port. When the virtual port is low (inactive), any request to set the virtual port results in a command from the main node to update all of the [A2B_GPIODDAT](#) registers across the system.

When the virtual port is high (active), any request to clear the virtual port results in a special command from the main node to notify all of the subordinate nodes of the request. If any of the subordinate nodes reject the request, the main node sees the rejection of the request, and the [A2B_GPIODDAT](#) registers retain their values. If none of the subordinate nodes reject the request, the main node sees an acceptance of the request and follows up with the updated [A2B_GPIODDAT](#) value.

GPIO Over Distance Programming Examples

The following procedures describe pin mapping cases to use GPIO over distance.

NOTE: Programming GPIO over distance must be done after the nodes have been discovered. For more information on node discovery, see the [Simple Discovery Flow](#) and [Appendix A: Additional Discovery Flow Examples](#) sections.

Mapping the Main Node DRX1/IO6 Pin to the Subordinate Node 2 ADR1/IO1 Pin

The following procedure describes how to map the main node DRX1/IO6 pin to the subordinate 2 ADR1/IO1 pin.

1. Write 0x04 to the main node [A2B_GPIOD6MSK](#) register to map the DRX1/IO6 pin to virtual port 2.
2. Write 0x40 to the main node [A2B_GPIODEN](#) register to enable GPIO over distance access on the DRX1/IO6 pin.
3. Write 0x02 to the subordinate node 2 [A2B_GPIOOEN](#) register to enable GPIO output for the ADR1/IO1 pin.
4. Write 0x04 to the subordinate node 2 [A2B_GPIOD1MSK](#) register to map virtual port 2 to the ADR1/IO1 pin.
5. Write 0x02 to the subordinate node 2 [A2B_GPIODEN](#) register to enable GPIO over distance access on the ADR1/IO1 pin.

Mapping the Subordinate Node 1 DTX1/IO4 Pin to the Main Node ADR1/IO1 Pin

The following procedure describes how to map the subordinate 1 DTX1/IO4 pin to the main node ADR1/IO1 pin.

1. Write 0x10 to the subordinate node 1 [A2B_GPIOD4MSK](#) register to map the DTX1/IO4 pin to bus GPIO port 4.
2. Write 0x10 to the subordinate node 1 [A2B_GPIODEN](#) register to enable GPIO over distance access on the DTX1/IO4 pin.
3. Write 0x02 to the main node [A2B_GPIOOEN](#) register to enable GPIO output for the ADR1/IO1 pin.
4. Write 0x10 to the main node [A2B_GPIOD1MSK](#) register to map bus GPIO port 4 to the ADR1/IO1 pin.
5. Write 0x02 to the main node [A2B_GPIODEN](#) register to enable GPIO over distance access on the ADR1/IO1 pin.

Mapping the ADR1/IO1 Pins on Subordinate Node 0 Through 2 to the Main Node ADR1/IO1 Pin

The following procedure describes how to map the ADR1/IO1 pin on subordinates 0 through 2 to the main node ADR1/IO1 pin.

1. For subordinate nodes 2, 1, and 0, write 0x01 to the `A2B_GPIOD1MSK` register to map the ADR1/IO1 pin of each subordinate to bus GPIO port 0.
2. For subordinate nodes 2, 1, and 0, write 0x02 to the `A2B_GPIODEN` register to enable GPIO over distance access on the ADR1/IO1 pin of each subordinate.
3. Write 0x02 to the main node `A2B_GPIOOEN` register to enable GPIO output for the ADR1/IO1 pin.
4. Write 0x01 to the main node `A2B_GPIOD1MSK` register to map bus GPIO port 0 to the ADR1/IO1 pin.
5. Write 0x02 to the main node `A2B_GPIODEN` register to enable GPIO over distance access on the ADR1/IO1 pin.

Transceiver Identification

Every A²B transceiver has a vendor ID register (`A2B_VENDOR`), a product ID register (`A2B_PRODUCT`), and a version ID (`A2B_VERSION`) register to indicate to a host which A²B transceivers are present in a system. Every A²B transceiver vendor is assigned a unique vendor ID (Analog Devices A²B transceivers use 0xAD as the vendor ID). The `A2B_PRODUCT` and `A2B_VERSION` registers are assigned by the chip vendor to uniquely identify the chips and indicate A²B interoperability. The transceiver models use 0x26 (AD2426), 0x27 (AD2427), and 0x28 (AD2428) as their product ID.

Every A²B transceiver also has a `A2B_CAPABILITY` register to identify available control interfaces and, as such, the presence of an I²C interface (`A2B_CAPABILITY.I2CAVAIL=1`).

Auto-Configuration System Information in EEPROM

In an A²B system, the supplier and specific product ID of each A²B node can be determined for auto-configuration if the subordinate modules contain a configuration memory (I²C EEPROM) with organization and content as described in [Appendix C: Module ID and Module Configuration Memory](#). Use auto-configuration for discovery when the host has no prior knowledge of the exact system configuration. Specific configuration commands for a subordinate node can also be stored in the configuration memory by using the optional configuration blocks.

Standby Mode

In standby mode, there is no upstream traffic on the A²B bus. Only a minimal (19-bit) SCF exists to keep all of the subordinate nodes synchronized, and there is no SRF. Header count errors and CRC errors are ignored, and data slots are disabled. GPIO settings retain their values while in standby mode.

While in normal mode, the host can write to the main transceiver `A2B_DATCTL` register to go to standby mode, but the write does not take effect until a new structure is applied to the system. The host performs the following actions:

1. Set the `A2B_DATCTL.STANDBY` bit in the main transceiver to generate a broadcast write of 0x80 to set the `A2B_DATCTL.STANDBY` bit in all of the discovered subordinate nodes. Writing 0x80 to the `A2B_DATCTL` register ensures that the data slots are disabled.

- Set the `A2B_CONTROL.NEWSTRCT` bit in the main transceiver to apply the new structure.

After the new structure is applied, the system transitions into standby mode. The host can move the system back to normal mode by writing 0x00 to the `A2B_DATCTL` register in the main node. This instruction generates a broadcast write of 0x00 to the `A2B_DATCTL` register in all of the subordinate nodes. The main node provides the standby done interrupt to the host (`A2B_INTTYPE = 0xFE`) when the system is back in normal mode.

Bus Monitor Support

Bus monitor mode enables the transceiver to act as a passive automotive audio bus monitor, also referred to as a *sniffer*. The A²B test equipment uses this mode. Only the host processor can allow bus monitors on A²B bus segments to monitor the synchronous data content. To permit this synchronous data monitoring, the host must set the `A2B_DATCTL.ENDSNIFF` bit in the main transceiver. This configuration triggers an A²B bus broadcast of the information to the attached bus monitor devices.

The *Bus Monitor Behavior* figure shows a bus monitor node inserted between subordinates 0 and 1 in an A²B network.

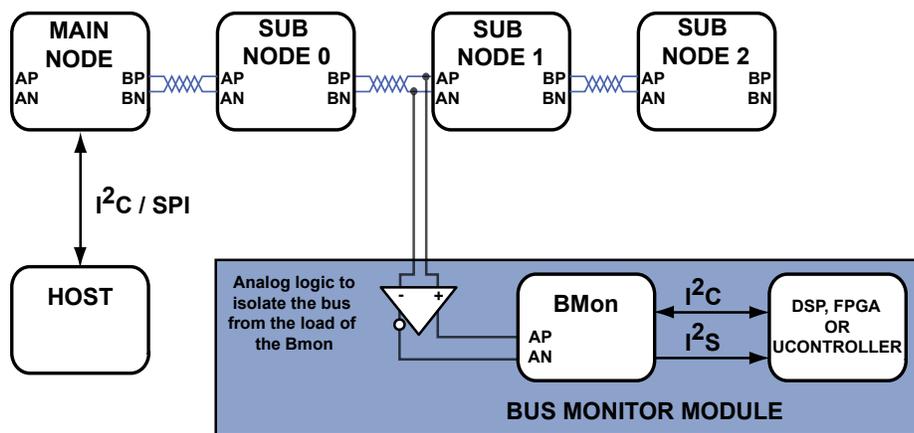


Figure 3-17: Bus Monitor Behavior

A bus monitor is passive in the system; it does not respond to bus synchronization control frames (SCFs) or contribute any data to the bus. It only uses the A-side transceiver while the B-side transceiver is deactivated. When in bus monitor mode, the transceiver synchronizes itself to SCFs and may snoop SCF control writes to configure its bus interface to match the downstream node being monitored. The A²B bus monitor transceiver uses its I²S/TDM port to transmit A²B bus traffic to a protocol analyzer circuit.

A bus monitor node behaves as follows:

- The B-Side (downstream) transceiver is disabled.
- The A-Side (upstream) transceiver is enabled to receive only (not to transmit).
- SRF generation is disabled.

4. The I²S/TDM interface is configured for 32-bit data width:

- Downstream SCFs are transmitted on the DTX0 pin
- Upstream SRFs are transmitted on the DTX1 pin
- Data slot bits can only stream out of the DTXn pins if the A²B bus main is programmed to enable this feature
 - Downstream slots are streamed out on the DTX0 pin
 - Upstream slots are streamed out on the DTX1 pin
 - If there are more data slots on the A²B bus than there are available I²S/TDM channels, then a programmable offset determines which data slots to monitor on the I²S/TDM channels

NOTE: When the bus monitor receiver is disabled, an external switch must be used to control the LVDS traffic going to the A-side of a transceiver in bus monitor mode.

A bus monitor that attaches to an A²B bus after discovery and initialization can miss the broadcast and therefore has the monitoring of synchronous data slots disabled. The preferred method is to attach bus monitors before initialization and discovery. Alternatively, for full support of bus monitors that must see data slots but attach after discovery, the host can perform regular writes to the [A2B_DATCTL](#) register to generate the enable data slot sniffing broadcasts. The bus monitor node microcontroller must set the `A2B_BMMCFG.BMMEN` bit to enable bus monitor mode and can further configure the bus monitor transceiver when attaching to and detaching from the A²B bus:

- The `A2B_BMMCFG.BMMNDSC` bit determines whether the bus monitor attaches before or after system discovery and initialization. When cleared (= 0), the monitor attaches before A²B discovery, so the discovery sequence sets the bus timing properties automatically. When set (= 1), the bus timing properties must be set by the bus monitor node microcontroller using local I²C register writes.
- The `A2B_BMMCFG.BMMRXEN` bit is used to keep the LVDS A-side transceiver input static while the bus monitor is being attached. It is also used to reinitiate the bus monitor lock sequence without physically detaching the bus monitor node.

Besides configuring and enabling bus monitor mode in the [A2B_BMMCFG](#) register, the use of bus monitor mode affects the meaning and settings of bits in the following A²B registers:

- I²S Global Configuration register ([A2B_I2SGCFG](#))
 - The `A2B_I2SGCFG.INV`, `A2B_I2SGCFG.EARLY` and `A2B_I2SGCFG.ALT` bits must be programmed to match the interface of the protocol analyzer
 - The `A2B_I2SGCFG.TDMSS` bit must be programmed to 0 for 32-bit TDM slot size
 - The `A2B_I2SGCFG.TDMMODE` field must be set to match the protocol analyzer's capabilities:
 - TDM2 allows monitoring of SCF and SRF frames

- TDM4 allows monitoring of SCF and SRF frames, as well as up to two upstream and two downstream data slots simultaneously
- TDM8 allows parallel monitoring of SCF and SRF frames, as well as up to six upstream and six downstream data slots simultaneously
- TDM16 allows parallel monitoring of SCF and SRF frames, as well as up to 14 upstream and 14 downstream data slots simultaneously
- TDM32 allows parallel monitoring of SCF and SRF frames, as well as up to 30 upstream and 30 downstream data slots simultaneously
- I²S Configuration register ([A2B_I2SCFG](#))
 - Setting the `A2B_I2SCFG.TX0EN` bit enables output of downstream data on the DTX0 pin
 - Setting the `A2B_I2SCFG.TX1EN` bit enables output of upstream data on the DTX1 pin
 - Set the `A2B_I2SCFG.TXBCLKINV` bit to match the interface of the protocol analyzer
 - The `A2B_I2SCFG.TX2PINTL`, `A2B_I2SCFG.RXBCLKINV`, and `A2B_I2SCFG.RX0EN` bits must be programmed to 0
- Local Upstream Slots Offset register ([A2B_UPOFFSET](#)) – determines the offset in number of data slots between upstream data slots received on the A²B bus and upstream data slots driven onto the DTX1 pin as I²S/TDM channels. The register programs between monitoring of higher or lower index slots if the number of slots exceeds the number of transmit channels available in the selected TDM format.
- Local Downstream Slots Offset register ([A2B_DNOFFSET](#)) – determines the offset in number of data slots between the downstream data slots received on the A²B bus and the downstream data slots driven onto the DTX0 pin as I²S/TDM channels. The register programs between monitoring of higher or lower index slots if the number of slots exceeds the number of transmit channels available in the selected TDM format.

I²S/TDM Channel Format

In standby mode, there is no upstream traffic on the A²B bus. Only a minimal (19-bit) SCF exists to keep all of the subordinate nodes synchronized, and there is no SRF. Header count errors and CRC errors are ignored, and data slots are disabled. GPIO settings retain their values while in standby mode.

The following examples describe the I²S/TDM output format in bus monitor mode.

The DTX0 pin transmits in the first two 32-bit I²S/TDM transmit channels' downstream frame status bits, followed by the downstream control frame information. Further I²S/TDM channels, if available and allowed, carry the downstream synchronous data. The [A2B_DNOFFSET](#) register provides an offset between the downstream data slots and the data slots produced on DTX0.

The DTX1 pin transmits in the first two 32-bit I²S/TDM transmit channels' upstream frame status bits, followed by the upstream response frame information. Further I²S/TDM channels, if available and allowed, carry the

upstream synchronous data. The `A2B_UPOFFSET` register provides an offset between upstream data slots and the data slots produced on DTX1.

During discovery and initialization, the host programs the data slot format register (`A2B_SLOTFMT`) in the main transceiver, which auto-broadcasts this information to the subordinates. An attached bus monitor can listen to this control message and derive the slot size settings (32 bits maximum).

Data is always transmitted MSB-aligned.

The parity bits are not included in the I²S/TDM channels, but the `A2B_ERRMGMT` register may be used to indicate data slot errors in the LSB of the data, below the LSB of the data, or in an additional error channel following the data channels.

The *TDM16 Downstream Example (DTX0 Pin) Registers* figure shows the downstream data produced for monitoring on the DTX0 pin with `A2B_I2SGCFG.TDMMODE = TDM16`.

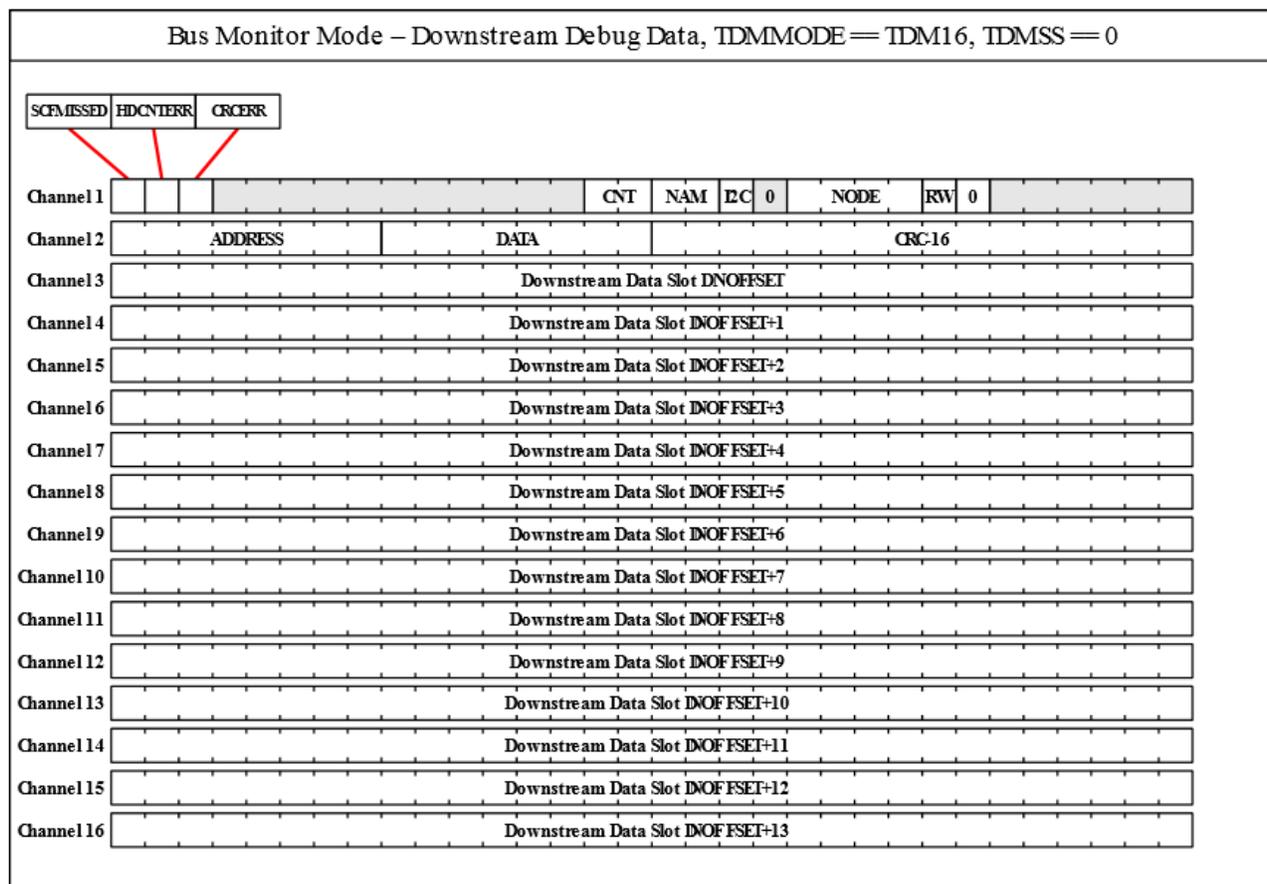


Figure 3-18: TDM16 Downstream Example (DTX0 Pin) Registers

The *TDM16 Upstream Example (DTX1 Pin) Registers* figure shows the upstream data produced for monitoring on the DTX1 pin with `A2B_I2SGCFG.TDMMODE = TDM16`.

6. The I²S/TDM interface starts transmitting after the bus monitor node locks its PLL. The IRQ pin on the bus monitor node goes high to indicate that the node found lock. This event should occur before the next-in-line node starts responding.
7. Writes to the `A2B_BCDNSLOTS`, `A2B_LDNSLOTS`, `A2B_LUPSLOTS`, `A2B_DNSLOTS`, `A2B_UPSLOTS`, `A2B_SLOTFMTA2B_DATCTL`, `A2B_TESTMODE`, `A2B_I2SRRATE`, `A2B_I2SRRCTL`, `A2B_UPMASK0` through `A2B_UPMASK3`, and `A2B_DNMASK0` through `A2B_DNMASK3` registers in the next-in-line subordinate node on the probed bus segment are mirrored in the bus monitor node, where they are locally accessible over the I²C interface. Application of a new data structure on the bus (when the host sets the `A2B_CONTROL.NEWSTRCT` bit in the main transceiver) is also applied to the bus monitor node.
8. The DTX[1:0] pins do not transmit data slot content unless the bus monitor has seen a broadcast write resulting from the host setting the `A2B_DATCTL.ENDSNIFF` bit in the main transceiver.

After Discovery

When `A2B_BMMCFG.BMMNDSC = 1`, the following sequence of events occurs:

1. The downstream subordinate node of the probed bus segment is already DC-biased and discovered.
2. Set `A2B_BMMCFG.BMMEN = A2B_BMMCFG.BMMNDSC = 1` through the I²C interface.
3. Physically attach the bus monitor to the bus segment (probe point).
4. Set `A2B_BMMCFG.BMMRXEN = 1` through the I²C interface. After the bus monitor transceiver correctly locks to the SCFs, the IRQ pin goes high.
5. Initialize the `A2B_RESPCYCS` register to 0x20 through the I²C interface. The appropriate value for `A2B_RESPCYCS` is determined from the SRF timing and updates automatically.
6. Configure I²S/TDM transmit settings through the I²C interface in the `A2B_I2SGCFG`, `A2B_I2SCFG`, `A2B_I2SRATE`, `A2B_SYNCOFFSET` and `A2B_ERRMGMT` registers to match the desired timing and format characteristics.
7. If the monitoring of control and response frames alone is desired, this step can be skipped. If monitoring of data slots is desired (and the host allows access to them), configure the `A2B_DNSLOTS`, `A2B_UPSLOTS`, `A2B_SLOTFMT` and `A2B_DATCTL` registers through the I²C interface. The correct values for these registers can come from values previously stored in memory after sniffing the same bus segment during discovery and initialization. If values are completely unknown, then software can try different values to find suitable settings.
 - The `A2B_DNSLOTS` register represents the number of downstream data slots at the A-side transceiver circuit of the next-in-line downstream subordinate.
 - The `A2B_UPSLOTS` register represents the number of upstream data slots at the A-side transceiver circuit of the next-in-line downstream subordinate.
 - The `A2B_SLOTFMT` register represents the data slot format.

- The `A2B_DATCTL.DNS` and `A2B_DATCTL.UPS` bits must match the committed values in the downstream subordinate node. The DTX0 and DTX1 pins do not transmit data slots in I²S/TDM channels if these bits are not set.
8. The DTX[1:0] pins do not transmit data slot content unless the bus monitor has seen a broadcast write resulting from the host setting the `A2B_DATCTL.ENDSNIFF` bit in the main transceiver.

Optimizing EMC Performance

EMC performance is critical in an A²B transceiver system design. The transceivers have several programmable features that can be utilized to optimize EMC performance:

- [Spread-Spectrum Clocking](#)
- [Programmable LVDS Transmit Levels](#)
- [Data-Only and Power-Only Bus Operation](#)

Spread-Spectrum Clocking

Spread-spectrum clocking can be used to reduce narrowband emissions on a PCB. By default, spread-spectrum clocking is disabled on the transceiver, but writes to the `A2B_PLLCTL` register can enable spread-spectrum clocking during discovery. The `A2B_PLLCTL` register contains settings that enable spread-spectrum clocking for clocks that are internal to the transceiver.

If spread-spectrum clocking support is enabled for the internal clocks, spread-spectrum clocking can also be enabled for both the I²S interface and the programmed CLKOUTs. Enabling spread-spectrum clocking for internal clocks, CLKOUTs, and the I²S interface may reduce narrowband emissions by several dB on a particular node.

ATTENTION: When spread-spectrum clocking is enabled on a clock output, the TIE jitter on that clock increases.

To enable an A²B network with spread-spectrum clocking, all nodes must be set to the same depth and frequency. Follow this sequence to set the nodes:

1. Discover all subordinates.
2. Configure spread-spectrum for all nodes (including the main node) with a broadcast write to the `A2B_PLLCTL` register of each node.

For a single node with spread spectrum (including systems with the AD2421/AD2422/AD2425 models), follow this sequence:

1. Discover all nodes.
2. Configure spread spectrum (by setting the `A2B_PLLCTL` register) for each subordinate, one at a time.
 - a. The `A2B_PLLCTL.SSDEPTH` bit is limited to setting 0x0.
 - b. Adjacent nodes must have the same `A2B_PLLCTL.SSFREQ` setting.

NOTE: A broadcast write to the `A2B_PLLCTL` register is mandatory when all the nodes in the system are enabled with spread spectrum. Set the `A2B_NODEADR.BRCST` bit and initiate a write to the `A2B_PLLCTL` register with `A2B_BUS_ADDR`. A broadcast write effects all nodes. It occurs in the main node first and then in the subordinate nodes during the next SCF.

Sequential programming of spread spectrum must follow the single node guidelines. The `A2B_PLLCTL.SSDEPTH` bit is limited to setting 0x0 for sequential programming of spread spectrum clocking, as well as in a system with a single node with spread spectrum clocking enabled.

The `A2B_PLLCTL.SSMODE` field can be set to protocol only or I²S+ protocol, whether or not spread spectrum clocking is enabled.

Programmable LVDS Transmit Levels

The LVDS transmitter can be set to transmit the signal at high, medium, or low levels. Higher transmit levels yield greater immunity to EMI, while lower transmit levels can reduce emissions from the twisted-pair cables that link A²B bus nodes together.

The LVDS transmit levels can be changed by adjusting the settings in the `A2B_TXACTL` (A-side) or `A2B_TXBCTL` (B-side) register. If a non-default transmit level is desired, `A2B_TXxCTL` must be written on each node (during discovery) before setting the `A2B_SWCTL.ENSX` bit. The `A2B_TXACTL.TXAOVREN` enable bit must be set in order for the `TXxLEVEL` setting to take effect.

Data-Only and Power-Only Bus Operation

The A²B bus can be operated without closing the switch to send a DC bias downstream. This requires that the `A2B_CONTROL.SWBYP` bit is set instead of the `A2B_SWCTL.ENSX` bit during discovery.

Conversely, the `A2B_SWCTL.DISNXT` bit allows a DC bias to be sent downstream without the presence of data. This setting should be applied at the same time as the write to set the `A2B_SWCTL.ENSX` bit during discovery. These modes are used primarily for debug purposes.

Cross-Over or Straight-Through Cabling

Straight-through cables can be supported by swapping the DC-coupling circuitry at the B-side connector. For hardware designed to support straight-through cables, the `A2B_CONTROL.XCVRBINV` bit must be set during discovery to ensure proper operation. This is done before setting the `A2B_SWCTL.ENSX` bit for each subordinate node that is connected with a straight-through cable.

IMPORTANT: Ensure that the `A2B_CONTROL.XCVRBINV` bit is not overwritten while doing other operations such as writing to the `A2B_CONTROL.NEWSTRCT` bit (which applies a new structure).

4 A²B Event Control

The A²B protocol engine contains a set of registers that provide support for interrupts to the host. These registers include:

- `A2B_INTSTAT`
- `A2B_INTSRC`
- `A2B_INTTYPE`
- `A2B_INTPND0` through `A2B_INTPND2`
- `A2B_INTMSK0` through `A2B_INTMSK2`

To register subordinate interrupt requests in the main node, unmask the subordinate node interrupts in the `A2B_INTMSK0` and `A2B_INTMSK1` registers. In main nodes only, also unmask interrupts in the `A2B_INTMSK2` register.

The active polarity of the `A2B_IRQ` pin is set using the `A2B_PINCFG` register. By default, interrupt requests are indicated with a high level on the `A2B_IRQ` pin and the setting of the `A2B_INTSTAT.IRQ` bit. An active interrupt request in the main transceiver is cleared and revised on a host read of the main transceiver `A2B_INTTYPE` register. This process also applies to the main node receiving an interrupt request from a subordinate node.

The main transceiver register (`A2B_INTSRC`) indicates whether the active interrupt is generated by the main node or by a subordinate node (where it also supplies the ID of the subordinate node). The `A2B_INTTYPE` register in the main transceiver contains information that the host uses to determine the interrupt cause. Priority logic automatically determines the value of the `A2B_INTSRC` and `A2B_INTTYPE` registers. Other pending interrupt requests can appear after reading the `A2B_INTTYPE` register. The `A2B_IRQ` pin goes low for one f_{SYSBCLK} cycle (~20 ns) when the `A2B_INTTYPE` register is read. The pin immediately transitions to high if there are pending interrupt requests.

When masked interrupts occur, they are registered as sticky bits in the `A2B_INTPND0` through `A2B_INTPND2` registers but do not trigger interrupt requests. Once unmasked, any pending interrupts trigger interrupt requests following this order of priority:

- Main node interrupts have priority over subordinate node interrupts.
- Lower subordinate node ID numbers take priority over higher numbers.

- Lower number `A2B_INTTYPE` has priority over higher number.
 - `A2B_INTPND0` takes priority over `A2B_INTPND1`, which takes priority over `A2B_INTPND2`.
 - Lower numbered bits in the pending registers `A2B_INTPND0` to `A2B_INTPND2` take priority over higher numbered bits.

The IRQ signal is immediately asserted when the main transceiver receives an interrupt request from a subordinate.

Host Response to Interrupt Requests

When the host receives an interrupt request from the main node (indicated by the IRQ signal going high), the host can read the `A2B_INTSRC` and `A2B_INTTYPE` registers to obtain the subordinate node ID that generated the interrupt request and the type of interrupt request, respectively. This can be accomplished by performing a single 2-byte read, starting at the `A2B_INTSRC` address, which reads both registers. At the completion of the `A2B_INTTYPE` register read, the active interrupt is cleared and IRQ goes low if there are no further pending interrupts.

Interrupt Latency

Interrupts are signaled upstream from a subordinate transceiver to the main transceiver within the Synchronization Response Frame (SRF). Interrupts that engage after the beginning of the SRF (after the subordinate node starts driving the AP and AN pins) are signaled to the main node in the SRF of the next superframe. Assuming there are no other interrupts with a higher priority that mask the IO pin interrupt in question, the latency between a subordinate node IO pin and main node IRQ is the sum of:

- Four SYSBCLK cycles for pin interrupt generation (81.4 ns) +
- One superframe latency to get into SRF (20,833.3 μ s) +
- 64 SYSBCLK cycles for the length of the SRF (1,302.1 ns) +
- Five SYSBCLK cycles for main Rx latency (101.7 ns) +
- Two SYSBCLK cycles for IRQ logic in the main node (40.7 ns)

In addition to this total latency of 22.36 μ s, there is an additional nine SYSBCLK cycles of latency for each subordinate that the SRF must pass through ($N \times 183.1$ ns). For example, in a system with three subordinates, a GPIO interrupt from subordinate 2 to the main node has a maximum latency of $22.36 \mu\text{s} + (2 \times 0.183) \mu\text{s} = 22.73 \mu\text{s}$.

Error Management

The following sections provide information about error management. All data transmitted over the A²B bus is checked for line code violations (DDERR) at the receiving end. Additionally, the SCF and SRF use cyclic redundant codes (CRC), and every synchronous data slot uses a parity bit for extra error detection certainty, as shown in the *Frame Structure Details* figure.

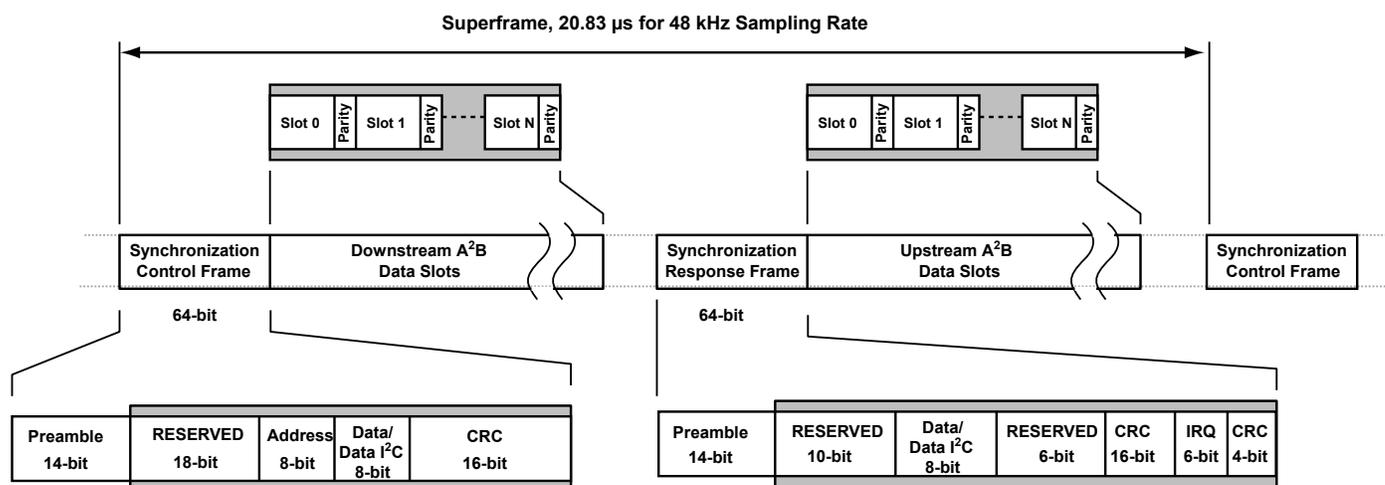


Figure 4-1: Frame Structure Details

Downstream Data Error Detection

A 16-bit cyclic redundant code (CRC) is part of any downstream control data inside the SCF. This CRC determines on the receiving side if the SCF data is corrupted during transmission.

The SCF has a preamble to indicate the start of a superframe. It provides a bit pattern that subordinate nodes use for clock and frame synchronization. If this frame sync is not detected by a subordinate node, the error is treated as a CRC error.

Upstream Data Error Detection

A 16-bit cyclic redundant code is part of any upstream response data inside the SRF. The CRC determines on the receiving side if the SRF data is corrupted during transmission. Interrupt request fields have an extra CRC (ICRC) inside the SRF to prevent incorrect interrupts from being triggered.

The SRF has a preamble to indicate the start of the response frame. It provides a bit pattern that is used for clock and frame synchronization. If this frame sync is not detected by an upstream node, the error is captured as an SRFCRCERR in the subordinate nodes and as a CRCERR in the main node.

Data Slot Error Correction

Possible cases for automatic correction of received data slots in a node are described as follows.

- If the frame sync preamble is not seen, all data slots received from the bus are automatically replaced with previous, good values.
- If a CRC error is detected in the SCF by a subordinate (`A2B_INT_PND0.CRCERR = 1`), all downstream data slots received from the bus are replaced with previous, good values.

- If a CRC error is detected in the SRF by the main node (`A2B_INTPND0.CRCERR = 1`), all upstream data slots received from the bus are replaced with previous good values.
- If a data decoding error (`A2B_INTPND0.DDERR = 1`) or a data parity error (`A2B_INTPND0.DPERR = 1`) is detected within a data slot, the received erroneous data slot is replaced automatically with a previous, good slot value.

Control and Response Error Handling

When a host accesses registers over I²C and A²B (for example, I²C over distance), the synchronization control frame (SCF) and the synchronization response frame (SRF) carry this data exchange. If there is a communication error in the control frame or the response frame, the main node automatically initiates a retry of the register access. The main node retries multiple times until either the access is successful or an I²C timeout occurs in the main node. During the retry time, I²C clock stretching is applied, which signals to the host that the transaction is not completed. If there is an I²C timeout (the I²C timeout occurs after 30 superframes), the main node flags an I2CERR interrupt, to which the host can respond.

Corrupted received interrupt requests in the main node are ignored. If a real interrupt event occurs, an interrupt is automatically re-generated by the subordinate node since it is not cleared.

Error Signaling

Any communication error, flagged in the `A2B_INTPND0` or `A2B_INTPND2` register, triggers an interrupt request when the corresponding interrupt is enabled in the `A2B_INTMSK0` or `A2B_INTMSK2` register, respectively. These interrupt requests use the `A2B_IRQ` pin or the `A2B_INTSTAT.IRQ` bit to signal the request to the host. The host can then read the `A2B_INTSRC` and `A2B_INTTYPE` registers to determine what the error is and where it occurred.

The `A2B_BECCCTL` register selects which communication errors are counted and what counter threshold must be exceeded for an interrupt request to be generated. Using this feature, certain single-bit communication errors do not have to generate an interrupt unless they significantly accumulate over the time period since the `A2B_BECCNT` register was last cleared. Additionally, three different methods to signal data slot errors over the I²S/TDM interface can be selected using the `A2B_ERRMGMT` register (see [Error Management Register](#) for details).

A²B Communication and Bit Errors

The A²B communication and bit errors are:

- HDCNTERR (`A2B_INTTYPE = 0`)

The SCF and SRF fields contains a 2-bit field CNT. In the SCF, the CNT field is incremented (modulo 4) from the value used in the previous superframe. In the SRF, the received value of the CNT field in the SCF is transmitted back to the main node. HDCNTERR indicates that the current node has detected a header count error. For the main node, this means that the synchronization response frame has a different CNT value than expected. For a subordinate node, this means that the synchronization control frame has a different value than expected.

- DDERR (`A2B_INTTYPE = 1`)

The DDERR error indicates a missing clock edge in the Differential Manchester data stream on the A²B bus. The data decoding error is reported only on data slots that are being consumed by the particular node. A data decode error in an SCF/SRF results in a CRC error and does not raise a data decoding error.

- CRCERR (`A2B_INTTYPE = 2`)

The CRCERR error indicates that a subordinate node detects a CRC error in the received SCF field. For the main node, the error indicates a CRC error in the received SRF field.

- DPERR (`A2B_INTTYPE = 3`)

A data slot on the A²B bus is protected by a parity bit. The DPERR error is reported only on data slots that are consumed by the particular node. Nodes do not check parity for the slots that are just passed through it.

- BECOVF (`A2B_INTTYPE = 4`)

The `A2B_BECCTL.THRESHLD` field configures the number of bit errors (HDCNTERR, DDERR, CRCERR, DPERR, and ICRCERR) to be counted before the `A2B_INTPND0.BECOVF` bit is set. This threshold is useful if it is not desired to signal the interrupt for every bit error. The threshold can be set based on acceptable noise and robustness over a particular period of time. The bit error counter should be cleared periodically. Excessive bit errors set the `A2B_INTPND0.BECOVF` bit and signal the interrupt. Bit error thresholds can be independently set at main and subordinate nodes.

- SRFERR (`A2B_INTTYPE = 5`)

The SRFERR error indicates that the SRF of a later node is not received prior the local timing window being expired, and the affected node generates its own SRF, which is being up-streamed to any earlier nodes. The error is valid for main and subordinate nodes.

- SRFCRCERR (`A2B_INTTYPE = 6`)

The SRFCRCERR error indicates that the current subordinate node detected a CRC error in the SRF field. Usually, when a subordinate node detects a CRC error in the SRF, it flags the SRFCRCERR error bit of the node. The subordinate does not try to correct the error and passes the SRF as-is upward. But the exception is, in case of a response to a command, the subordinate node inserts its own SRF, including the CRC. In the main node, the CRCERR field is used to indicate a CRC error in the SRF.

- PWRERR (`A2B_INTTYPE = 9-15`)

PWRERR is the mask bit for errors from the DLPS (digital line power switch) block; LDO2 is used internally for powering the DLPS block.

Subordinate Node Interrupt Handling

This section describes how subordinate node interrupts are internally handled by the main node. When an interrupt occurs in a subordinate node, the following sequence of events happens in response:

1. After the subordinate node interrupt occurs, the related bits in the `A2B_INTPND0`, `A2B_INTPND1`, `A2B_MBOX0STAT`, and `A2B_MBOX1STAT` registers of the subordinate transceiver are set.
2. If the `A2B_INTSTAT . IRQ` bit is low, it gets set. The highest priority pending interrupt type is then written to the `A2B_INTTYPE` register.
3. The subordinate node begins signaling the IRQ in the interrupt field of the SRF. Any upstream subordinates without an active interrupt passes this field upstream.
4. When the main node receives an interrupt field with a valid CRC and the IRQ field set, then main node sets its `A2B_INTSTAT . IRQ` bit if the bit is not already set. The main then updates the `A2B_INTSRC` register with the subordinate number and sets the `A2B_INTTYPE` register to 0x80. At this point, the IRQ pin of the main node is driven active.
5. The main automatically reads the `A2B_INTTYPE` register from the appropriate subordinate and updates its `A2B_INTTYPE` register. This is held off if there is a new structure being applied (`A2B_CONTROL . NEWSTRCT` set within the last five superframes), or if a remote I²C stop command needs to be sent.
6. Once the `A2B_INTTYPE` is read, the main automatically performs a write to the appropriate subordinate to clear the interrupt. This is held off if there is a new structure being applied, or if a remote I²C stop command needs to be sent. At this point, the subordinate node stops signaling the interrupt in the SRF.
7. When the IRQ pin of the main node is asserted as a result of the subordinate interrupt, the host processor reads the `A2B_INTSTAT` and `A2B_INTTYPE` registers to ascertain the interrupt type and identify which subordinate node raised the interrupt.

If the host reads the `A2B_INTTYPE` register from the main node after step 4 but before step 5 completes, the host may read 0x80 from `A2B_INTTYPE`. If the subordinate node does not drop off the bus, the `A2B_INTTYPE` bit field eventually updates.

When the host reads the `A2B_INTTYPE` register = 0x80, an additional read of the `A2B_INTTYPE` register is recommended to confirm the interrupt type. If a subordinate signals an interrupt and then drops off the bus (presumably, due to a switch fault), the next upstream subordinate eventually switches to being the last subordinate after 32 frames of missed SRFs. At this point, if the main node (not the host processor) is still internally attempting to read `A2B_INTTYPE` from the missing subordinate node, the newly last subordinate sends a special SRF, indicating to the main node that the read cannot go through. This causes `A2B_INTTYPE` to be set to 0xFD and the interrupt identification process to terminate. Since the missed SRF timeout is 32 superframes (after which the upstream node becomes the last node), the error type 0xFD is unlikely.

In other words, the subordinate node `A2B_INTTYPE` read error (0xFD) interrupt occurs when the main node is attempting to read `A2B_INTTYPE` from a subordinate based on a received interrupt but receives a response from an upstream subordinate node indicating that subordinate is now the last subordinate node. The main difference between `A2B_INTTYPE` = 0xFD and `A2B_INTTYPE` = 0x80 is that `A2B_INTTYPE` = 0x80 can be seen while the main node is still attempting to read `A2B_INTTYPE`, so it may subsequently resolve, whereas when the `A2B_INTTYPE` bit field = 0xFD cannot resolve.

If a subordinate node just reports an interrupt to the main node, without any additional line failures after that. If after step 4, the host reads `A2B_INTTYPE` too fast, it reads `A2B_INTTYPE = 0x80`, resulting in the IRQ to be cleared. The main node does not reassert the IRQ if the `A2B_INTTYPE` register is read before the register value is updated from a subordinate. If the host reads the `A2B_INTSRC` register, then the `A2B_INTTYPE` register after seeing the IRQ (which is recommended), then the `A2B_INTTYPE` value is valid (unless there is a line error).

If a subordinate node with no pending interrupt disconnects from the rest of the bus, the upstream subordinate generates the SRFERR in 32 consecutive superframes.

Error Management Register

When A²B data slots are not received correctly (detected by a parity error or a data decode error on any bit in the slot), the last good sample received for that slot is repeated. The `A2B_ERRMGMT` register also controls the ways in which bad data slots can be indicated across the I²S/TDM interface.

When the `A2B_ERRMGMT.ERRLSB` bit is set, the LSB of each data slot is used to indicate whether the slot is received correctly or not. For example, in the main node with a 24-bit upstream slot size, the 24th data bit sent over DTX0 or DTX1 is low when the data is valid; it is high when the data is not valid. This method changes the meaning of the LSB in the received I²S/TDM data words.

When the `A2B_ERRMGMT.ERRSIG` bit is set, all bits below the LSB of each data slot are used to indicate whether the slot is received correctly. With a 24-bit slot size, the last 8 bits in each 32-cycle data slot are low when the data is valid; the bits are high when the data is not valid. If the `A2B_ERRMGMT.ERRSIG` bit is not set, the extra eight bits are always low. This method preserves the meaning of the LSB in the received I²S/TDM data words, but the data word size must be smaller than the data channel size for this method to work. Data channel width is usually 32 bits, but it can be programmed to 16 bits.

When the `A2B_ERRMGMT.ERRSLOT` bit is set, the number of slots generated on the A²B bus is incremented by 1. In the main node, the protocol engine normally writes `A2B_UPSLOTS` pieces of data to the frame buffer in each superframe. In a subordinate node, the number of slots written is normally `A2B_LDNSLOTS + A2B_BCDNSLOTS`. The additional data slot enabled by using this method is appended to the end of the configured A²B traffic. It contains a single bit of error information for each of the preceding data slots in that superframe. The MSB of the extra slot indicates that an error occurred in data slot 0. The next bit indicates an error in data slot 1, and so on. For example, `0x80000000` indicates that there was an error in slot 0, while `0xfffff00` indicates that slots 0 through 23 all contained errors. If the `A2B_I2SGCFG.TDMSS` bit is set for a channel size of 16 bits, only the first 16 data channels can be reported. If the `A2B_I2SGCFG.TDMSS` bit is set for a channel size of 32 bits, up to 32 data channels can be reported for errors.

Bit Error Control Register

The `A2B_BECONT` register controls bit error counting, including interrupt thresholds of 2^n , where n ranges from 1 to 8. It selects which communication errors enter a counter and at what counter-threshold an interrupt request is generated. Using this feature, certain single-bit communication errors do not have to generate an interrupt unless they significantly accumulate over the time period when the `A2B_BECONT` register is last cleared.

Testing and Debugging

For testing and debugging, the transceiver allows generation of interrupts and bit errors using the raise `A2B_RAISE` and generate error `A2B_GENERR` registers.

Raise (`A2B_RAISE`) Register

The `A2B_RAISE` register allows the host to generate an interrupt in any node in the system via software. The register must be written over the A²B bus, as writes to the register from the local I²C port have no effect.

Generate Error (`A2B_GENERR`) Register

- *0x01 Generate Header Count Error* (`A2B_GENERR.GENHCERR`)

1. When the main node generates the header count error:

The main node changes the 2-bit CNT field in the SCF for one frame only. In the subsequent frame, it sends the correct CNT field.

Because each subordinate node receives the SCF, all subordinates detect the (`A2B_INTPND0.HDCNTERR`) error.

2. When a subordinate node generates the header count error:

The subordinate node changes the 2-bit CNT field in the SRF. Generally, the subordinate node passes the received SRF as-is from a downstream subordinate node. In this case (because the subordinate node is receiving the write to the `A2B_GENERR` command in the frame), it is already generating a response with its own SRF, but with the wrong CNT field, as the command indicated.

Though the upstream subordinate nodes receive the SRF, the nodes do not check whether the CNT field is correct or not. The subordinate nodes only generate `A2B_INTPND0.HDCNTERR` on checks of the SCF. Therefore, when the subordinate node generates this error, only the main node detects it.

- *0x02 Generate Data Decoding Error* (`A2B_GENERR.GENDDERR`)

Generating a data decoding error requires a Manchester coding violation to be applied to data slots, not to the SCF and SRF fields.

1. When the main node generates the data decoding error:

The main node induces a Manchester encoding error on the first downstream data slot (slot 0 only). It does not inject the error on any other data slots. Since nodes report the data decode error only on data slots that are consumed, only the subordinate nodes that consume slot 0 detect the error when the main node generates it. When a subordinate node passes (without consuming) the data downstream, it sends the same bit stream that it receives and does not detect the error.

2. When a subordinate node generates the data decoding error:

The subordinate node induces a Manchester encoding error on the first upstream data slot it contributes, not on any passed data slot(s). If the subordinate contributes more than one upstream slot, it only induces the error on the first one. subordinate nodes do not induce encoding errors on downstream data.

Since data decode errors are only reported on data slots which are consumed, only the upstream nodes that consume the first contributed upslot detect the error. If an upstream subordinate node or a main node does not consume the first contributed data slot, then it does not detect the error.

- **0x04 Generate CRC Error** (`A2B_GENERR.GENCRRCERR`)

1. When the main node generates the CRC error:

The main node induces the error in the CRC field of the SCF for one frame only. Because each subordinate node receives the SCF, all subordinates detect the error in the CRC.

2. When a subordinate node generates the CRC error:

Subordinate nodes induce the error in the CRC field of the SRF for one frame only. Since all upstream subordinate nodes receive the SRF and check the CRC, all of them detect this error when any downstream subordinate generates it. The subordinate nodes report the SRF CRC errors in the `A2B_INTPND0.SRFRCRCERR` field, but these errors are not counted by the bit error counter. The main node detects the error as `A2B_INTPND0.CRCERR` and increments the bit error counter, if enabled.

- **0x08 Generate Data Parity Error**, `A2B_GENERR.GENDPERR`

1. When the main node generates the data decoding error:

The main node induces the data parity error on the first downstream data slot (slot 0). It does not induce the error on other data slots. When the main node generates the data parity error, only the subordinate nodes that consume slot 0 detect it. Subordinate nodes that do not consume slot 0 do not detect it.

2. When a subordinate node generates the data decoding error:

A subordinate node induces the data parity error on only the first upstream data slot it contributes. It does not induce the error in the downstream portion of the superframe. When a subordinate node generates the error, all of the upstream nodes that consume the first contributed slot detect it. If an upstream subordinate node or a main node does not consume the first contributed data slot, it does not detect it.

- **0x10 Generate Interrupt Frame CRC Error** (`A2B_GENERR.GENICRCERR`)

1. The main node cannot generate the interrupt frame CRC error.
2. When a subordinate node generates the interrupt frame CRC error, only the main node is able to detect that error. Other upstream subordinate nodes do not check the CRC in the interrupt frame.

Unique ID

Each transceiver contains a 48-bit unique ID. Read the `A2B_CHIPID0` through `A2B_CHIPID5` registers to obtain the unique ID. If a read of the unique ID fails, an interrupt is generated (`A2B_INTTYPE = 0xFC`), which indicates that the unique ID cannot be recovered. If this occurs, return the transceiver to Analog Devices.

5 A²B System Debug

The following sections provide information on system diagnostics for fault isolation and correction. In addition to the A²B line fault detection, a loop back test mode is provided to validate the I²S/TDM connections in main and subordinate nodes.

Line Fault Diagnostics

This section discusses the A²B line fault diagnostics. It provides descriptions of the different faults and programming instructions for responding to line fault events in software.

Diagnostics During Discovery

The *Line Faults* table shows the different types of line faults and the pins affected by the faults. All faults can be detected and localized during discovery of the bus. When a fault is detected during discovery, the switches that enable bias current to the next in line node are disconnected automatically.

Open wires are indicated by the `A2B_INTTYPE` register value of 0x0C.

Wires accidentally connected to the wrong port of the next node (port B instead of port A) also can create the same response or flag 0x0D to the `A2B_INTTYPE` register.

Reverse wire faults occur when the positive wire of one node accidentally connects to the negative input of the next in line node. This event is flagged with the `A2B_INTTYPE` register value 0x0D or indicated by a timeout while waiting for the discovery done response (`A2B_INTTYPE` = 0x18).

Timeout during discovery also occurs when an invalid value is programmed to the `A2B_DISCVRY.DRESPCYC` bit field, or if the next in line node has a physical defect that prevents the node from responding.

The faults that require specific software flow for detection and localization are shown with shading in the table.

NOTE: The `A2B_SWCTL.ENS` bit is not cleared automatically when a line fault opens the bias switches; this has to be done in software. The `A2B_SWCTL.ENS` bit in the main/LPS transceiver should be set to 0 in the event of a critical line fault to disconnect bus bias to affected bus segments.

Table 5-1: Line Faults

| Wires | Affected Pins | Detect | Localize | INTTYPE | Remarks |
|---|-----------------------|--------|----------|--|--|
| <i>Partial Bus Operation May Continue for Nodes Upstream of the Fault</i> | | | | | |
| Open | BP | Yes | Yes | 0x0C or 0x0D ^{*1} | Open wires (BP and BN are the B-side positive and negative connector pins) |
| | BN | | | | |
| | BN and BP | | | | |
| Wrong Port | B to B' port | Yes | Yes | 0x0C or 0x0D | B' is B-side of next in line node |
| Reverse Wires | BN to AP and BP to AN | Yes | Yes | 0x0D | Wrong port or reverse wires (AP is the A side positive connector pin of the next in line node) |
| | | | | No 0x18 timeout (no DSCDONE interrupt) | Reverse wires undetected by hardware diagnostics |
| Defective Node | NA | Yes | Yes | No 0x18 timeout (no DSCDONE interrupt) | Defective node or wrong software parameter value for A2B_DISCVRY.DRESPCYC |
| Short of Wires | BP with BN | Yes | Yes | 0x0B | Wires shorted together |
| <i>Critical Faults</i> | | | | | |
| Short to Ground | BP | Yes | Yes | 0x09 | Positive wire shorted to ground |
| | BN | | Yes | 0x29 | Software routine localizes fault |
| Short to V _{BAT} | BN | Yes | Yes | 0x0A | Negative wire shorted to V _{BAT} |
| | BP | | Yes | 0x2A | Software routine localizes fault |

*1 At lower VIN voltages, open wire faults may provide the 0x0D fault code for AD242x devices.

CAUTION: The short to ground and short to V_{BAT} faults are critical faults for which the whole bus shuts off up to the next local powered node. Normal A²B bus operation should always be discontinued, including removal of bus power by the main/LPS node (independent of line fault location). Program the A2B_SWCTL.ENSX bit = 0 to the main/LPS transceiver until the fault is corrected.

For the following faults, partial A²B bus operation can continue between the main and subordinate nodes which are upstream of the line fault location.

- Open
- Wrong port
- Reverse wires
- Defective node
- Wrong discovery parameter for next-in-line node

- Wires shorted together

NOTE: The A-port of a LPS (Local Power Subordinate Node) is isolated from its B-port. This prevents the fault that occurred at the B-port or the fault occurred at a downstream bus-powered node to back-propagate to previous nodes. In case of a critical fault, the bus can partially operate until the LPS node, which reports the fault or is the last LPS before the fault. However, if the downstream nodes are contributing upstream slots to the main node or any other upstream node present between the main and the LPS, then parity errors occur in cases of partial bus operation.

Registers for Line Diagnostics

The following registers are used to diagnose line faults on the A²B bus. Refer to the *Register Descriptions* section for details.

- The [A2B_SWCTL](#) register controls the bias voltage to be switched onto the B-side A²B bus link for the next in line nodes. The register also provides special line fault sensing modes.
- The [A2B_SWSTAT](#) register provides line diagnostics status information.
- The [A2B_INTSRC](#) register contains information about the source of an active interrupt, which subordinate node generated it, or whether the interrupt originates from the main node. Line errors can be located with this register.
- The [A2B_INTTYPE](#) register stores information about the type of the current interrupt request. A read of this register clears the corresponding interrupt.

Open Wire Fault

The *Open Wire Fault* figure shows an open wire fault between subordinate node 0 and subordinate node 1 transceivers. Communication continues between the main node and subordinate node 0 transceivers when this fault occurs between subordinate node 0 and subordinate node 1.

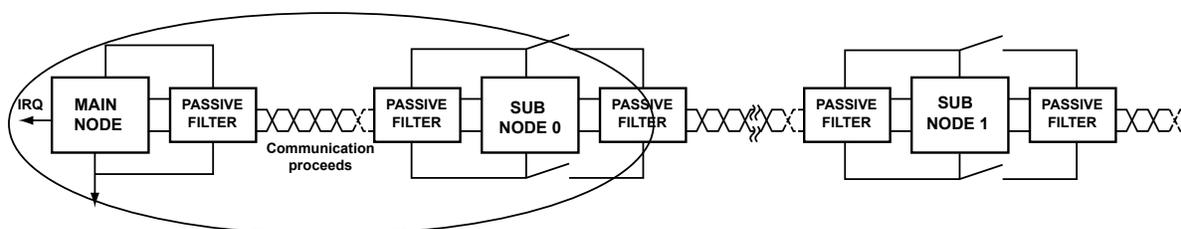


Figure 5-1: Open Wire Fault

Short of Wires Fault

The *Short of Wires Fault* figure shows a short of wires line fault between subordinate node 0 and subordinate node 1. Communication continues between the main node and subordinate node 0 transceivers when this fault occurs between subordinate node 0 and subordinate node 1.

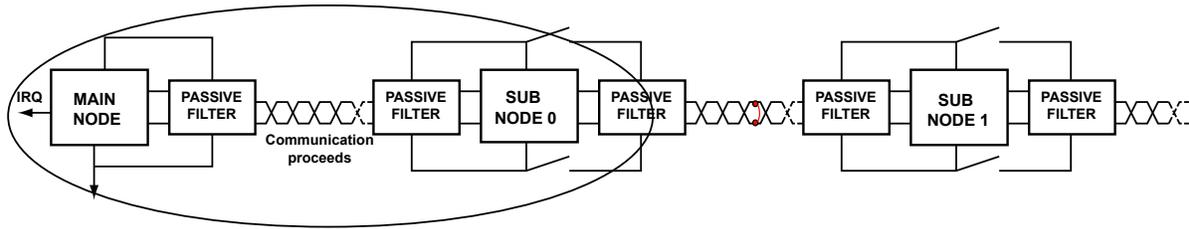


Figure 5-2: Short of Wires Fault

Short To GND BP

The *Short to GND BP Fault* figure shows the BP wire shorted to ground between subordinate node 0 and subordinate node 1 transceivers. All bus communication stops when this fault occurs between subordinate node 0 and subordinate node 1 transceivers.

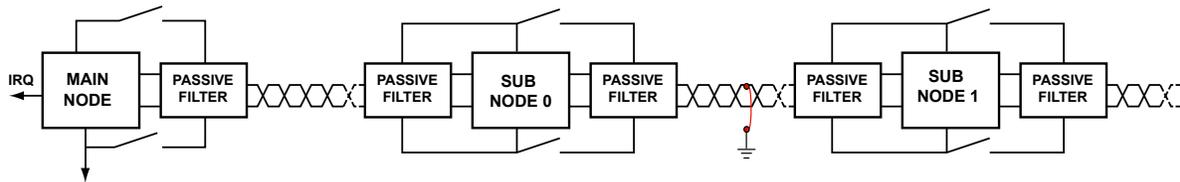
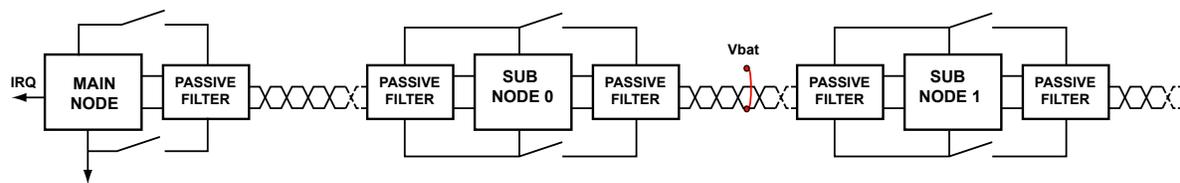


Figure 5-3: Short To GND BP

Short to V_{BAT} BN

The *Short to V_{BAT} BN* figure shows the BN wire shorted to V_{BAT} between sub node 0 and sub node 1 transceivers. All bus communication stops when this fault occurs.


 Figure 5-4: Short to V_{BAT} BN

Short To GND BN

The *Short to GND BN* figure shows the BN wire shorted to ground between sub node 0 and sub node 1 transceivers. Bus communication can continue without an immediate fault when this fault occurs between sub node 0 and sub node 1.

NOTE: This line fault is a special diagnostic case because it propagates to earlier nodes as the FET switches have reverse diodes. During discovery or rediscovery of the bus, this fault is identified as not localized with specific `A2B_INTTYPE` code (0x29). In order to localize the fault, set the `A2B_SWCTL.DIAGMODE` bit = 1.

See the [Diagnostics Software Flow](#) section and the [Localizing Concealed Faults](#) table for the fault diagnostics software flow.

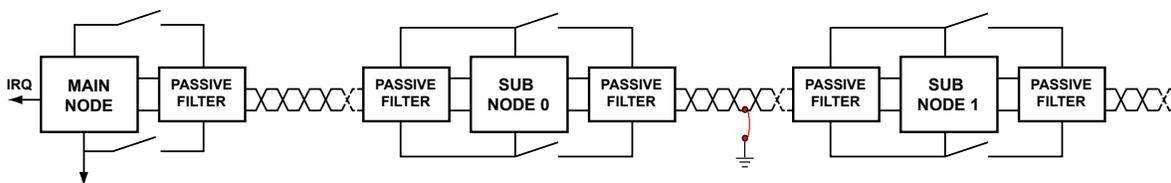


Figure 5-5: Short to GND BN

Short to V_{BAT} BP

The *Short to V_{BAT} BP* figure shows the BP wire shorted to V_{BAT} between sub node 0 and sub node 1. Bus communication can continue without an immediate fault when the short to V_{BAT} BP fault occurs between sub node 0 and sub node 1.

NOTE: This line fault is a special diagnostic case because it propagates to earlier nodes as the FET switches have reverse diodes. During discovery or rediscovery of the bus, this fault is identified as not localized with specific `A2B_INTTYPE` code (0x2A). In order to localize the fault, set the `A2B_SWCTL.DIAGMODE` bit = 1. See the [Diagnostics Software Flow](#) section and the [Localizing Concealed Faults](#) figure for more information about the fault diagnostics software flow.

NOTE: In some cases, a BP short to VBAT fault occurs at a downstream bus-powered-node. This fault can back-propagate to the upstream nodes and cause an indeterminate fault (for example `A2B_INTTYPE` = 0x0F). If this fault occurs during discovery, follow the procedure described in [Diagnostics Software Flow](#)

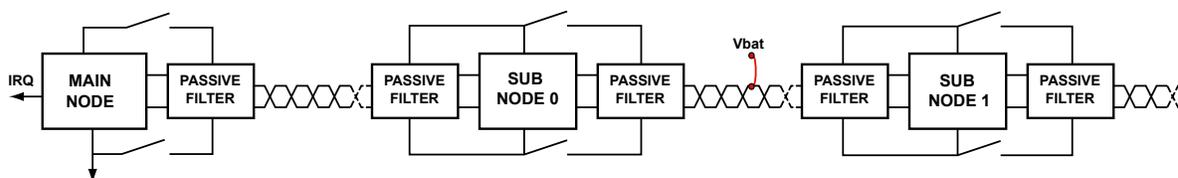


Figure 5-6: Short to V_{BAT} BP

Line Diagnostics After Discovery

Full line diagnostics are only performed during discovery. However, certain interrupts (if enabled) after discovery can indicate line faults during operation. Rediscovery detects the cause and location of the faults which are can be localized.

After discovery, any of the following interrupt types (`A2B_INTTYPE`) indicate that there is a line fault:

- 0x0A (10: PWRERR)
- 0x0F (15: PWRERR)

- 0x2A (42: PWRERR)
- 0x80 (128: interrupt messaging error)
- 0xFD (253: subordinate node INTTYPE read error)

When a subordinate node detects the SRF missed error (SRFMISSEERR) in 32 consecutive frames, the node assumes a downstream bus drop and sets its last node bit (`A2B_NODE.LAST = 1`) to become the last node in the system. A bus drop condition resulting from a line fault occurring after discovery can be detected in a last node (`A2B_NODE.LAST = 1`) that has the SRFMISSEERR latched.

Excessive accumulation of bit errors can happen if there is a slot configuration mismatch between nodes. This can also happen when positive A²B wire shorts to a noisy V_{BAT} or negative A²B wire shorts to a noisy GND. The bus can operate under these conditions but is more susceptible to impairments (for example, electromagnetic interference).

Use the `A2B_BECNT` register to count accumulated errors as follows.

- Set the `A2B_BECCTL` register to 0xE4 (interrupt after 256 CRC errors). Acceptable audio noise and robustness is subjective and needs to be determined in vehicle tests. Adjust the threshold accordingly.
- Periodically write 0 to the `A2B_BECNT` register (once every second) to reset the error counter. Acceptable audio noise and robustness is subjective and needs to be determined in vehicle tests. Adjust time for the `A2B_BECNT` register accordingly.
- The bit error counter overflow (0x04: BECOVF) interrupt indicates bus issues.

Diagnostics Software Flow

Use the following software flow and the *Diagnostics Software Flow* figure for node discovery with diagnostics.

1. Set the `A2B_SWCTL` register = 0x00.
2. Enable the power error interrupts and the `A2B_INTPND2.DSCDONE` interrupt in the main node. Set the `A2B_SWCTL` register = 0x01 to enable the power switch.
3. Wait for interrupt to occur. If the `A2B_INTTYPE` register = 0x18 for `A2B_INTPND2.DSCDONE` (indicating a successful node discovery), proceed to step 7.
4. If the `A2B_INTTYPE` register = 0x29 or 0x2A or 0x0F, configure the `A2B_SWCTL.ENS` bit = 0 in the next upstream local powered node and wait 50-100ms. Proceed to rediscovery with the `A2B_SWCTL.DIAGMODE` bit = 1 in the [Localizing Concealed Faults](#) section (step 8).

ADDITIONAL INFORMATION: If the `A2B_INTTYPE` register is any other `A2B_INTPND0.PWRERR` type or if the discovery operation times out, proceed to step 5.

5. If the `A2B_INTTYPE` register = 0x0B, 0x0C, or 0x0D, the `A2B_INTSRC` register can be read to determine the location. If the operation times out, then by process of elimination the bus wires to the node being discovered are most likely reversed. Proceed to step 6.

- If the `A2B_INTTYPE` register = 0x09 or 0x0A, disable the entire bus until the next upstream local powered node by clearing the `A2B_SWCTL.ENS` bit in the next upstream local powered node (after the `A2B_INTSRC` and `A2B_INTTYPE` register values have been communicated to the host).

ADDITIONAL INFORMATION: Once any other localized fault has been detected, halt the discovery process. Retry the discovery process periodically by software to determine if the fault is cleared. There is no automatic retry mechanism within the transceiver.

- If this is not the last node, reprogram the `A2B_SWCTL.MODE` bits = 2. This setting ignores fluctuation on VIN due to downstream current draw and prevents incorrect localization on errors that occur on nodes that are located further downstream. Program the downstream node register settings and repeat step 1 on next node.

ADDITIONAL INFORMATION: Continue this cycle until all nodes are discovered. Once all nodes are discovered, configure the `A2B_SWCTL.MODE` bits = 0 to all nodes while keeping the `A2B_SWCTL.ENS` bit = 1. Full A²B bus discovery is complete now.

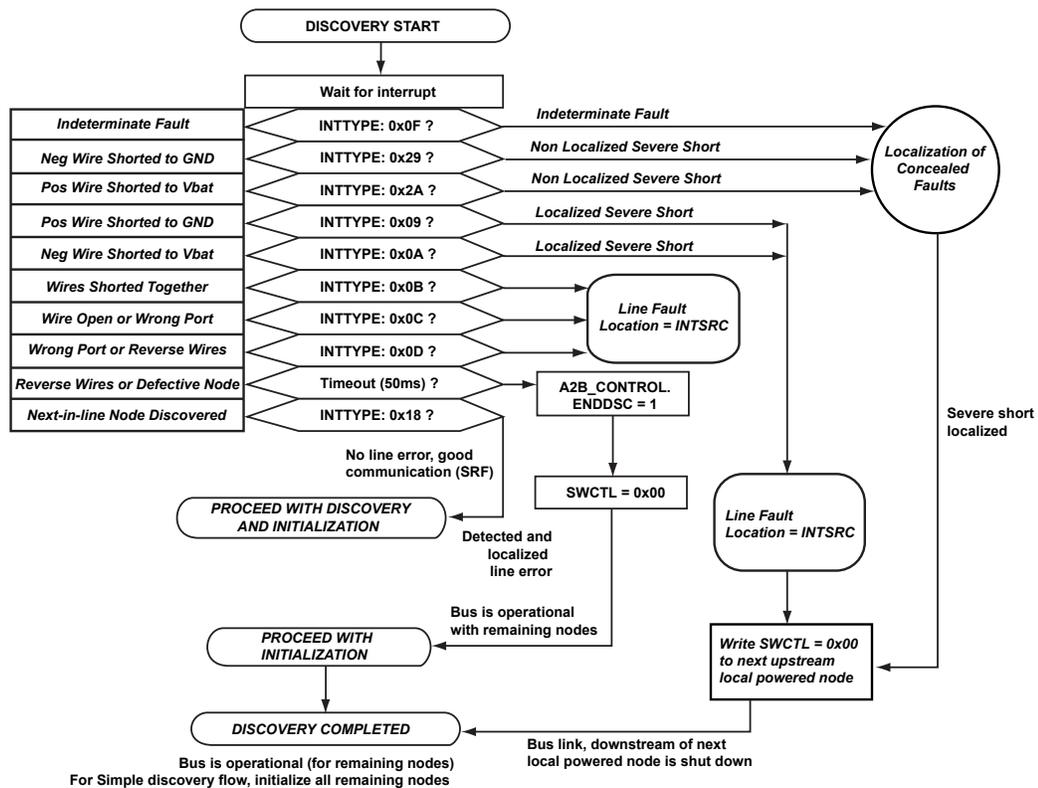


Figure 5-7: Diagnostics Software Flow

- If there is a timeout while discovering a subordinate node, stop the discovery process by setting the `A2B_CONTROL.ENDDSC` bit.

Localizing Concealed Faults

This section describes the procedure to determine the location of the potentially concealed faults to either V_{BAT} or GND.

While the bus should not be operated long term in the presence of such a fault, run the following procedure on a short term basis to establish the location of a concealed fault before disabling the bus at the main node. This process is shown in the *Localization of Concealed Faults* figure.

1. Following from step 4 in [Diagnostics Software Flow](#), where the `A2B_SWCTL.ENS` bit = 0 in the main/LPS transeiver, set variables to keep track of the current node position and the last known good node. Also set a variable, for example `PriorFault = 0`. This keeps track of whether or not a fault is reported in a prior node discovery. After waiting at least 100 ms to allow for the electrical steady state of the bus to settle, proceed to step 2.

ADDITIONAL INFORMATION: Start fault localization from the LPS node if partial bus operation is needed. Otherwise start from the main node.

2. In the current node, set the `A2B_INTMSK0` register = 0x10 and the `A2B_SWCTL` register = 0x09. This sets the `A2B_SWCTL.ENS` bits initiating rediscovery in diagnostic mode. Define two variables for keeping track of whether or not a fault and/or discovery is completed in the current discovery attempt; for example, `Disc` and `Fault`. Clear both variables in this step.

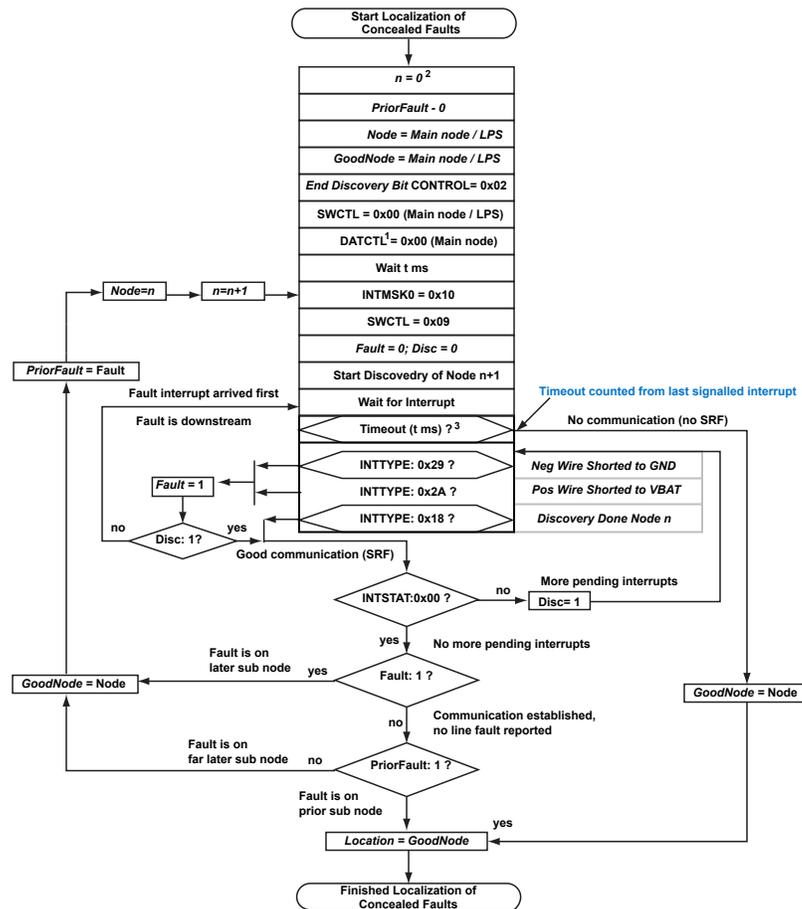
ADDITIONAL INFORMATION: Wait for an interrupt from this operation, allow 100 ms for a timeout. This timeout provides sufficient time for bus diagnostics and, possibly, full discovery to complete. This process can take longer than usual when `A2B_SWCTL.ENS = 1` in the presence of a fault.

3. If no interrupt is received before the timeout expires, the fault is located immediately downstream of the current node. Set `GoodNode` to the current node. Proceed to step 9.
4. If the `A2B_INTTYPE` register = 0x29 or 0x2A or 0x0F, then an error occurred somewhere downstream of the current node. This means a fault is detected so set `Fault = 1`. If `DISC = 0`, return to step 3 with a 100 ms timeout waiting to see if discovery completes. If `DISC = 1` and the discovery process was reported previously as complete (`A2B_INTTYPE` register = 0x18), proceed to step 6.
5. If the `A2B_INTTYPE` register = 0x18, the downstream node was successfully discovered and communication has been established. In diagnostic mode, this can occur even in the presence of a detected fault of `A2B_INTTYPE = 0x29` or `0x2A`. These faults can occur when the physical fault exists on only one of the wires between the two nodes. Proceed to step 6.
6. Check the value of the `A2B_INTSTAT` register for other pending interrupts. If the `A2B_INTSTAT` register is non-zero, the fault and discovery completion both occurred faster than the interrupt service routine response. In this case, the 0x18 DSCDONE interrupt is a higher priority. Set `DISC = 1` and return to step 3. If the `A2B_INTSTAT` register = 0, there are no more pending interrupts. Proceed to step 7.
7. In order to reach this step, discovery must have been completed successfully. If a fault was also detected, then `Fault = 1`, and it is necessary to continue the bus discovery to localize the fault. Set `GoodNode = Node`,

PriorFault = Fault, and Node = n. Increment n in preparation for the discovery of the next node and return to step 2. If Fault = 0 (Fault was not detected), proceed to step 8.

8. To reach this step, discovery must have been completed and no fault was detected. This can occur for one of two reasons. Either the current node is too far upstream of the fault to detect it yet, or the node is already downstream of the fault where fault is no longer present. If PriorFault = 1, then it is the latter case, so proceed to step 9. If PriorFault = 0, then the fault has yet to be detected. In this case, continue bus discovery to localize the fault. Set GoodNode = Node, PriorFault = Fault, and Node = n. Increment n in preparation for the discovery of the next node and return to step 2.
9. Report the fault location as being immediately downstream of the last recorded GoodNode. The location of the error is after the current node unless in step 8 it was detected that the line fault is before the current node. In this case, the last GoodNode is one node up. Localization of the concealed fault is complete.

Figure 5-8: Localization of Concealed Faults



LEGEND

1. Step only needed when localizing fault from the main node.
2. Node numbering starts from the main node/LPS. For example, if GoodNode = 0, then the fault is downstream to the main node/LPS.
3. t = 100ms

Bus Drop Detection

The *Diagnostics Software Flow* figure describes the flow of bus drop detection in the A²B system.

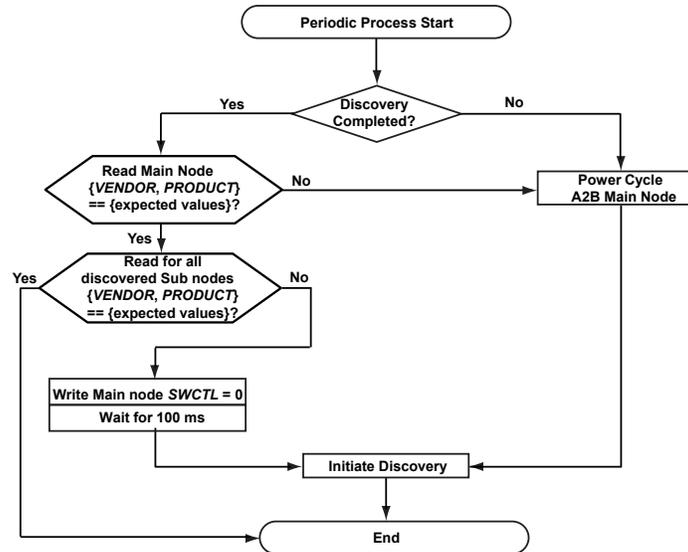


Figure 5-9: Diagnostics Software Flow

I²S Loopback

I²S loopback occurs inside the transceiver. Data driven to the DTX0 pad is sampled as A²B receive data instead of the data on the DRX0 pin. Data driven to the DTX1 pad is sampled as A²B receive data instead of the data on the DRX1 pin.

The `A2B_I2STEST.BUSLOOPBK` bit enables loopback from the DTX0 pins to the serial RX blocks. The values of the `A2B_I2STEST.SELRX1`, `A2B_I2STEST.RX2LOOPBK`, and `A2B_I2STEST.LOOPBK2TX` bits are ignored if this bit is set. If the `A2B_I2STEST.PATTRN2TX` bit is set, a fixed pattern (0xB38F0E32) is driven on the DTX0 and DTX1 pins instead of transmit data from the A²B bus.

If I²S Loopback mode is enabled, program the value of the `A2B_I2SCFG.RX0EN` bit to match the value of the `A2B_I2SCFG.TX0EN` bit, and the value of the `A2B_I2SCFG.RX1EN` bit to match the value of the `A2B_I2SCFG.TX1EN` bit.

The number of data slots received and transmitted on the A²B bus by each node is controlled by a number of registers.

If the `A2B_SLOTFMT.UPSIZE` and `A2B_SLOTFMT.DNSIZE` bit field values are different, looped back data, which changes direction on the bus, is either truncated to a smaller bit width or zero-filled to a larger bit width.

When this mode is enabled, the program is responsible for ensuring that the data received from the A²B bus and looped back through the serial blocks can be transmitted on the A²B bus.

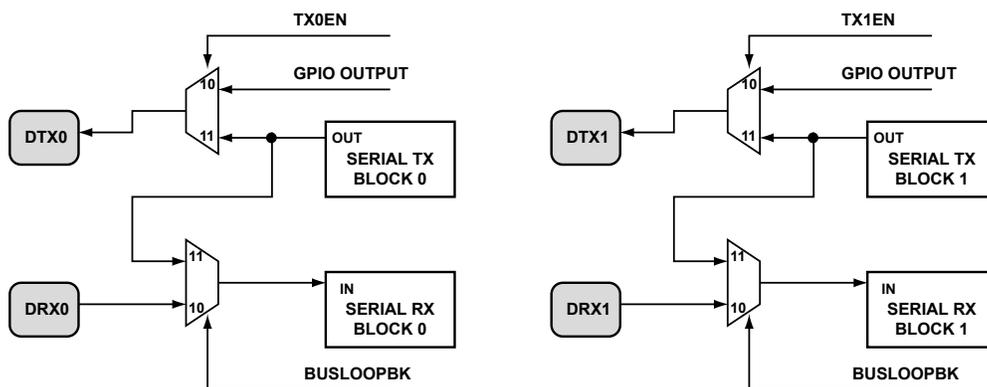


Figure 5-10: Serial TX Block to Serial RX Block

I²S TDM Test Mode (I²S Loopback)

Pattern generation and loopback test modes are provided for easy validation of I²S TDM connectivity in main and subordinate nodes. The transmit pattern generator uses the default bit pattern 1011_0011_1000_1111__0000_1110_0011_0010 on all channels, where 1011 is the most significant nibble and 0010 is the least significant nibble.

Use the following procedure for I²S TDM testing.

1. For main node to host link verification, set the `A2B_I2STEST.PATTRN2TX` bit in the main transceiver and verify that the TX interface with the default bit pattern matches the expected timing (possibly using a scope, logic analyzer, or other device).
2. For host to main node link verification. Set the `A2B_I2STEST.RX2LOOPBK` and `A2B_I2STEST.LOOPBK2TX` bits in the main transceiver, wait one cycle, and verify that the DTX data received at the host matches the DRX data sent from the previous frame.

ADDITIONAL INFORMATION: The RX to TX loopback does not working correctly when the main node is also receiving TX data from the bus. The `A2B_DATCTL` register must be 0x00 while looping back from RX to TX.

3. For sub node to peripheral link verification, if a sub node is connected to a DAC (for example, to send to a speaker), set the `A2B_I2STEST.PATTRN2TX` bit in the sub node and verify the expected DTX timing.
4. For peripheral to sub node link verification, if a sub node has a peripheral that provides input signals over the I²S/ TDM interface, set the `A2B_I2STEST.RX2LOOPBK` and `A2B_I2STEST.LOOPBK2TX` bits. Verify that the DTX interface matches the DRX interface with a one frame delay. Alternatively (without using the `A2B_I2STEST` register) check the RX data at the previously verified main node I²S/TDM DTX interface.
5. System verification with external loopback. Connect the DTX0/DTX1 pins with the DRX0/DRX1 pins in a sub node to generate a digital loopback. The default bit pattern can be verified at the main node DTX pins when the `A2B_I2STEST.PATTRN2TX` bit is set at the sub node.

ADDITIONAL INFORMATION: If the `A2B_I2STEST.RX2LOOPBK` bit is cleared while the `A2B_I2STEST.LOOPBK2TX` bit is set, then the last received frame is repeated on the TX pins. This behavior persists until the `A2B_I2STEST.RX2LOOPBK` bit is set or the `A2B_I2STEST.LOOPBK2TX` bit is cleared. If the `A2B_I2STEST.LOOPBK2TX` bit is enabled after reset, the default pattern is generated until the `A2B_I2STEST.RX2LOOPBK` bit is set.

ADDITIONAL INFORMATION: The *Frame Buffer* figure shows the TX frame buffer that is used for loop-back tests.

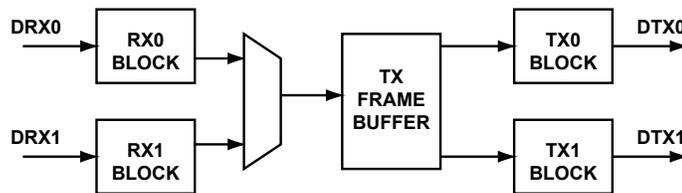


Figure 5-11: Frame Buffer

6 Register Summary

The following table provides the map of the AD2426/AD2427/AD2428 registers and bits.

| Reg. Addr. | Name | Bit 7 | Bit 6 | Bit 5 | Bit 4 | Bit 3 | Bit 2 | Bit 1 | Bit 0 | Reset | RW | | |
|------------|------------|------------|------------|----------|-----------|----------|----------|----------|----------|-------|------|------|-----|
| 0x00 | CHIP | Reserved | | CHIPADR | | | | | | | | 0x50 | R/W |
| 0x01 | NODEADR | BRCST | Reserved | PERI | Reserved | NODE | | | | 0x00 | R/W | | |
| 0x02 | VENDOR | VENDOR | | | | | | | | | 0xAD | R/NW | |
| 0x03 | PRODUCT | PRODUCT | | | | | | | | | 0x25 | R/NW | |
| 0x04 | VERSION | VERSION | | | | | | | | | 0x00 | R/NW | |
| 0x05 | CAPABILITY | Reserved | | | | | | | I2CAVAIL | 0x01 | R/NW | | |
| 0x09 | SWCTL | Reserved | DISNXT | MODE | | DIAGMODE | Reserved | | ENSW | 0x00 | R/W | | |
| 0x0A | BCDNSLOTS | Reserved | | | BCDNSLOTS | | | | | | 0x00 | R/W | |
| 0x0B | LDNSLOTS | DNMASKEN | Reserved | LDNSLOTS | | | | | | 0x00 | R/W | | |
| 0x0C | LUPSLOTS | Reserved | | | LUPSLOTS | | | | | | 0x00 | R/W | |
| 0x0D | DNSLOTS | Reserved | | | DNSLOTS | | | | | | 0x00 | R/W | |
| 0x0E | UPSLOTS | Reserved | | | UPSLOTS | | | | | | 0x00 | R/W | |
| 0x0F | RESPCYCS | RESPCYCS | | | | | | | | | 0x40 | R/W | |
| 0x10 | SLOTFMT | UPFMT | UPSIZE | | | DNFMT | DNSIZE | | | | 0x00 | R/W | |
| 0x11 | DATCTL | STANDBY | Reserved | ENDSNIFF | Reserved | | | UPS | DNS | 0x00 | R/W | | |
| 0x12 | CONTROL | MSTR | Reserved | | XCVRBINV | SWBYP | SOFTTRST | ENDDSC | NEWSTRCT | 0x00 | R/W | | |
| 0x13 | DISCVRY | DRESPCYC | | | | | | | | | 0x00 | R/W | |
| 0x14 | SWSTAT | FAULT_NLOC | FAULT_CODE | | | Reserved | | FAULT | FIN | 0x00 | R/NW | | |
| 0x15 | INTSTAT | Reserved | | | | | | | IRQ | 0x00 | R/NW | | |
| 0x16 | INTSRC | MSTINT | SLVINT | Reserved | | INODE | | | | 0x00 | R/NW | | |
| 0x17 | INTTYPE | TYPE | | | | | | | | | 0x00 | R/NW | |
| 0x18 | INTPND0 | SRFCRCERR | SRFERR | BECOVF | PWRERR | DPERR | CRCERR | DDERR | HDCNTERR | 0x00 | R/W | | |
| 0x19 | INTPND1 | IO7PND | IO6PND | IO5PND | IO4PND | IO3PND | IO2PND | IO1PND | IO0PND | 0x00 | R/W | | |
| 0x1A | INTPND2 | Reserved | | | | SLVIRQ | ICRCERR | I2CERR | DSCDONE | 0x00 | R/W | | |
| 0x1B | INTMSK0 | SRFCRCIEEN | SRFEIEN | BECIEN | PWREIEN | DPEIEN | CRCEIEN | DDEIEN | HCEIEN | 0x00 | R/W | | |
| 0x1C | INTMSK1 | IO7IRQEN | IO6IRQEN | IO5IRQEN | IO4IRQEN | IO3IRQEN | IO2IRQEN | IO1IRQEN | IO0IRQEN | 0x00 | R/W | | |
| 0x1D | INTMSK2 | Reserved | | | | SLVIRQEN | ICRCIEEN | I2CEIEN | DSCDIEN | 0x00 | R/W | | |
| 0x1E | BECCTL | THRESHLD | | | ENICRC | ENDP | ENCRC | ENDD | ENHDCNT | 0x00 | R/W | | |
| 0x1F | BECNT | BECNT | | | | | | | | | 0x00 | R/W | |

Register Summary

| Reg. Addr. | Name | Bit 7 | Bit 6 | Bit 5 | Bit 4 | Bit 3 | Bit 2 | Bit 1 | Bit 0 | Reset | RW | |
|------------|-------------|-----------------|------------|------------|------------|------------|------------|------------|------------|-------|------|------|
| 0x20 | TESTMODE | Reserved | | RXDPH | | Reserved | PRBSN2N | PRBSDN | PRBSUP | 0x00 | R/W | |
| 0x21 | ERRCNT0 | ERRCNT0 [7:0] | | | | | | | | | 0x00 | R/NW |
| 0x22 | ERRCNT1 | ERRCNT1 [15:8] | | | | | | | | | 0x00 | R/NW |
| 0x23 | ERRCNT2 | ERRCNT2 [23:16] | | | | | | | | | 0x00 | R/NW |
| 0x24 | ERRCNT3 | ERRCNT3 [31:24] | | | | | | | | | 0x00 | R/NW |
| 0x29 | NODE | LAST | NLAST | DISCVD | Reserved | NUMBER | | | | 0x80 | R/NW | |
| 0x2B | DISCSTAT | DSCACT | Reserved | | | DNODE | | | | 0x00 | R/NW | |
| 0x2E | TXACTL | TXAOVREN | Reserved | | | | TXALEVEL | | | 0x00 | R/W | |
| 0x30 | TXBCTL | TXBOVREN | Reserved | | | | TXBLEVEL | | | 0x00 | R/W | |
| 0x3E | LINTTYPE | LINTTYPE | | | | | | | | | 0x00 | R/NW |
| 0x3F | I2CCFG | Reserved | | | | | FRAMERATE | EACK | DATARATE | 0x00 | R/W | |
| 0x40 | PLLCTL | SSMODE | | Reserved | | | SSDEPTH | Reserved | SSFREQ | 0x00 | R/W | |
| 0x41 | I2SGCFG | INV | EARLY | ALT | TDMSS | RXONDTX1 | TDMMODE | | | 0x00 | R/W | |
| 0x42 | I2SCFG | RXBCLKINV | RX2PINTL | RX1EN | RX0EN | TXBCLKINV | TX2PINTL | TX1EN | TX0EN | 0x00 | R/W | |
| 0x43 | I2SRATE | SHARE | REDUCE | BCLKRATE | | | I2SRATE | | | 0x00 | R/W | |
| 0x44 | I2STXOFFSET | TSBEFORE | TSAFTER | TXOFFSET | | | | | | 0x00 | R/W | |
| 0x45 | I2SRXOFFSET | Reserved | | RXOFFSET | | | | | | 0x00 | R/W | |
| 0x46 | SYNCOFFSET | SYNCOFFSET | | | | | | | | | 0x00 | R/W |
| 0x47 | PDMCTL | Reserved | PDMRATE | | HPFEN | PDM1SLOTS | PDM1EN | PDM0SLOTS | PDM0EN | 0x00 | R/W | |
| 0x48 | ERRMGMT | Reserved | | | | | ERRSLOT | ERRSIG | ERRLSB | 0x00 | R/W | |
| 0x4A | GPIODAT | IO7DAT | IO6DAT | IO5DAT | IO4DAT | IO3DAT | IO2DAT | IO1DAT | IO0DAT | 0x00 | R/W | |
| 0x4B | GPIODATSET | IO7DSET | IO6DSET | IO5DSET | IO4DSET | IO3DSET | IO2DSET | IO1DSET | IO0DSET | 0x00 | R/W | |
| 0x4C | GPIODATCLR | IO7DCLR | IO6DCLR | IO5DCLR | IO4DCLR | IO3DCLR | IO2DCLR | IO1DCLR | IO0DCLR | 0x00 | R/W | |
| 0x4D | GPIOOEN | IO7OEN | IO6OEN | IO5OEN | IO4OEN | IO3OEN | IO2OEN | IO1OEN | IO0OEN | 0x00 | R/W | |
| 0x4E | GPIOEN | IO7IEN | IO6IEN | IO5IEN | IO4IEN | IO3IEN | IO2IEN | IO1IEN | IO0IEN | 0x00 | R/W | |
| 0x4F | GPIOIN | IO7IN | IO6IN | IO5IN | IO4IN | IO3IN | IO2IN | IO1IN | IO0IN | 0x00 | R/NW | |
| 0x50 | PINTEN | IO7IE | IO6IE | IO5IE | IO4IE | IO3IE | IO2IE | IO1IE | IO0IE | 0x00 | R/W | |
| 0x51 | PINTINV | IO7INV | IO6INV | IO5INV | IO4INV | IO3INV | IO2INV | IO1INV | IO0INV | 0x00 | R/W | |
| 0x52 | PINCFG | Reserved | | IRQTS | IRQINV | Reserved | | | DRVSTR | 0x01 | R/W | |
| 0x53 | I2STEST | Reserved | | | BUSLOOPBK | SELRX1 | RX2LOOPBK | LOOPBK2TX | PATTRN2TX | 0x00 | R/W | |
| 0x54 | RAISE | RTYPE | | | | | | | | | 0x00 | R/W |
| 0x55 | GENERR | Reserved | | | GENICRCERR | GENDPERR | GENCRCERR | GENDDERR | GENHCERR | 0x00 | R/W | |
| 0x56 | I2SRRATE | RBUS | Reserved | RRDIV | | | | | | | 0x00 | R/W |
| 0x57 | I2SRRCTL | Reserved | | STRBDIR | ENSTRB | Reserved | | ENXBIT | ENVLSB | 0x00 | R/W | |
| 0x58 | I2SRRSOFFS | Reserved | | | | | | RRSOFFSET | | 0x00 | R/W | |
| 0x59 | CLK1CFG | CLK1EN | CLK1INV | CLK1PDIV | Reserved | CLK1DIV | | | | 0x00 | R/W | |
| 0x5A | CLK2CFG | CLK2EN | CLK2INV | CLK2PDIV | Reserved | CLK2DIV | | | | 0x00 | R/W | |
| 0x5B | BMMCFG | Reserved | | | | | BMMNDSC | BMMRXEN | BMMEN | 0x00 | R/W | |
| 0x5C | SUSCFG | Reserved | | SUSDIS | SUSOE | Reserved | SUSSEL | | | 0x00 | R/W | |
| 0x5D | PDMCTL2 | Reserved | | PDMINVCLK | PDMALTCLK | PDM1FFRST | PDM0FFRST | PDMDEST | | 0x00 | R/W | |
| 0x60 | UPMASK0 | RXUPSL0T07 | RXUPSL0T06 | RXUPSL0T05 | RXUPSL0T04 | RXUPSL0T03 | RXUPSL0T02 | RXUPSL0T01 | RXUPSL0T00 | 0x00 | R/W | |

| Reg. Addr. | Name | Bit 7 | Bit 6 | Bit 5 | Bit 4 | Bit 3 | Bit 2 | Bit 1 | Bit 0 | Reset | RW |
|------------|-----------|---------------|-------------|-------------|-------------|-------------|-------------|-------------|-------------|-------|------|
| 0x61 | UPMASK1 | RXUP SLOT15 | RXUP SLOT14 | RXUP SLOT13 | RXUP SLOT12 | RXUP SLOT11 | RXUP SLOT10 | RXUP SLOT09 | RXUP SLOT08 | 0x00 | R/W |
| 0x62 | UPMASK2 | RXUP SLOT23 | RXUP SLOT22 | RXUP SLOT21 | RXUP SLOT20 | RXUP SLOT19 | RXUP SLOT18 | RXUP SLOT17 | RXUP SLOT16 | 0x00 | R/W |
| 0x63 | UPMASK3 | RXUP SLOT31 | RXUP SLOT30 | RXUP SLOT29 | RXUP SLOT28 | RXUP SLOT27 | RXUP SLOT26 | RXUP SLOT25 | RXUP SLOT24 | 0x00 | R/W |
| 0x64 | UPOFFSET | Reserved | | | | UPOFFSET | | | | 0x00 | R/W |
| 0x65 | DNMASK0 | RXDNSLOT07 | RXDNSLOT06 | RXDNSLOT05 | RXDNSLOT04 | RXDNSLOT03 | RXDNSLOT02 | RXDNSLOT01 | RXDNSLOT00 | 0x00 | R/W |
| 0x66 | DNMASK1 | RXDNSLOT15 | RXDNSLOT14 | RXDNSLOT13 | RXDNSLOT12 | RXDNSLOT11 | RXDNSLOT10 | RXDNSLOT09 | RXDNSLOT08 | 0x00 | R/W |
| 0x67 | DNMASK2 | RXDNSLOT23 | RXDNSLOT22 | RXDNSLOT21 | RXDNSLOT20 | RXDNSLOT19 | RXDNSLOT18 | RXDNSLOT17 | RXDNSLOT16 | 0x00 | R/W |
| 0x68 | DNMASK3 | RXDNSLOT31 | RXDNSLOT30 | RXDNSLOT29 | RXDNSLOT28 | RXDNSLOT27 | RXDNSLOT26 | RXDNSLOT25 | RXDNSLOT24 | 0x00 | R/W |
| 0x69 | DNOFFSET | Reserved | | | | DNOFFSET | | | | 0x00 | R/W |
| 0x6A | CHIPID0 | CHIPID[7:0] | | | | | | | | 0xXX | R/NW |
| 0x6B | CHIPID1 | CHIPID[15:8] | | | | | | | | 0xXX | R/NW |
| 0x6C | CHIPID2 | CHIPID[23:16] | | | | | | | | 0xXX | R/NW |
| 0x6D | CHIPID3 | CHIPID[31:24] | | | | | | | | 0xXX | R/NW |
| 0x6E | CHIPID4 | CHIPID[39:32] | | | | | | | | 0xXX | R/NW |
| 0x6F | CHIPID5 | CHIPID[47:40] | | | | | | | | 0xXX | R/NW |
| 0x80 | GPIODEN | IOD7EN | IOD6EN | IOD5EN | IOD4EN | IOD3EN | IOD2EN | IOD1EN | IOD0EN | 0x00 | R/W |
| 0x81 | GPIOD0MSK | IOD0MSK | | | | | | | | 0x00 | R/W |
| 0x82 | GPIOD1MSK | IOD1MSK | | | | | | | | 0x00 | R/W |
| 0x83 | GPIOD2MSK | IOD2MSK | | | | | | | | 0x00 | R/W |
| 0x84 | GPIOD3MSK | IOD3MSK | | | | | | | | 0x00 | R/W |
| 0x85 | GPIOD4MSK | IOD4MSK | | | | | | | | 0x00 | R/W |
| 0x86 | GPIOD5MSK | IOD5MSK | | | | | | | | 0x00 | R/W |
| 0x87 | GPIOD6MSK | IOD6MSK | | | | | | | | 0x00 | R/W |
| 0x88 | GPIOD7MSK | IOD7MSK | | | | | | | | 0x00 | R/W |
| 0x89 | GPIODDAT | IOD7DAT | IOD6DAT | IOD5DAT | IOD4DAT | IOD3DAT | IOD2DAT | IOD1DAT | IOD0DAT | 0x00 | R/W |
| 0x8A | GPIODINV | IOD7INV | IOD6INV | IOD5INV | IOD4INV | IOD3INV | IOD2INV | IOD1INV | IOD0INV | 0x00 | R/W |
| 0x90 | MBOX0CTL | Reserved | | MBOX0LEN | | MBOX0FIEN | MBOX0E1EN | MBOX0DIR | MBOX0EN | 0x00 | R/W |
| 0x91 | MBOX0STAT | Reserved | | MBOX0E1IRQ | MBOX0F1IRQ | Reserved | | MBOX0EMPTY | MBOX0FULL | 0x02 | R/W |
| 0x92 | MBOX0B0 | MBOX0[7:0] | | | | | | | | 0x00 | R/W |
| 0x93 | MBOX0B1 | MBOX0[15:8] | | | | | | | | 0x00 | R/W |
| 0x94 | MBOX0B2 | MBOX0[23:16] | | | | | | | | 0x00 | R/W |
| 0x95 | MBOX0B3 | MBOX0[31:24] | | | | | | | | 0x00 | R/W |
| 0x96 | MBOX1CTL | Reserved | | MBOX1LEN | | MBOX1FIEN | MBOX1E1EN | MBOX1DIR | MBOX1EN | 0x02 | R/W |
| 0x97 | MBOX1STAT | Reserved | | MBOX1E1IRQ | MBOX1F1IRQ | Reserved | | MBOX1EMPTY | MBOX1FULL | 0x02 | R/W |
| 0x98 | MBOX1B0 | MBOX1[7:0] | | | | | | | | 0x00 | R/W |
| 0x99 | MBOX1B1 | MBOX1[15:8] | | | | | | | | 0x00 | R/W |
| 0x9A | MBOX1B2 | MBOX1[23:16] | | | | | | | | 0x00 | R/W |
| 0x9B | MBOX1B3 | MBOX1[31:24] | | | | | | | | 0x00 | R/W |

7 AD2428 A2B Register Descriptions

The transceiver (A2B) contains the following registers.

Table 7-1: AD2428 A2B Register List

| Name | Description |
|----------------|--|
| A2B_CHIP | I2C Chip Address Register (Target Only) |
| A2B_NODEADR | Node Address Register (Main Node Only) |
| A2B_VENDOR | Vendor ID Register |
| A2B_PRODUCT | Product ID Register |
| A2B_VERSION | Version ID Register |
| A2B_CAPABILITY | Capability ID Register |
| A2B_SWCTL | Switch Control Register |
| A2B_BCDNSLOTS | Broadcast Downstream Slots Register (Subordinate Only) |
| A2B_LDNSLOTS | Local Downstream Slots Register (Subordinate Only) |
| A2B_LUPLSLOTS | Local Upstream Slots Register (Subordinate Only) |
| A2B_DNSLOTS | Downstream Slots Register |
| A2B_UPSLOTS | Upstream Slots Register |
| A2B_RESPCYCS | Response Cycles Register |
| A2B_SLOTFMT | Slot Format Register (Main Node Only, Auto-Broadcast) |
| A2B_DATCTL | Data Control Register (Main Node Only, Auto-Broadcast) |
| A2B_CONTROL | Control Register |
| A2B_DISCVRY | Discovery Register (Main Node Only) |
| A2B_SWSTAT | Switch Status Register |
| A2B_INTSTAT | Interrupt Status Register |
| A2B_INTSRC | Interrupt Source Register (Main Node Only) |
| A2B_INTTYPE | Interrupt Type Register (Main Node Only) |
| A2B_INTPND0 | Interrupt Pending 0 Register |

Table 7-1: AD2428 A2B Register List (Continued)

| Name | Description |
|-----------------|---|
| A2B_INTPND1 | Interrupt Pending 1 Register |
| A2B_INTPND2 | Interrupt Pending 2 Register (Main Node Only) |
| A2B_INTMSK0 | Interrupt Mask 0 Register |
| A2B_INTMSK1 | Interrupt Mask 1 Register |
| A2B_INTMSK2 | Interrupt Mask 2 Register (Main Node Only) |
| A2B_BECCTL | Bit Error Count Control Register |
| A2B_BECCNT | Bit Error Count Register |
| A2B_TESTMODE | Testmode Register |
| A2B_ERRCNT0 | PRBS Error Count Byte 0 Register |
| A2B_ERRCNT1 | PRBS Error Count Byte 1 Register |
| A2B_ERRCNT2 | PRBS Error Count Byte 2 Register |
| A2B_ERRCNT3 | PRBS Error Count Byte 3 Register |
| A2B_NODE | Node Register |
| A2B_DISCSTAT | Discovery Status Register (Main Node Only) |
| A2B_TXACTL | LVDSA TX Control Register |
| A2B_TXBCTL | LVDSB TX Control Register |
| A2B_LINTTYPE | Local Interrupt Type (Subordinate Only) |
| A2B_I2CCFG | I2C Configuration Register |
| A2B_PLLCTL | PLL Control Register |
| A2B_I2SGCFG | I2S Global Configuration Register |
| A2B_I2SCFG | I2S Configuration Register |
| A2B_I2SRATE | I2S Rate Register (Target Only) |
| A2B_I2STXOFFSET | I2S Transmit Data Offset Register (Controller Only) |
| A2B_I2SRXOFFSET | I2S Receive Data Offset Register (Controller Only) |
| A2B_SYNCOFFSET | SYNC Offset Register (Subordinate Only) |
| A2B_PDMCTL | PDM Control Register |
| A2B_ERRMGMT | Error Management Register |
| A2B_GPIODAT | GPIO Output Data Register |
| A2B_GPIODATSET | GPIO Output Data Set Register |
| A2B_GPIODATCLR | GPIO Output Data Clear Register |
| A2B_GPIOOEN | GPIO Output Enable Register |

Table 7-1: AD2428 A2B Register List (Continued)

| Name | Description |
|----------------|--|
| A2B_GPIOIEN | GPIO Input Enable Register |
| A2B_GPIOIN | GPIO Input Value Register |
| A2B_PINTEN | Pin Interrupt Enable Register |
| A2B_PINTINV | Pin Interrupt Invert Register |
| A2B_PINCFG | Pin Configuration Register |
| A2B_I2STEST | I2S Test Register |
| A2B_RAISE | Raise Interrupt Register |
| A2B_GENERR | Generate Bus Error |
| A2B_I2SRRATE | I2S Reduced Rate Register (Controller Node Only, Auto-Broadcast) |
| A2B_I2SRRCTL | I2S Reduced Rate Control Register |
| A2B_I2SRRSOFFS | I2S Reduced Rate SYNC Offset Register (Target Only) |
| A2B_CLK1CFG | CLKOUT1 Configuration Register |
| A2B_CLK2CFG | CLKOUT2 Configuration Register |
| A2B_BMMCFG | Bus Monitor Mode Configuration Register |
| A2B_SUSCFG | Sustain Configuration Register (Subordinate Only) |
| A2B_PDMCTL2 | PDM Control 2 Register |
| A2B_UPMASK0 | Upstream Data RX Mask 0 Register (Target Only) |
| A2B_UPMASK1 | Upstream Data RX Mask 1 Register (Target Only) |
| A2B_UPMASK2 | Upstream Data RX Mask 2 Register (Target Only) |
| A2B_UPMASK3 | Upstream Data RX Mask 3 Register (Target Only) |
| A2B_UPOFFSET | Local Upstream Channel Offset Register (Target Only) |
| A2B_DNMASK0 | Downstream Data RX Mask 0 Register (Target Only) |
| A2B_DNMASK1 | Downstream Data RX Mask 1 Register (Target Only) |
| A2B_DNMASK2 | Downstream Data RX Mask 2 Register (Target Only) |
| A2B_DNMASK3 | Downstream Data RX Mask 3 Register (Target Only) |
| A2B_DNOFFSET | Local Downstream Channel Offset Register (Target Only) |
| A2B_CHIPID0 | Chip ID Register 0 |
| A2B_CHIPID1 | Chip ID Register 1 |
| A2B_CHIPID2 | Chip ID Register 2 |
| A2B_CHIPID3 | Chip ID Register 3 |
| A2B_CHIPID4 | Chip ID Register 4 |

Table 7-1: AD2428 A2B Register List (Continued)

| Name | Description |
|---------------|---|
| A2B_CHIPID5 | Chip ID Register 5 |
| A2B_GPIODEN | GPIO Over Distance Enable Register |
| A2B_GPIOD0MSK | GPIO Over Distance Mask 0 Register |
| A2B_GPIOD1MSK | GPIO Over Distance Mask 1 Register |
| A2B_GPIOD2MSK | GPIO Over Distance Mask 2 Register |
| A2B_GPIOD3MSK | GPIO Over Distance Mask 3 Register |
| A2B_GPIOD4MSK | GPIO Over Distance Mask 4 Register |
| A2B_GPIOD5MSK | GPIO Over Distance Mask 5 Register |
| A2B_GPIOD6MSK | GPIO Over Distance Mask 6 Register |
| A2B_GPIOD7MSK | GPIO Over Distance Mask 7 Register |
| A2B_GPIODDAT | GPIO Over Distance Data Register |
| A2B_GPIODINV | GPIO Over Distance Invert Register |
| A2B_MBOX0CTL | Mailbox 0 Control Register (Subordinate Only) |
| A2B_MBOX0STAT | Mailbox 0 Status Register (Subordinate Only) |
| A2B_MBOX0B0 | Mailbox 0 Byte 0 Register (Subordinate Only) |
| A2B_MBOX0B1 | Mailbox 0 Byte 1 Register (Subordinate Only) |
| A2B_MBOX0B2 | Mailbox 0 Byte 2 Register (Subordinate Only) |
| A2B_MBOX0B3 | Mailbox 0 Byte 3 Register (Subordinate Only) |
| A2B_MBOX1CTL | Mailbox 1 Control Register (Subordinate Only) |
| A2B_MBOX1STAT | Mailbox 1 Status Register (Subordinate Only) |
| A2B_MBOX1B0 | Mailbox 1 Byte 0 Register (Subordinate Only) |
| A2B_MBOX1B1 | Mailbox 1 Byte 1 Register (Subordinate Only) |
| A2B_MBOX1B2 | Mailbox 1 Byte 2 Register (Subordinate Only) |
| A2B_MBOX1B3 | Mailbox 1 Byte 3 Register (Subordinate Only) |

I2C Chip Address Register (Target Only)

The `A2B_CHIP` register stores a 7-bit I²C chip address. It is used during I²C transactions to address a remote peripheral device connected to a target node. The A²B target node acts as the I²C controller in I²C transactions with peripherals. This register only has an effect on I²C when it is programmed in a target node. The register can be written to and read from in a controller node without any influence on the chip's functionality.

Address: 0x00

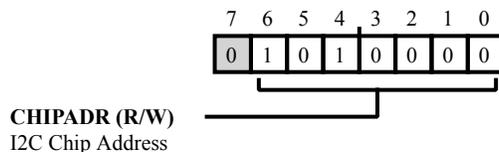


Figure 7-1: A2B_CHIP Register Diagram

Table 7-2: A2B_CHIP Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 6:0 (R/W) | CHIPADR | I2C Chip Address. The <code>A2B_CHIP.CHIPADR</code> bit field stores the I ² C address used by a target transceiver for I ² C accesses to a locally-connected peripheral. The A ² B target node acts as the I ² C controller in I ² C transactions with peripherals. |

Node Address Register (Main Node Only)

The `A2B_NODEADR` register provides control bits for addressing subordinate nodes through the A²B bus. This register can only be written in the main node. A write to this address in a subordinate node has no effect.

Address: 0x01

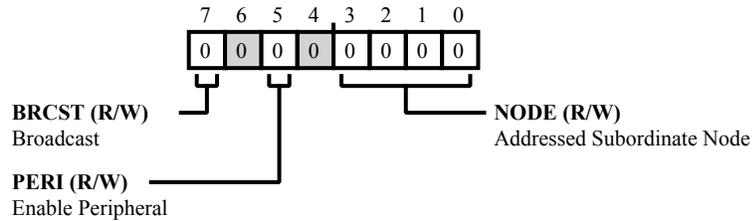


Figure 7-2: A2B_NODEADR Register Diagram

Table 7-3: A2B_NODEADR Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 7 (R/W) | BRCST | Broadcast. The <code>A2B_NODEADR.BRCST</code> bit enables broadcast mode. When an I ² C write with <code>BUS_ADDR</code> occurs in broadcast mode, the same control data is written to all nodes (controller and targets) simultaneously. The broadcast allows simultaneous control of all discovered A ² B transceivers, but not their respective I ² C peripherals. Therefore, clear the <code>A2B_NODEADR.PERI</code> bit (=0) when the <code>A2B_NODEADR.BRCST</code> bit is set to 1. |
| | | 0 Normal, directed register access |
| | | 1 Write to all nodes handled as broadcast access |
| 5 (R/W) | PERI | Enable Peripheral. The <code>A2B_NODEADR.PERI</code> bit enables register access (over I ² C) of peripheral devices on target nodes. The <code>A2B_NODEADR.BRCST</code> bit must be cleared (=0) when the <code>A2B_NODEADR.PERI</code> bit is set. When accessing target node registers through <code>BUS_ADDR</code> , the <code>A2B_NODEADR.PERI</code> bit must be cleared. |
| | | 0 Remote peripheral access disabled |
| | | 1 Remote peripheral access enabled |

Table 7-3: A2B_NODEADR Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration | |
|---------------------|----------|---|-------------|
| 3:0 (R/W) | NODE | <p>Addressed Subordinate Node.</p> <p>The A2B_NODEADR.NODE bit field selects a subordinate node by its address. Addresses are assigned based on the position in the A²B topology, starting with address 0 for the node connected directly to the main node. The value of the A2B_NODEADR.NODE field is irrelevant when the A2B_NODEADR.BRCST bit is set.</p> | |
| | | 0-9 | Node number |
| | | 10-15 | Reserved |

Vendor ID Register

The `A2B_VENDOR` register identifies the part as manufactured by Analog Devices.

Address: 0x02

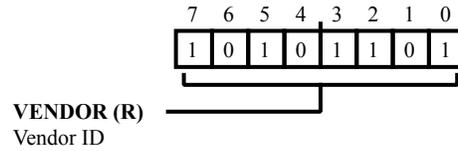


Figure 7-3: A2B_VENDOR Register Diagram

Table 7-4: A2B_VENDOR Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7:0 (R/NW) | VENDOR | Vendor ID. The <code>A2B_VENDOR.VENDOR</code> bit field contains the vendor identification number of the transceiver chip. |

Product ID Register

The `A2B_PRODUCT` register identifies the last two digits of the part number in hexadecimal format (for example, `0x26=AD2426W`).

Address: `0x03`

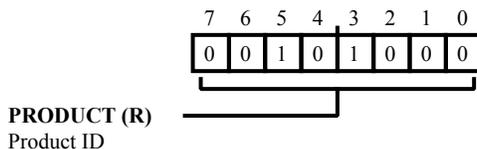


Figure 7-4: A2B_PRODUCT Register Diagram

Table 7-5: A2B_PRODUCT Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7:0 (R/NW) | PRODUCT | Product ID. The <code>A2B_PRODUCT</code> . <code>PRODUCT</code> bit field contains the product identification number of the transceiver. |

Version ID Register

The `A2B_VERSION` register identifies the version of the part.

Address: 0x04

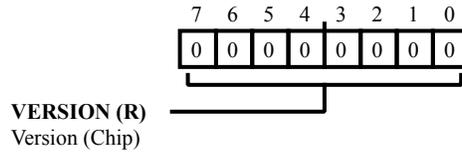


Figure 7-5: A2B_VERSION Register Diagram

Table 7-6: A2B_VERSION Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 7:0 (R/NW) | VERSION | Version (Chip). The <code>A2B_VERSION.VERSION</code> bit field contains the production version number of the chip. Bits 7:4 indicate major product revisions, while bits 3:0 are for minor revisions. |

Capability ID Register

The `A2B_CAPABILITY` register identifies available control interfaces. Transceivers that have an EEPROM storage device connected can store specific descriptor information in the EEPROM module.

Address: 0x05

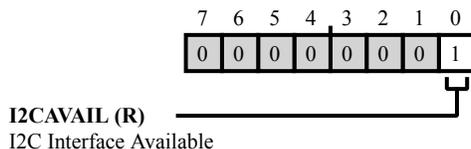


Figure 7-6: A2B_CAPABILITY Register Diagram

Table 7-7: A2B_CAPABILITY Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 0 (R/NW) | I2CAVAIL | I2C Interface Available. The <code>A2B_CAPABILITY.I2CAVAIL</code> bit signals availability of the I ² C interface on the transceiver for access to peripheral devices. If this bit is set (=1), module descriptor information can be accessible through the I ² C interface. A connected EEPROM (for example, an AT24C01) with module descriptor information must have an I ² C device address of 0x50. |
| | | 0 No I ² C interface is available |
| | | 1 I ² C interface is available |

Switch Control Register

The `A2B_SWCTL` register controls the switching of A²B bus power onto the downstream B-side of the A²B bus. This register must be written over the A²B bus. A write to this register from the local I²C port has no effect.

Note: The `A2B_SWCTL.DIAGMODE` bit must only be set when localizing the faults. Under all other conditions, the bit must be cleared to ensure proper operation of the device.

Address: 0x09

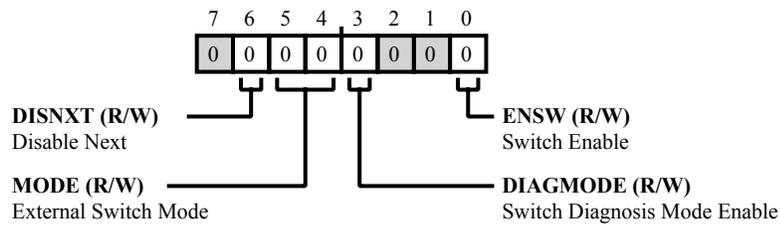


Figure 7-7: A2B_SWCTL Register Diagram

Table 7-8: A2B_SWCTL Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 6 (R/W) | DISNXT | <p>Disable Next.</p> <p>The <code>A2B_SWCTL.DISNXT</code> bit controls when packets are sent to the next node after the switch is enabled (<code>A2B_SWCTL.ENSW=1</code>). When <code>A2B_SWCTL.DISNXT</code> is cleared, synchronization packets are automatically passed to the next node immediately after the <code>A2B_SWSTAT.FIN</code> bit is set by the transceiver (signaling successful switching).</p> <p>When set, synchronization packets are not sent to the next node. A²B bus activity does not commence until discovery frames are issued when the <code>A2B_DISCVRY</code> register is programmed.</p> |
| | | 0 Enable synchronization packets |
| | | 1 Disable synchronization packets |

Table 7-8: A2B_SWCTL Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 5:4 (R/W) | MODE | External Switch Mode. The A2B_SWCTL.MODE bit field defines the diagnostic fault detection method for biasing the B-side A ² B bus with bus power for the next node. The setting depends on the external hardware configuration. When A2B_SWCTL.MODE= 0, the internal switch is configured for negative bias on the VSSN pin, and an external switch is required on the SWP pin for full line diagnostics, as well as localization and automatic line isolation after a fault is detected. When A2B_SWCTL.MODE= 1, the downstream node is not using A ² B bus power and is not properly terminating the bias. In this mode, open and reverse wire faults are not diagnosed, but all other fault types are diagnosed as long as the hardware configuration of the local node is as described for mode 0. When A2B_SWCTL.MODE=2, the voltage on the VIN pin (for example, 5 V) differs significantly from the bias voltage (8 V) on the SENSE pin. This applies when an extra regulator feeds the VIN pin. |
| | | 0 Use internal switch for VSSN pin and external switch for SWP pin |
| | | 1 Downstream node not using A ² B bus power and not properly terminating the bias |
| | | 2 Voltage on the VIN pin |
| | | 3 Reserved |
| 3 (R/W) | DIAGMODE | Switch Diagnosis Mode Enable. The A2B_SWCTL.DIAGMODE bit enables switch diagnosis mode. |
| | | 0 Switch diagnosis mode disabled |
| | | 1 Switch diagnosis mode enabled |
| 0 (R/W) | ENSW | Switch Enable. The A2B_SWCTL.ENSW bit enables A ² B bus power switching. |
| | | 0 Switch disabled |
| | | 1 Switch enabled |

Broadcast Downstream Slots Register (Subordinate Only)

In a subordinate node, the `A2B_BCDNSLOTS` register defines the number of data slots that are captured by the node and also passed downstream (B-side) as broadcast data to the next node. If any bits are set in the `A2B_DNMASK0` through `A2B_DNMASK3` registers, the value of the `A2B_BCDNSLOTS` register is ignored. Changes to this register only take effect after setting the `A2B_CONTROL.NEWSTRCT` bit of the main node. The `A2B_BCDNSLOTS` register is not used in the main node.

Address: 0x0A

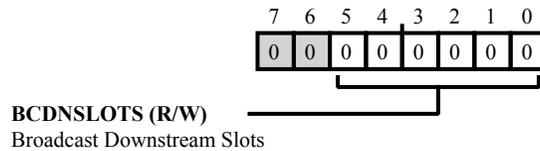


Figure 7-8: A2B_BCDNSLOTS Register Diagram

Table 7-9: A2B_BCDNSLOTS Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|-----------|---|
| 5:0 (R/W) | BCDNSLOTS | Broadcast Downstream Slots. The <code>A2B_BCDNSLOTS.BCDNSLOTS</code> bit field configures the number of broadcast downstream slots. This field must be programmed with a value between 0 and 32. |

Local Downstream Slots Register (Subordinate Only)

In a subordinate node, the meaning of the `A2B_LDNSLOTS` register changes depending on whether or not the downstream broadcast mask enable bit (`A2B_LDNSLOTS.DNMaskEN`) is set. If

`A2B_LDNSLOTS.DNMaskEN=0` (default), the `A2B_LDNSLOTS` register defines the number of data slots which are captured by the local node during the downstream portion of the superframe. These data slots are consumed by the node and are not passed downstream to the next node. If `A2B_LDNSLOTS.DNMaskEN=1`, the `A2B_LDNSLOTS` register defines the number of data slots that are added by the local node during the downstream portion of the superframe after `A2B_DNSLOTS.DNSLOTS` data slots are passed downstream by the transceiver. Changes to this register only take effect after setting the `A2B_CONTROL.NEWSTRCT` bit in the main node.

Address: 0x0B

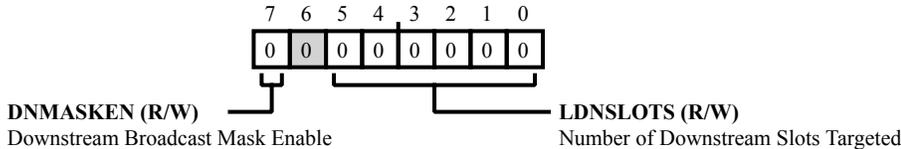


Figure 7-9: A2B_LDNSLOTS Register Diagram

Table 7-10: A2B_LDNSLOTS Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7 (R/W) | DNMaskEN | Downstream Broadcast Mask Enable. The <code>A2B_LDNSLOTS.DNMaskEN</code> bit enables the downstream mask enable bits in the <code>A2B_DNMask0</code> through <code>A2B_DNMask3</code> registers. |
| | | 0 Downstream data slot masks disabled |
| | | 1 Downstream data slot masks enabled |
| 5:0 (R/W) | LDNSLOTS | Number of Downstream Slots Targeted. When <code>A2B_LDNSLOTS.DNMaskEN=0</code> , the <code>A2B_LDNSLOTS.LDNSLOTS</code> bit field defines the number of data slots which are captured by the local node during the downstream portion of the superframe. When <code>A2B_LDNSLOTS.DNMaskEN=1</code> , the <code>A2B_LDNSLOTS.LDNSLOTS</code> bit field defines the number of data slots which are added by the local node during the downstream portion of the superframe. This field must be programmed with a value between 0 and 32 and be sufficient to accommodate all the data relative to its mode of TDM operation and the number of enabled data pins. |

Local Upstream Slots Register (Subordinate Only)

In a subordinate node, the `A2B_LUPSLOTS` register defines the number of data slots which are added by the local node during the upstream portion of the superframe. Changes to this register only take effect after setting the `A2B_CONTROL.NEWSTRCT` bit in the main node. The `A2B_LUPSLOTS` register is not used in the main node.

Address: 0x0C

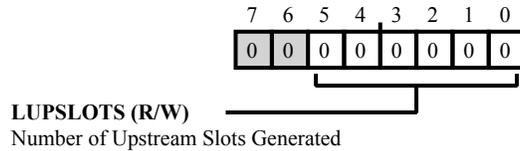


Figure 7-10: A2B_LUPSLOTS Register Diagram

Table 7-11: A2B_LUPSLOTS Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 5:0 (R/W) | LUPSLOTS | Number of Upstream Slots Generated. The <code>A2B_LUPSLOTS.LUPSLOTS</code> bit field defines the number of data slots which are added by the transceiver during the upstream portion of the superframe. These bits must be programmed with a value between 0 and 32. |

Downstream Slots Register

In a subordinate node, the [A2B_DNSLOTS](#) register defines the number of data slots (not including broadcast slots) that are passed downstream (B-side) after the transceiver begins to capture data slots. In the main node, the [A2B_DNSLOTS](#) register defines the total number of downstream data slots (including broadcast slots). Changes to this register only take effect after setting the [A2B_CONTROL.NEWSTRCT](#) bit in the main node.

Address: 0x0D

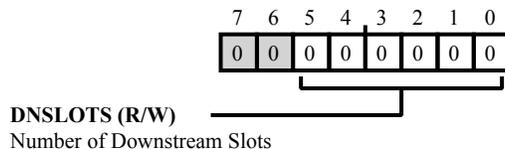


Figure 7-11: A2B_DNSLOTS Register Diagram

Table 7-12: A2B_DNSLOTS Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 5:0 (R/W) | DNSLOTS | <p>Number of Downstream Slots.</p> <p>In a main node, the <code>A2B_DNSLOTS.DNSLOTS</code> bit field is the number of downstream slots, including broadcast data slots. It must be sufficient to accommodate the data intended for downstream devices, which is a function of the TDM mode and the number of enabled data pins.</p> <p>In a subordinate node, the <code>A2B_DNSLOTS.DNSLOTS</code> bit field sets the number of data slots which are passed downstream. When calculating the value to program to this field, the same guidance as in the main node applies. But, subordinate nodes must also include any broadcast downstream slots, as programmed in the A2B_BCDNSLOTS register.</p> <p>Valid programming values are between 0 and 32.</p> |

Upstream Slots Register

In a subordinate node, the `A2B_UPSLOTS` register defines the number of data slots which are passed upstream by the B-side transceiver before the transceiver begins to add data slots. In the main node, the `A2B_UPSLOTS` register defines the total number of upstream data slots. Changes to this register only take effect after setting the `A2B_CONTROL.NEWSTRCT` bit in the main node.

Address: 0x0E

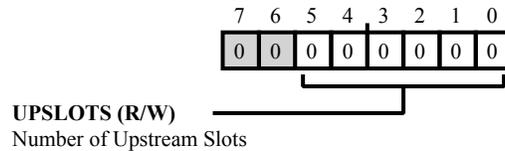


Figure 7-12: A2B_UPSLOTS Register Diagram

Table 7-13: A2B_UPSLOTS Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 5:0 (R/W) | UPSLOTS | <p>Number of Upstream Slots.</p> <p>In a main node, the <code>A2B_UPSLOTS.UPSLOTS</code> bit field is the number of upstream slots being received from the first-in-line subordinate node. It must be sufficient to accommodate all data intended for upstream devices, which is a function of TDM serial mode and the number of enabled data pins.</p> <p>In a subordinate node, the <code>A2B_UPSLOTS.UPSLOTS</code> bit field defines the number of data slots which are received from the next-in-line subordinate node and passed upstream before the transceiver begins to add data slots.</p> <p>Valid programming values are between 0 and 32.</p> |

Response Cycles Register

The `A2B_RESPCYCS` register defines the time between the start of the downstream header (the first SCF preamble bit) and the start of the upstream header (the first SCF preamble bit). The value in the register represents the number of bus bit times multiplied by 4. 1024 bit counts are in an A²B superframe between SCFs. One bus bit time = $1/(f_{\text{SYSBCLK}})$.

The `A2B_DISCVRY` register in the main transceiver is programmed with the `A2B_RESPCYCS` register value during discovery. Changes to this register only take effect after setting the `A2B_CONTROL.NEWSTRCT` bit in the main node. This register must be written over the A²B bus, as writes to this register from the local I²C port have no effect.

Address: 0x0F

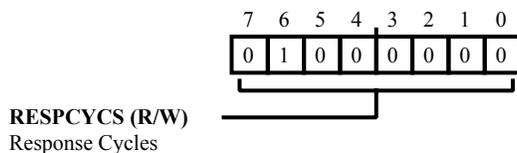


Figure 7-13: A2B_RESPCYCS Register Diagram

Table 7-14: A2B_RESPCYCS Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 7:0 (R/W) | RESPCYCS | Response Cycles. The <code>A2B_RESPCYCS.RESPCYCS</code> bit field is one-fourth the time (in terms of bus bits) from the start of a downstream frame to the start of an upstream frame. |

Slot Format Register (Main Node Only, Auto-Broadcast)

The `A2B_SLOTFMT` register defines the size and format of the downstream and upstream data slots. Floating-point compression of A²B data can be enabled to reduce bandwidth using this register, and ECC protection of A²B data can alternately be enabled. All nodes in an A²B system are subject to the same upstream and downstream slot format setting. Changes to this register only take effect after setting the `A2B_CONTROL.NEWSTRCT` bit in the main node.

When the `A2B_SLOTFMT` register is written in the main node, the new setting is automatically broadcast to all discovered subordinate nodes over the A²B bus. Local host writes to this register in a subordinate node have no effect.

Address: 0x10

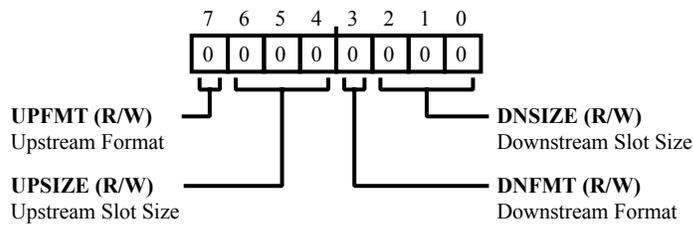


Figure 7-14: A2B_SLOTFMT Register Diagram

Table 7-15: A2B_SLOTFMT Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration | |
|---------------------|----------|--|-------------------------------------|
| 7 (R/W) | UPFMT | <p>Upstream Format.</p> <p>The A2B_SLOTFMT.UPFMT bit configures the format of the upstream data on the A²B bus bus. When A2B_SLOTFMT.UPFMT= 0, the format of the data on the A²B bus bus is normal (no compression, no ECC protection, and one parity bit). When A2B_SLOTFMT.UPFMT = 1, an alternate data format is utilized, depending on the upstream data width (A2B_SLOTFMT.UPSIZE).</p> <p>When the A2B_SLOTFMT.UPSIZE bit is programmed for 12-, 16-, or 20-bit data, setting the A2B_SLOTFMT.UPFMT bit enables floating-point compression of upstream data. When this compression is used, the I²S/TDM or PDM data is 4 bits wider than the A²B data, which is compressed to reduce A²B bus bandwidth, and the data is protected by a parity bit.</p> <p>When the A2B_SLOTFMT.UPSIZE bit is programmed for 24- or 32-bit data, setting the A2B_SLOTFMT.UPFMT bit enables ECC protection on upstream data slots, where ECC bits are added to each data slot instead of a parity bit (6 ECC bits for 24-bit data, 7 ECC bits for 32-bit data).</p> <p>Setting the A2B_SLOTFMT.UPFMT bit when A2B_SLOTFMT.UPSIZE is programmed for 8- or 28-bit data has no effect.</p> | |
| | | 0 | Normal upstream data slot format |
| | | 1 | Alternate upstream data slot format |
| 6:4 (R/W) | UPSIZE | <p>Upstream Slot Size.</p> <p>The A2B_SLOTFMT.UPSIZE bit field selects the upstream data slot size.</p> | |
| | | 0 | 8 bits |
| | | 1 | 12 bits |
| | | 2 | 16 bits |
| | | 3 | 20 bits |
| | | 4 | 24 bits |
| | | 5 | 28 bits |
| | | 6 | 32 bits |

Table 7-15: A2B_SLOTFMT Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 3 (R/W) | DNFMT | <p>Downstream Format.</p> <p>The A2B_SLOTFMT.DNFMT bit configures the format of the downstream data on the A²B bus bus. When A2B_SLOTFMT.DNFMT= 0, the format of the data on the A²B bus bus is normal (no compression, no ECC protection, and one parity bit). When A2B_SLOTFMT.DNFMT = 1, an alternate data format is utilized, depending on the downstream data width (A2B_SLOTFMT.DNSIZE).</p> <p>When the A2B_SLOTFMT.DNSIZE field is programmed for 12-, 16-, or 20-bit data, setting the A2B_SLOTFMT.DNFMT bit enables floating-point compression of downstream data. When this compression is used, the I²S/TDM or PDM data is 4 bits wider than the A²B data, which is compressed to reduce A²B bus bandwidth, and the data is protected by a parity bit.</p> <p>When the A2B_SLOTFMT.DNSIZE bit is programmed for 24- or 32-bit data, setting the A2B_SLOTFMT.DNFMT bit enables ECC protection on downstream data slots, where ECC bits are added to each data slot instead of a parity bit (6 ECC bits for 24-bit data, 7 ECC bits for 32-bit data).</p> <p>Setting the A2B_SLOTFMT.DNFMT bit when A2B_SLOTFMT.DNSIZE is programmed for 8- or 28-bit data has no effect.</p> |
| | | 0 Normal downstream data slot format |
| | | 1 Alternate downstream data slot format |
| 2:0 (R/W) | DNSIZE | <p>Downstream Slot Size.</p> <p>The A2B_SLOTFMT.DNSIZE bit field selects the downstream data slot size.</p> |
| | | 0 8 bits |
| | | 1 12 bits |
| | | 2 16 bits |
| | | 3 20 bits |
| | | 4 24 bits |
| | | 5 28 bits |
| | | 6 32 bits |
| | | 7 Reserved |

Data Control Register (Main Node Only, Auto-Broadcast)

The `A2B_DATCTL` register is used to enable data slots and standby mode on the A²B bus. Changes to this register only take effect after setting the `A2B_CONTROL.NEWSTRCT` bit in the main node. When the `A2B_DATCTL` register is written in the main node, the new setting is automatically broadcast to all discovered subordinate node over the A²B bus. Local host writes to this register in a subordinate node have no effect.

NOTE: To switch back to normal operation, first exit the standby mode by clearing the `A2B_DATCTL.STANDBY` bit, then write to the `A2B_DATCTL` register to enable the upstream and downstream slots.

Address: 0x11

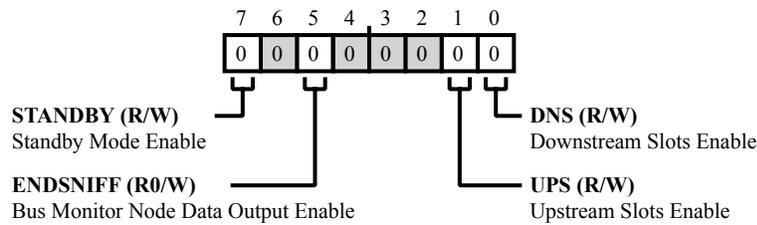


Figure 7-15: A2B_DATCTL Register Diagram

Table 7-16: A2B_DATCTL Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7 (R/W) | STANDBY | Standby Mode Enable. The <code>A2B_DATCTL.STANDBY</code> bit globally enables power saving mode for all nodes and minimizes bus activity. The only traffic required is a minimal downstream preamble to keep all of the PLLs in the subordinate nodes synchronized. Reads and writes across the A ² B bus are not supported in this mode. |
| | | 0 Disabled |
| | | 1 Enabled |
| 5 (R0/W) | ENDSNIFF | Bus Monitor Node Data Output Enable. The <code>A2B_DATCTL.ENDSNIFF</code> bit controls whether or not an attached Bus Monitor Node produces data slots as output. |
| | | 0 Disabled |
| | | 1 Enabled |
| 1 (R/W) | UPS | Upstream Slots Enable. The <code>A2B_DATCTL.UPS</code> bit globally enables upstream synchronous data to be sent over the bus. |
| | | 0 Disabled |
| | | 1 Enabled |

Table 7-16: A2B_DATCTL Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration | |
|---------------------|----------|--|----------|
| 0 (R/W) | DNS | Downstream Slots Enable. The A2B_DATCTL.DNS bit globally enables downstream synchronous data to be sent over the bus. | |
| | | 0 | Disabled |
| | | 1 | Enabled |

Control Register

The `A2B_CONTROL` register provides bits which control nodes on the bus.

Address: 0x12

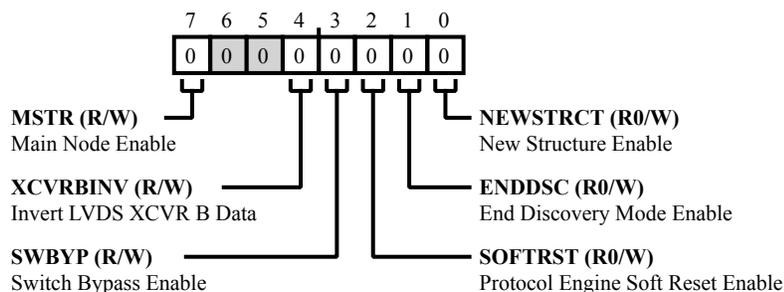


Figure 7-16: A2B_CONTROL Register Diagram

Table 7-17: A2B_CONTROL Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7 (R/W) | MSTR | Main Node Enable. The <code>A2B_CONTROL.MSTR</code> bit controls whether the current node is a subordinate node or a main node. |
| | | 0 Subordinate node |
| | | 1 Main node |
| 4 (R/W) | XCVRBINV | Invert LVDS XCVR B Data. The <code>A2B_CONTROL.XCVRBINV</code> bit controls an optional inversion of data to/from LVDS XCVR B. Data is inverted when this bit is set. |
| 3 (R/W) | SWBYP | Switch Bypass Enable. The <code>A2B_CONTROL.SWBYP</code> bit enables the downstream LVDS XCVR without waiting for the line switch to be turned on. When this bit is set the line switch will not be enabled even if <code>A2B_SWCTL.ENS</code> is set. |
| 2 (R0/W) | SOFTRST | Protocol Engine Soft Reset Enable. When the <code>A2B_CONTROL.SOFTRST</code> bit is set, the protocol engine in the bus node is reset, and all registers return to their respective reset states. |
| | | 0 No action |
| | | 1 Reset protocol engine |

Table 7-17: A2B_CONTROL Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 1 (R0/W) | ENDDSC | End Discovery Mode Enable. In the main node, setting the A2B_CONTROL.ENDDSC bit ends discovery attempts to a new subordinate node. |
| | | 0 No action |
| | | 1 End discovery |
| 0 (R0/W) | NEWSTRCT | New Structure Enable. The A2B_CONTROL.NEWSTRCT bit synchronously applies a new structure to all nodes. When the A2B_CONTROL.NEWSTRCT bit is set in the main node, a new structure is applied within 5 superframe cycles unless communication errors create delays. The A2B_CONTROL.NEWSTRCT bit must be set before changes to some registers take effect. Do not write to any of these registers for 5 superframe periods after setting the A2B_CONTROL.NEWSTRCT bit. See the register descriptions for details. |
| | | 0 No action |
| | | 1 Enable new structure |

Discovery Register (Main Node Only)

Programming the `A2B_DISCVRY` register with a response cycle value for a new node to be added allows the new subordinate node to be discovered. It triggers the start of full discovery frames being sent to the next-in-line subordinate node.

When the `A2B_DISCVRY` register is written in the main node, the new setting is automatically broadcast to all subordinate nodes over the A²B bus. Local host and direct `BUS_ADDR` writes to this register in a subordinate node have no effect.

Address: 0x13

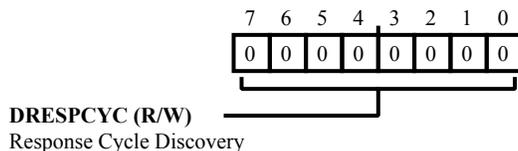


Figure 7-17: A2B_DISCVRY Register Diagram

Table 7-18: A2B_DISCVRY Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7:0 (R/W) | DRESPCYC | Response Cycle Discovery. The <code>A2B_DISCVRY.DRESPCYC</code> bit field is written with the value to be used for <code>A2B_RESPCYCS</code> by a to-be discovered subordinate node. |

Switch Status Register

The `A2B_SWSTAT` register provides line diagnostics status information. Line diagnostics are performed when bias is switched onto the A²B bus towards the next-in-line subordinate node.

Address: 0x14

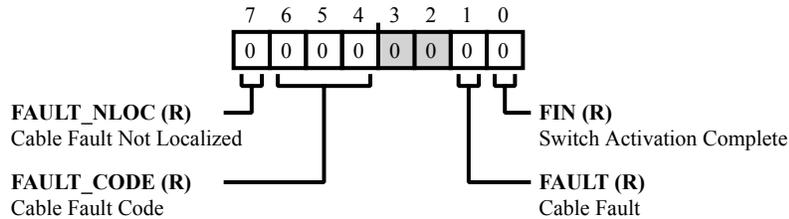


Figure 7-18: A2B_SWSTAT Register Diagram

Table 7-19: A2B_SWSTAT Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|------------|---|
| 7 (R/NW) | FAULT_NLOC | Cable Fault Not Localized. The <code>A2B_SWSTAT.FAULT_NLOC</code> bit indicates that the identified line fault is not localized. |
| | | 0 Switch fault localized |
| | | 1 Switch fault not localized |
| 6:4 (R/NW) | FAULT_CODE | Cable Fault Code. The <code>A2B_SWSTAT.FAULT_CODE</code> bit field contains downstream link cable diagnostic error codes. |
| | | 0 No fault detected |
| | | 1 Cable terminal shorted to GND |
| | | 2 Cable terminal shorted to VBAT |
| | | 3 Cable terminals shorted together |
| | | 4 Cable disconnected or open circuit |
| | | 5 Cable is reverse connected |
| | | 6 Reserved |
| | | 7 Undetermined fault |
| 1 (R/NW) | FAULT | Cable Fault. The <code>A2B_SWSTAT.FAULT</code> bit indicates a cable fault has been detected. |
| | | 0 Cable fault not detected |
| | | 1 Cable fault detected |

Table 7-19: A2B_SWSTAT Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration | |
|---------------------|----------|--|---|
| 0 (R/NW) | FIN | <p>Switch Activation Complete.</p> <p>The <code>A2B_SWSTAT.FIN</code> bit indicates the successful completion of the switch activation sequence for biasing of the downstream link. When this bit is set the transceiver begins passing SCFs to the next-in-line subordinate node, thus allowing it to begin locking its PLL, unless the switch is bypassed (<code>A2B_CONTROL.SWBYP = 1</code>)</p> | |
| | | 0 | Switch is open or has not completed closing |
| | | 1 | Switch completed closing |

Interrupt Status Register

The `A2B_INTSTAT` register contains interrupt status information for the node.

Address: 0x15

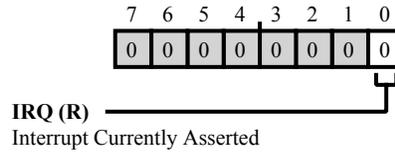


Figure 7-19: A2B_INTSTAT Register Diagram

Table 7-20: A2B_INTSTAT Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 0 (R/NW) | IRQ | Interrupt Currently Asserted. When the <code>A2B_INTSTAT</code> . <code>IRQ</code> bit is set, the node is signaling an interrupt request, either through the IRQ pin for a main node or over the A ² B bus for a subordinate node. |
| | | 0 No interrupt request |
| | | 1 Interrupt request |

Interrupt Source Register (Main Node Only)

The `A2B_INTSRC` register contains information about the current highest priority interrupt. It is updated when the `A2B_INTTYPE` register is read. A value of `0x00` in this register indicates that no interrupts are present.

Address: `0x16`

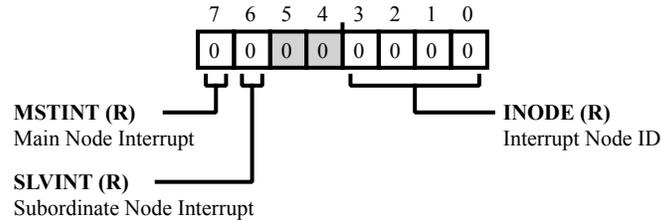


Figure 7-20: `A2B_INTSRC` Register Diagram

Table 7-21: `A2B_INTSRC` Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7 (R/NW) | MSTINT | Main Node Interrupt. When the <code>A2B_INTSRC.MSTINT</code> bit is set, the current interrupt is being generated by the main node. |
| 6 (R/NW) | SLVINT | Subordinate Node Interrupt. When the <code>A2B_INTSRC.SLVINT</code> bit is set, the current interrupt is being generated by a subordinate node. |
| 3:0 (R/NW) | INODE | Interrupt Node ID. The <code>A2B_INTSRC.INODE</code> bit field contains the node number of the subordinate node that asserted the current interrupt. |

Interrupt Type Register (Main Node Only)

The `A2B_INTTYPE` register contains information about the pending interrupt being generated by the node indicated in the `A2B_INTSRC` register and signaled with the IRQ pin. A host read of the `A2B_INTTYPE` register in the main node clears this pending interrupt in the main and deasserts the IRQ pin. If other interrupts are pending, the `A2B_INTSRC` and `A2B_INTTYPE` registers are updated to reflect the highest priority pending interrupt, and the IRQ pin is asserted again. Nodes closer to the main node have a higher priority when the same interrupt appears in more than one subordinate node.

Address: 0x17

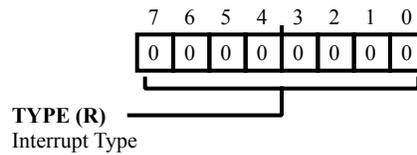


Figure 7-21: A2B_INTTYPE Register Diagram

Table 7-22: A2B_INTTYPE Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 7:0 (R/NW) | TYPE | Interrupt Type. The <code>A2B_INTTYPE</code> .TYPE bit field contains the current interrupt type. Interrupt types are described in the interrupt pending registers (<code>A2B_INTPND0</code> through <code>A2B_INTPND2</code>). |
| | | 0 HDCNTERR - Header count error |
| | | 1 DDERR - Data decoding error |
| | | 2 CRCERR - CRC error |
| | | 3 DPERR - Data parity error |
| | | 4 BECOVF - Bit error counter overflow error |
| | | 5 SRFERR - SRF miss error |
| | | 6 SRFCRCERR - SRF CRC error (subordinate node only) |
| | | 9 PWRERR - Positive terminal BP shorted to GND |
| | | 10 PWRERR - Negative terminal BN shorted to VBAT |
| | | 11 PWRERR - BP shorted to BN |
| | | 12 PWRERR - Cable disconnected or open circuit or wrong port |
| | | 13 PWRERR - Cable is reverse connected or wrong port |
| | | 15 PWRERR - Undetermined fault |

Table 7-22: A2B_INTTYPE Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| | | 16 IO0PND - GP input IO0 interrupt (subordinate node only) |
| | | 17 IO1PND - GP input IO1 interrupt |
| | | 18 IO2PND - GP input IO2 interrupt |
| | | 19 IO3PND - GP input IO3 interrupt |
| | | 20 IO4PND - GP input IO4 interrupt |
| | | 21 IO5PND - GP input IO5 interrupt |
| | | 22 IO6PND - GP input IO6 interrupt |
| | | 23 IO7PND - GP input IO7 interrupt |
| | | 24 DSCDONE - Discovery done interrupt (main node only) |
| | | 25 I2CERR - I2C error (main node only) |
| | | 26 ICRERR - Interrupt CRC error (main node only) |
| | | 41 PWRERR - Non-localized negative terminal BN short to GND |
| | | 42 PWRERR - Non-localized positive terminal BP short to VBAT |
| | | 48 Mailbox 0 full |
| | | 49 Mailbox 0 empty |
| | | 50 Mailbox 1 full |
| | | 51 Mailbox 1 empty |
| | | 128 Interrupt messaging error |
| | | 252 Startup error - Return to factory |
| | | 253 Subordinate INTTYPE read error - Main node only |
| | | 254 Standby done - Main node only |
| | | 255 MSTR_RUNNING - Main node PLL locked |

Interrupt Pending 0 Register

The `A2B_INTPND0` register contains interrupt pending bits for the node.

Address: 0x18

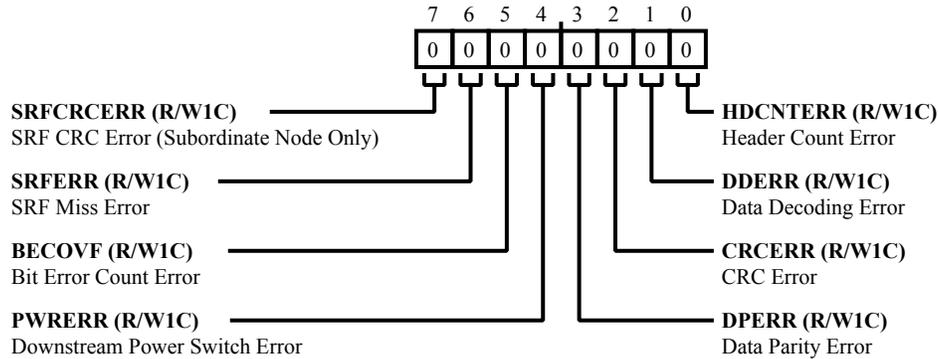


Figure 7-22: A2B_INTPND0 Register Diagram

Table 7-23: A2B_INTPND0 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|-----------|---|
| 7 (R/W1C) | SRFCRCERR | SRF CRC Error (Subordinate Node Only). The <code>A2B_INTPND0.SRFCRCERR</code> bit indicates that the current subordinate node has detected a CRC error in the SRF. |
| | | 0 No SRF CRC error |
| | | 1 SRF CRC error detected |
| 6 (R/W1C) | SRFERR | SRF Miss Error. The <code>A2B_INTPND0.SRFERR</code> bit indicates that the node has not received the SRF from the downstream node at the specified time. |
| | | 0 No SRF miss error |
| | | 1 SRF miss error detected |
| 5 (R/W1C) | BECOVF | Bit Error Count Error. The <code>A2B_INTPND0.BECOVF</code> bit indicates that the number of errors programmed into the bit error count control register has been exceeded. |
| | | 0 No BEC error pending |
| | | 1 BEC error pending |

Table 7-23: A2B_INTPND0 Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 4 (R/W1C) | PWRERR | Downstream Power Switch Error. The A2B_INTPND0.PWRERR bit indicates an error reported from the downstream power switch. |
| | | 0 No power error |
| | | 1 Downstream power switch error |
| 3 (R/W1C) | DPERR | Data Parity Error. The A2B_INTPND0.DPERR bit indicates that the current node has detected a data parity error. The error is detected only if the node consumes the data slot with a data parity error. |
| | | 0 No data parity error |
| | | 1 Data parity error detected |
| 2 (R/W1C) | CRCERR | CRC Error. The A2B_INTPND0.CRCERR bit indicates that the current node has detected a CRC error. For the main node, this applies to a CRC error in the SRF. For a subordinate node, this applies to a CRC error in the SCF. |
| | | 0 No CRC error |
| | | 1 CRC error detected |
| 1 (R/W1C) | DDERR | Data Decoding Error. The A2B_INTPND0.DDERR bit indicates that the current node has detected a data decoding error. The error is detected only if the node consumes the data slot with a data decoding error. |
| | | 0 No data decoding error |
| | | 1 Data decoding error detected |
| 0 (R/W1C) | HDCNTERR | Header Count Error. The A2B_INTPND0.HDCNTERR bit indicates the current node has detected a header count error. For the main node, this means that the SRF has a different count value than expected. For a subordinate node, this means that the SRF has a different value than expected. |
| | | 0 No header count error |
| | | 1 Header count error detected |

Interrupt Pending 1 Register

The `A2B_INTPND1` register contains interrupt pending bits for the node.

Address: 0x19

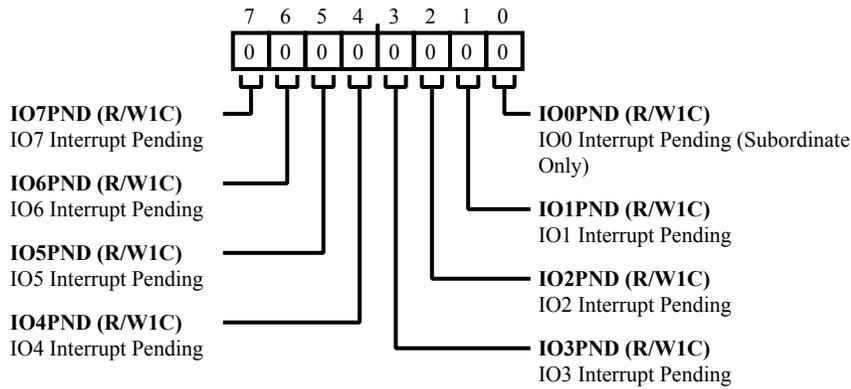


Figure 7-23: A2B_INTPND1 Register Diagram

Table 7-24: A2B_INTPND1 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 7 (R/W1C) | IO7PND | IO7 Interrupt Pending. The <code>A2B_INTPND1 . IO7PND</code> bit indicates that a pin interrupt request from IO7 is pending. |
| | | 0 No interrupt pending |
| | | 1 Interrupt pending <inherit> |
| 6 (R/W1C) | IO6PND | IO6 Interrupt Pending. The <code>A2B_INTPND1 . IO6PND</code> bit indicates that a pin interrupt request from IO6 (DRX1) is pending. |
| | | 0 No interrupt pending |
| | | 1 Interrupt pending |
| 5 (R/W1C) | IO5PND | IO5 Interrupt Pending. The <code>A2B_INTPND1 . IO5PND</code> bit indicates that a pin interrupt request from IO5 (DRX0) is pending. |
| | | 0 No interrupt pending |
| | | 1 Interrupt pending |

Table 7-24: A2B_INTPND1 Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 4 (R/W1C) | IO4PND | IO4 Interrupt Pending. The A2B_INTPND1 . IO4PND bit indicates that a pin interrupt request from IO4 (DTX1) is pending. |
| | | 0 No interrupt pending |
| | | 1 Interrupt pending |
| 3 (R/W1C) | IO3PND | IO3 Interrupt Pending. The A2B_INTPND1 . IO3PND bit indicates that a pin interrupt request from IO3 (DTX0) is pending. |
| | | 0 No interrupt pending |
| | | 1 Interrupt pending |
| 2 (R/W1C) | IO2PND | IO2 Interrupt Pending. The A2B_INTPND1 . IO2PND bit indicates that a pin interrupt request from IO2 (ADR2) is pending. |
| | | 0 No interrupt pending |
| | | 1 Interrupt pending |
| 1 (R/W1C) | IO1PND | IO1 Interrupt Pending. The A2B_INTPND1 . IO1PND bit indicates that a pin interrupt request from IO1 (ADR1) is pending. |
| | | 0 No interrupt pending |
| | | 1 Interrupt pending |
| 0 (R/W1C) | IO0PND | IO0 Interrupt Pending (Subordinate Only). The A2B_INTPND1 . IO0PND bit indicates that a pin interrupt request from IO0 (IRQ) is pending. On main nodes, this bit always reads 0. |
| | | 0 No interrupt pending <inherit> |
| | | 1 Interrupt pending |

Interrupt Pending 2 Register (Main Node Only)

The `A2B_INTPND2` register contains interrupt pending bits relevant only to main nodes.

Address: 0x1A

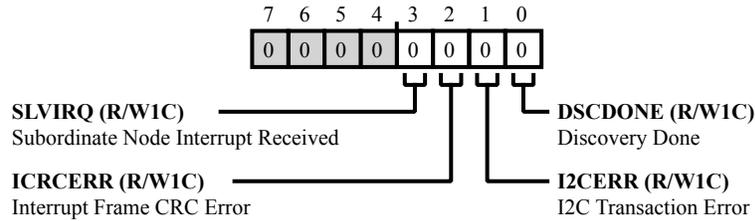


Figure 7-24: A2B_INTPND2 Register Diagram

Table 7-25: A2B_INTPND2 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 3 (R/W1C) | SLVIRQ | Subordinate Node Interrupt Received. In the main mode, the <code>A2B_INTPND2.SLVIRQ</code> bit indicates that a subordinate node has signaled an interrupt to the main node. This bit always reads zero in a subordinate node. |
| | | 0 No interrupt |
| | | 1 Subordinate node has signaled an interrupt |
| 2 (R/W1C) | ICRCERR | Interrupt Frame CRC Error. In the main mode, the <code>A2B_INTPND2.ICRCERR</code> bit indicates that the main node has detected an interrupt frame CRC error. |
| | | 0 No error |
| | | 1 Interrupt frame CRC error detected |
| 1 (R/W1C) | I2CERR | I2C Transaction Error. The <code>A2B_INTPND2.I2CERR</code> bit indicates that an I ² C access error has occurred. Examples of this are an I ² C write to a target node with early acknowledge that did not complete or a broadcast write that timed out. |
| | | 0 No error |
| | | 1 An I ² C access error occurred |
| 0 (R/W1C) | DSCDONE | Discovery Done. The <code>A2B_INTPND2.DSCDONE</code> bit indicates that a new subordinate node has been discovered. This bit always reads zero in subordinate nodes. |
| | | 0 No new subordinate node discovered |
| | | 1 New subordinate node discovered |

Interrupt Mask 0 Register

The `A2B_INTMSK0` register determines which `A2B_INTPND0` register bits generate interrupts.

Address: 0x1B

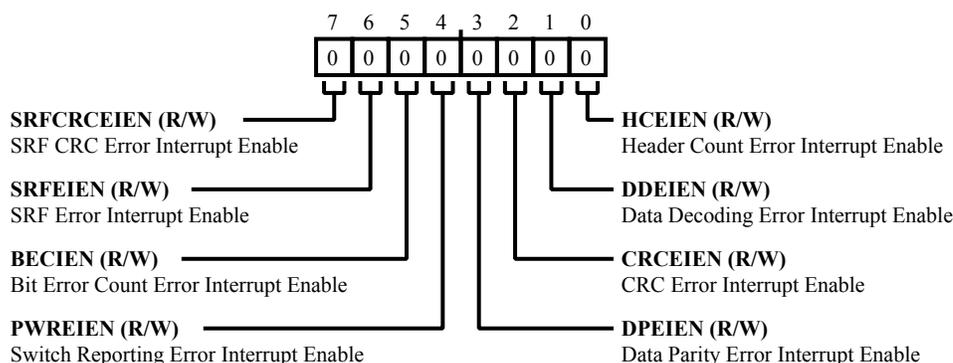


Figure 7-25: A2B_INTMSK0 Register Diagram

Table 7-26: A2B_INTMSK0 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|------------|--|
| 7 (R/W) | SRFCRCEIEN | SRF CRC Error Interrupt Enable. |
| 6 (R/W) | SRFEIEN | SRF Error Interrupt Enable. |
| 5 (R/W) | BECIEN | Bit Error Count Error Interrupt Enable. |
| 4 (R/W) | PWREIEN | Switch Reporting Error Interrupt Enable. |
| 3 (R/W) | DPEIEN | Data Parity Error Interrupt Enable. |
| 2 (R/W) | CRCEIEN | CRC Error Interrupt Enable. |
| 1 (R/W) | DDEIEN | Data Decoding Error Interrupt Enable. |
| 0 (R/W) | HCEIEN | Header Count Error Interrupt Enable. |

Interrupt Mask 1 Register

The `A2B_INTMSK1` register determines which `A2B_INTPND1` register bits generate interrupts.

Address: 0x1C

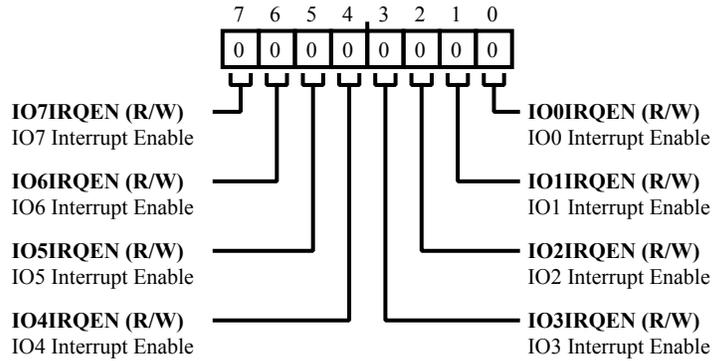


Figure 7-26: A2B_INTMSK1 Register Diagram

Table 7-27: A2B_INTMSK1 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|-------------------------|
| 7 (R/W) | IO7IRQEN | IO7 Interrupt Enable. |
| 6 (R/W) | IO6IRQEN | IO6 Interrupt Enable. |
| 5 (R/W) | IO5IRQEN | IO5 Interrupt Enable. |
| 4 (R/W) | IO4IRQEN | IO4 Interrupt Enable. |
| 3 (R/W) | IO3IRQEN | IO3 Interrupt Enable. |
| 2 (R/W) | IO2IRQEN | IO2 Interrupt Enable. |
| 1 (R/W) | IO1IRQEN | IO1 Interrupt Enable. |
| 0 (R/W) | IO0IRQEN | IO0 Interrupt Enable. |

Interrupt Mask 2 Register (Main Node Only)

The `A2B_INTMSK2` register determines which `A2B_INTPND2` register bits generate interrupts.

Address: 0x1D

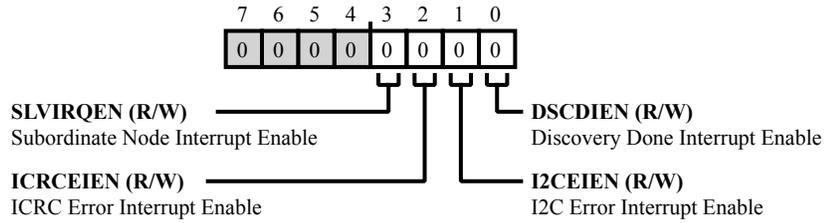


Figure 7-27: A2B_INTMSK2 Register Diagram

Table 7-28: A2B_INTMSK2 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|------------------------------------|
| 3 (R/W) | SLVIRQEN | Subordinate Node Interrupt Enable. |
| 2 (R/W) | ICRCEIEN | ICRC Error Interrupt Enable. |
| 1 (R/W) | I2CEIEN | I2C Error Interrupt Enable. |
| 0 (R/W) | DSCDIEN | Discovery Done Interrupt Enable. |

Bit Error Count Control Register

The `A2B_BECCTL` register controls bit error counting, including interrupt thresholds.

Address: 0x1E

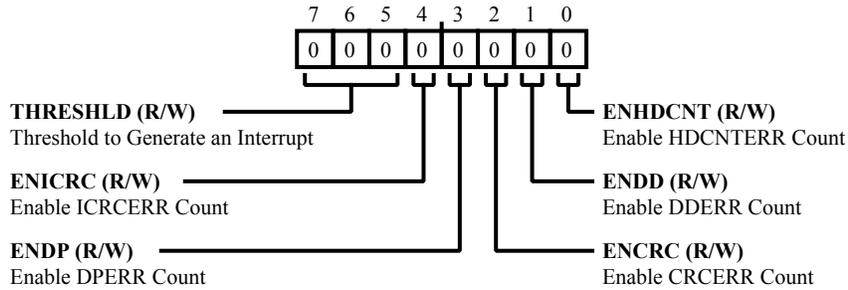


Figure 7-28: A2B_BECCTL Register Diagram

Table 7-29: A2B_BECCTL Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|--------------------------------|----------|---|
| 7:5 (R/W) | THRESHLD | Threshold to Generate an Interrupt. The <code>A2B_BECCTL.THRESHLD</code> bit field configures the number of errors counted before the <code>A2B_INTPND0.BECOVF</code> bit is set. |
| | | 0 Interrupt after 2 errors |
| | | 1 Interrupt after 4 errors |
| | | 2 Interrupt after 8 errors |
| | | 3 Interrupt after 16 errors |
| | | 4 Interrupt after 32 errors |
| | | 5 Interrupt after 64 errors |
| | | 6 Interrupt after 128 errors |
| 7 Interrupt after 256 errors | | |
| 4 (R/W) | ENICRC | Enable ICRERR Count. When the <code>A2B_BECCTL.ENICRC</code> bit is set, the bit error count register is incremented every time a CRC error is detected in the interrupt response frame. |
| | | 0 Disabled |
| | | 1 Enable bit error counting |

Table 7-29: A2B_BECCTL Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 3 (R/W) | ENDP | Enable DPERR Count. When the A2B_BECCTL.ENDP bit is set, the bit error count register is incremented on every parity error of the streaming data. |
| | | 0 No parity error |
| | | 1 Parity error |
| 2 (R/W) | ENCRC | Enable CRCERR Count. When the A2B_BECCTL.ENCRC bit is set, the bit error count register is incremented on every CRC error in a control or response frame. This excludes interrupt frame CRC errors. |
| | | 0 No CRC error |
| | | 1 CRC error |
| 1 (R/W) | ENDD | Enable DDERR Count. When the A2B_BECCTL.ENDD bit is set, the bit error count register is incremented on every data decoding error. |
| | | 0 Disabled |
| | | 1 Enabled |
| 0 (R/W) | ENHDCNT | Enable HDCNTERR Count. When the A2B_BECCTL.ENHDCNT bit is set, the bit error count register is incremented if there is a discrepancy between the actual and expected header count field. |
| | | 0 Disabled |
| | | 1 Enabled |

Bit Error Count Register

The `A2B_BECNT` register accumulates the error count of the error types selected in the `A2B_BECCTL` register. Any write to this register clears the count.

Address: 0x1F

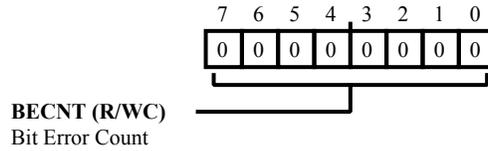


Figure 7-29: A2B_BECNT Register Diagram

Table 7-30: A2B_BECNT Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 7:0 (R/WC) | BECNT | Bit Error Count. The <code>A2B_BECNT.BECNT</code> bit field provides the number of bit errors counted, based on the value programmed into the <code>A2B_BECCTL</code> register. |

Testmode Register

The `A2B_TESTMODE` register provides control bits to be used in testing the A²B link. The `A2B_TESTMODE.PRBSDN` and `A2B_TESTMODE.PRBSUP` bits are used to enable the use of pseudo-random data in the downstream and upstream data slots on the A²B bus, respectively. Downstream data is checked in the last subordinate node based on the programming of the `A2B_DNSLOTS`, `A2B_LDNSLOTS`, and `A2B_BCDNSLOTS` registers. Upstream data is checked in the main node. Data mismatches increment a 32-bit counter (which can be read via the `A2B_ERRCNT0` through `A2B_ERRCNT3` registers). The `A2B_TESTMODE` register must be programmed via a broadcast write. subordinate to subordinate communications adversely affect a Bit Error Rate Test (BERT).

Address: 0x20

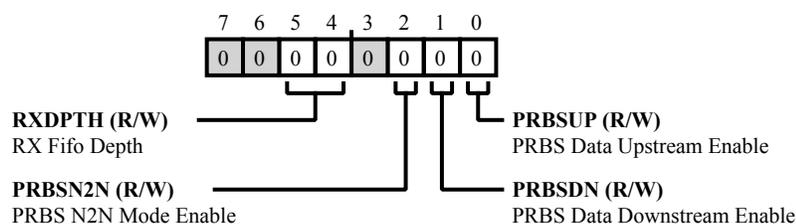


Figure 7-30: `A2B_TESTMODE` Register Diagram

Table 7-31: `A2B_TESTMODE` Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 5:4 (R/W) | RXDPATH | RX Fifo Depth. The <code>A2B_TESTMODE.RXDPATH</code> bits control the data recovery FIFO depth. |
| | | 0 Do Not Change FIFO Depth |
| | | 1 Increase FIFO Depth by 1 |
| | | 2 Increase FIFO Depth by 2 |
| 2 (R/W) | PRBSN2N | PRBS N2N Mode Enable. When the <code>A2B_TESTMODE.PRBSN2N</code> bit is set, each node checks all incoming data bits and transmits the expected data to the next node. This allows for better determination of where bus errors occur. This bit only takes effect when either or both of the <code>A2B_TESTMODE.PRBSDN</code> and <code>A2B_TESTMODE.PRBSUP</code> bits are set. |
| | | 0 Disabled |
| | | 1 Enabled |

Table 7-31: A2B_TESTMODE Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 1 (R/W) | PRBSDN | PRBS Data Downstream Enable. The A2B_TESTMODE.PRBSDN bit enables PRBS data to be sent downstream towards the last target node. |
| | | 0 Disable PRBS data |
| | | 1 PRBS data |
| 0 (R/W) | PRBSUP | PRBS Data Upstream Enable. The A2B_TESTMODE.PRBSUP bit enables PRBS data to be sent upstream towards the main node. |
| | | 0 Disable PRBS data |
| | | 1 PRBS data |

PRBS Error Count Byte 0 Register

The `A2B_ERRCNT0` register holds the least significant byte of the 32-bit error count accumulated during the PRBS bit error test.

Address: 0x21

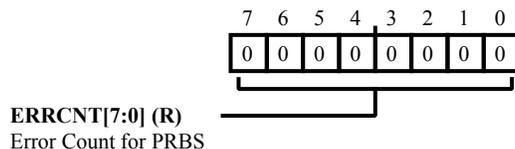


Figure 7-31: A2B_ERRCNT0 Register Diagram

Table 7-32: A2B_ERRCNT0 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7:0 (R/NW) | ERRCNT | Error Count for PRBS. The <code>A2B_ERRCNT0</code> .ERRCNT bit field contains one byte of the 32-bit PRBS bit error count. |

PRBS Error Count Byte 1 Register

The `A2B_ERRCNT1` register holds the second byte (bits 15:8) of the error count accumulated during the PRBS bit error test.

Address: 0x22

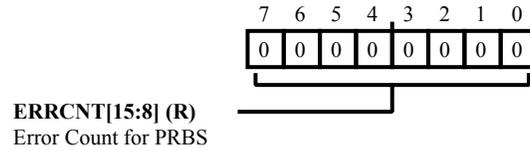


Figure 7-32: A2B_ERRCNT1 Register Diagram

Table 7-33: A2B_ERRCNT1 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 7:0 (R/NW) | ERRCNT | Error Count for PRBS. The <code>A2B_ERRCNT1 . ERRCNT</code> bit field contains one byte of the 32-bit PRBS bit error count. |

PRBS Error Count Byte 2 Register

The `A2B_ERRCNT2` register holds the third byte (bits 23:16) of the error count accumulated during the PRBS bit error test.

Address: 0x23

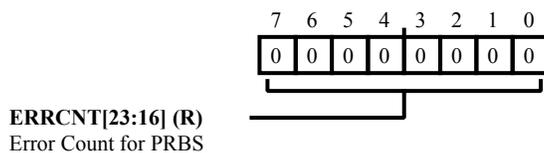


Figure 7-33: A2B_ERRCNT2 Register Diagram

Table 7-34: A2B_ERRCNT2 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7:0 (R/NW) | ERRCNT | Error Count for PRBS. The <code>A2B_ERRCNT2</code> .ERRCNT bit field contains one byte of the 32-bit PRBS bit error count. |

PRBS Error Count Byte 3 Register

The `A2B_ERRCNT3` register holds the most significant byte (bits 31:24) of the 32-bit error count accumulated during the PRBS bit error test. The `A2B_ERRCNT0` register is the least significant byte of the 32-bit error count.

Address: 0x24

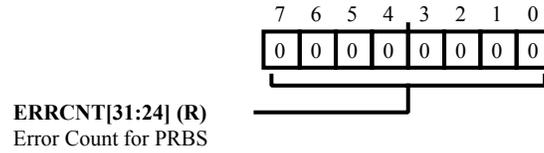


Figure 7-34: A2B_ERRCNT3 Register Diagram

Table 7-35: A2B_ERRCNT3 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7:0 (R/NW) | ERRCNT | Error Count for PRBS. The <code>A2B_ERRCNT3</code> . <code>ERRCNT</code> bit field contains one byte of the 32-bit PRBS bit error count. |

Node Register

The `A2B_NODE` register contains information required for node-to-node communication.

Address: 0x29

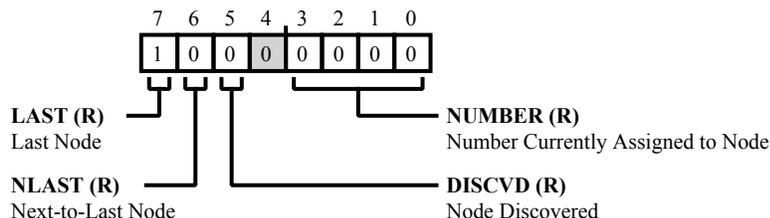


Figure 7-35: A2B_NODE Register Diagram

Table 7-36: A2B_NODE Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 7 (R/NW) | LAST | Last Node. The <code>A2B_NODE.LAST</code> bit indicates that this node is not connected to a downstream node. It is set by default at reset and cleared during discovery. |
| | | 0 Not Last Node |
| | | 1 Last Node |
| 6 (R/NW) | NLAST | Next-to-Last Node. The <code>A2B_NODE.NLAST</code> bit indicates that this node is directly upstream of the last node. It is set during discovery. |
| | | 0 Not Next-to-Last Node |
| | | 1 Next-to-Last Node |
| 5 (R/NW) | DISCVD | Node Discovered. The <code>A2B_NODE.DISCVD</code> bit indicates that this node has been discovered. This bit always reads as 0 in a main node. |
| | | 0 Not Discovered |
| | | 1 Discovered |
| 3:0 (R/NW) | NUMBER | Number Currently Assigned to Node. The <code>A2B_NODE.NUMBER</code> bit field reports the node number assigned to the node during discovery. This field always reads as 0 in a main node. |

Discovery Status Register (Main Node Only)

The `A2B_DISCSTAT` register provides status for discovery transactions on the A²B bus. An I²C write to the `A2B_DISCVRY` register sets the `A2B_DISCSTAT.DSCACT` bit and causes the `A2B_NODEADR.NODE` field to be written to this register. Discovery mode can be aborted by writing to the `A2B_CONTROL.ENDDSC` bit.

Address: 0x2B

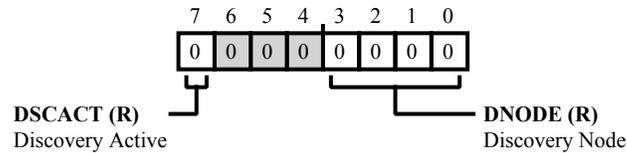


Figure 7-36: A2B_DISCSTAT Register Diagram

Table 7-37: A2B_DISCSTAT Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7 (R/NW) | DSCACT | Discovery Active. The <code>A2B_DISCSTAT.DSCACT</code> bit is set while the main node is in discovery mode. |
| 3:0 (R/NW) | DNODE | Discovery Node. When the <code>A2B_DISCSTAT.DSCACT</code> bit is set, the <code>A2B_DISCSTAT.DNODE</code> bit field shows the node being used for discovery frames. If <code>A2B_DISCSTAT.DSCACT</code> is cleared, the <code>A2B_DISCSTAT.DNODE</code> bit field retains the value of the last node discovered. |

LVDSA TX Control Register

The `A2B_TXACTL` register provides transmitter control for LVDS transceiver A. The values in this register are only applied when the `A2B_TXACTL.TXAOVREN` bit is set.

Address: 0x2E

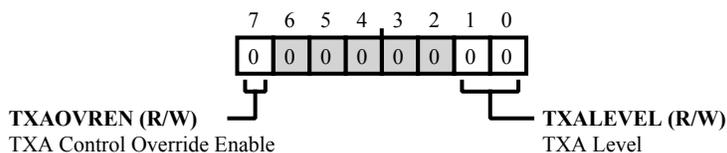


Figure 7-37: A2B_TXACTL Register Diagram

Table 7-38: A2B_TXACTL Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 7 (R/W) | TXAOVREN | TXA Control Override Enable. The <code>A2B_TXACTL.TXAOVREN</code> bit is used to force values from the <code>A2B_TXACTL</code> register to override the default values. |
| 1:0 (R/W) | TXALEVEL | TXA Level. The <code>A2B_TXACTL.TXALEVEL</code> bit field determines the transmitter output signal levels. |
| | 0 | High Transmit Power Level |
| | 1 | Reserved |
| | 2 | Medium Transmit Power Level |
| | 3 | Low Transmit Power Level |

LVDSB TX Control Register

The `A2B_TXBCTL` register provides transmitter control for LVDS transceiver B. The values in this register are only applied when the `A2B_TXBCTL.TXBOVREN` bit is set.

Address: 0x30

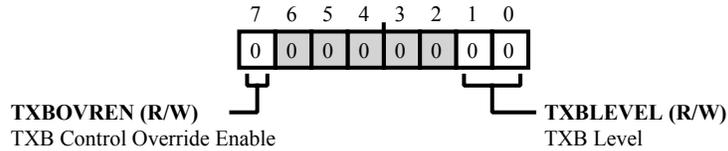


Figure 7-38: A2B_TXBCTL Register Diagram

Table 7-39: A2B_TXBCTL Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 7 (R/W) | TXBOVREN | TXB Control Override Enable. The <code>A2B_TXBCTL.TXBOVREN</code> bit is used to force values from the <code>A2B_TXBCTL</code> register to override the default values. |
| 1:0 (R/W) | TXBLEVEL | TXB Level. The <code>A2B_TXBCTL.TXBLEVEL</code> bit field determines the transmitter output signal levels. |
| | 0 | High Transmit Power Level |
| | 1 | Reserved |
| | 2 | Medium Transmit Power Level |
| | 3 | Low Transmit Power Level |

Local Interrupt Type (Subordinate Only)

The `A2B_LINTTYPE` register contains information about the pending local interrupt from a subordinate node to a local processor signaled with the IRQ pin. A read of the `A2B_LINTTYPE` register in a target node by a local processor has the effect of clearing this pending interrupt as well as deasserting the IRQ pin. A read of this register by the Host processor has no effect. This register is only used when signaling mailbox interrupts to a local processor in an A2B target.

Address: 0x3E

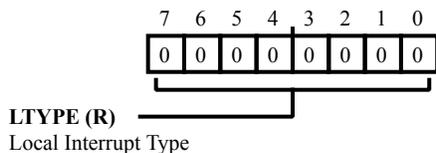


Figure 7-39: A2B_LINTTYPE Register Diagram

Table 7-40: A2B_LINTTYPE Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|-------------------------|
| 7:0 (R/NW) | LTYPE | Local Interrupt Type. |
| | | 48 Mailbox 0 Full |
| | | 49 Mailbox 0 Empty |
| | | 50 Mailbox 1 Full |
| | | 51 Mailbox 1 Empty |

I2C Configuration Register

The `A2B_I2CCFG` register controls the data rate of the I²C port in A²B target nodes and sets the I²C behavior in the A²B controller node.

Address: 0x3F

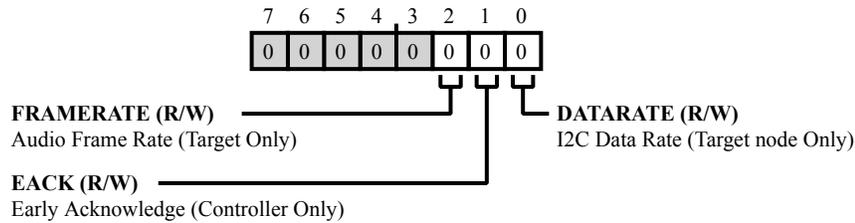


Figure 7-40: A2B_I2CCFG Register Diagram

Table 7-41: A2B_I2CCFG Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|-----------|--|
| 2 (R/W) | FRAMERATE | Audio Frame Rate (Target Only). The <code>A2B_I2CCFG.FRAMERATE</code> bit defaults to 48 kHz. This bit only affects the local clock generation for the I ² C interface to match standard I ² C clock speeds. |
| | | 0 48 kHz |
| | | 1 44.1 kHz |
| 1 (R/W) | EACK | Early Acknowledge (Controller Only). When <code>A2B_I2CCFG.EACK</code> is set, the I ² C interface provides an acknowledge to writes addressed to a target node before the write has completed on the A ² B bus. If there is an error (for example, a timeout or address error), the <code>A2B_INTPND2.I2CERR</code> bit is set. When <code>A2B_I2CCFG.EACK</code> is cleared, I ² C transactions are clock-stretched until they are complete in the system so that a correct ACK/NACK can be generated by the I ² C interface. The <code>A2B_I2CCFG.EACK</code> bit can be used for I ² C access of a target node. For accesses to peripherals connected to target nodes, the clock stretching feature is required for the I ² C interface of the host. |
| | | 0 Stretch Transactions |
| | | 1 Provide Write Acknowledge |
| 0 (R/W) | DATARATE | I2C Data Rate (Target node Only). The <code>A2B_I2CCFG.DATARATE</code> bit configures the I ² C data rate. |
| | | 0 100 kHz |
| | | 1 400 kHz |

PLL Control Register

The `A2B_PLLCTL` register provides control bits for the PLL.

Address: 0x40

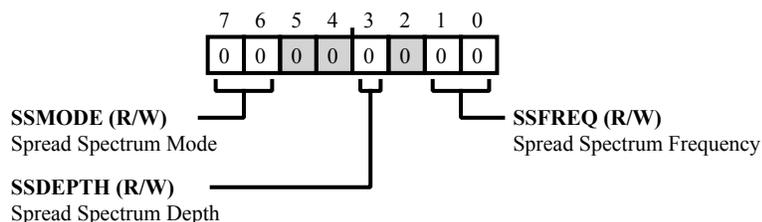


Figure 7-41: A2B_PLLCTL Register Diagram

Table 7-42: A2B_PLLCTL Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7:6 (R/W) | SSMODE | Spread Spectrum Mode. The <code>A2B_PLLCTL.SSMODE</code> bit field selects the spread spectrum mode. Spread-spectrum clocking support can be enabled for the internal clocks, the I ² S interfaces, and the programmed CLKOUTs. |
| | | 0 No Spread |
| | | 1 Spread A ² B Bus Clocks Only |
| | | 2 Spread A ² B Bus Clocks and I ² S Clocks |
| | | 3 Reserved |
| 3 (R/W) | SSDEPTH | Spread Spectrum Depth. The <code>A2B_PLLCTL.SSDEPTH</code> bit determines the spread spectrum depth of modulation. |
| | | 0 Low Spread Spectrum Depth of Modulation |
| | | 1 High Spread Spectrum Depth of Modulation |
| 1:0 (R/W) | SSFREQ | Spread Spectrum Frequency. The <code>A2B_PLLCTL.SSFREQ</code> bit determines the frequency modulation (multiples of f_{SYNCM}). |
| | | 0 4x |
| | | 1 5x |
| | | 2 6x |
| | | 3 7x |

I2S Global Configuration Register

The `A2B_I2SGCFG` register provides bits which control the operation of all I²S units. The `A2B_I2SGCFG` register must be programmed before any of the `A2B_I2SCFG.TX0EN`, `A2B_I2SCFG.TX1EN`, `A2B_I2SCFG.RX0EN`, `A2B_I2SCFG.RX1EN`, `A2B_PDMCTL.PDM0EN`, and `A2B_PDMCTL.PDM1EN` : "controllers bits are set.

For the controller node, the `A2B_I2SGCFG` register must be programmed before discovery and must not be modified after discovery.

Address: 0x41

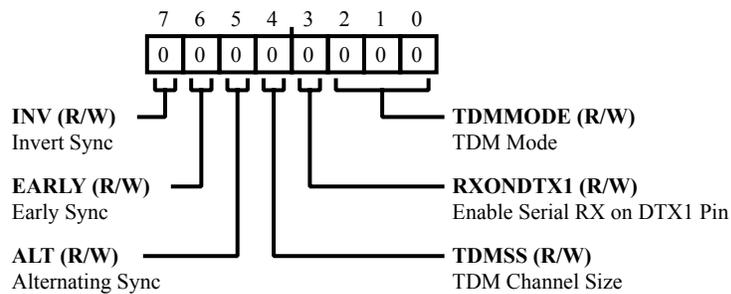


Figure 7-42: A2B_I2SGCFG Register Diagram

Table 7-43: A2B_I2SGCFG Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 7 (R/W) | INV | Invert Sync. The <code>A2B_I2SGCFG.INV</code> bit determines whether the rising edge or the falling edge of the <code>A2B_SYNC</code> pin corresponds to the start of an audio frame. If the <code>A2B_I2SGCFG.INV</code> bit is to be set in a controller node, it must be set before the <code>A2B_SWCTL.ENS</code> bit is set. |
| | | 0 Rising edge of SYNC pin at start of audio frame |
| | | 1 Falling edge of SYNC pin at start of audio frame |
| 6 (R/W) | EARLY | Early Sync. The <code>A2B_I2SGCFG.EARLY</code> bit determines whether the <code>A2B_SYNC</code> pin changes in the same cycle as the MSB of data slot 0 or one cycle before the MSB of data slot 0. |
| | | 0 Change SYNC pin in same cycle |
| | | 1 Change SYNC pin in previous cycle |

Table 7-43: A2B_I2SGCFG Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 5 (R/W) | ALT | Alternating Sync. The A2B_I2SGCFG.ALT bit determines whether the A2B_SYNC pin is pulsed high for one cycle at the start of each sampling period or driven high during right channel data and low during left channel data for I ² S stereo mode operation. |
| | | 0 Pulse SYNC Pin High for 1 Cycle |
| | | 1 Drive SYNC Pin for I ² S Operation |
| 4 (R/W) | TDMSS | TDM Channel Size. The A2B_I2SGCFG.TDMSS bit determines whether the slot size is 16 or 32 bits. |
| | | 0 32-Bit |
| | | 1 16-Bit |
| 3 (R/W) | RXONDTX1 | Enable Serial RX on DTX1 Pin. When the A2B_I2SGCFG.RXONDTX1 bit is set, the DTX1 pin is used for I ² S/TDM RX in place of the DRX1 pin, and the values of A2B_I2SCFG.TX1EN and A2B_I2SCFG.RX1EN are ignored. |
| 2:0 (R/W) | TDMMODE | TDM Mode. The A2B_I2SGCFG.TDMMODE bit field selects the mode for the I ² S/TDM units. |
| | | 0 TDM2 |
| | | 1 TDM4 |
| | | 2 TDM8 |
| | | 3 TDM12 (No target node support) |
| | | 4 TDM16 |
| | | 5 TDM20 (No target node support) |
| | | 6 TDM24 (No target node support) |
| | | 7 TDM32 (Single data pin support only) |

I2S Configuration Register

The `A2B_I2SCFG` register provides control over which I²S data pins are enabled, how the data associated with them is stored in the internal frame buffers, and the polarity of the BCLK signal.

IMPORTANT: If both `A2B_I2SCFG.RX1EN` and `A2B_I2SCFG.RX0EN` are set and both `A2B_PDMCTL.PDM1EN` and `A2B_PDMCTL.PDM0EN` are cleared, the received I²S data from the `A2B_DRX1` and `A2B_DRX0` pins is written into the A²B frame buffer independently.

IMPORTANT: If both `A2B_I2SCFG.TX1EN` and `A2B_I2SCFG.TX0EN` are set, the I²S transmit data for both the `A2B_DTX1` and `A2B_DTX0` pins is read from the A²B frame buffer independently.

Address: 0x42

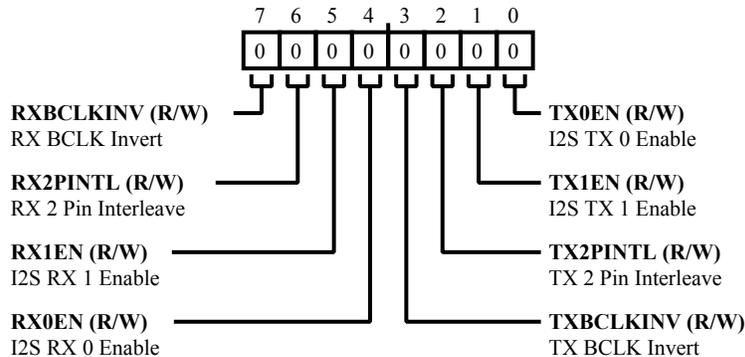


Figure 7-43: A2B_I2SCFG Register Diagram

Table 7-44: A2B_I2SCFG Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|-----------|--|
| 7 (R/W) | RXBCLKINV | RX BCLK Invert. The <code>A2B_I2SCFG.RXBCLKINV</code> bit controls the BCLK edge that the <code>A2B_DRX0</code> and <code>A2B_DRX1</code> pins are sampled on. For main nodes only, this is also the sampling edge for the <code>A2B_SYNC</code> pin. |
| | | 0 Sample on rising edge of BCLK |
| | | 1 Sample on falling edge of BCLK |

Table 7-44: A2B_I2SCFG Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|-----------|---|
| 6 (R/W) | RX2PINTL | RX 2 Pin Interleave. The A2B_I2SCFG.RX2PINTL bit is only used when TDM data is received on two pins simultaneously. When this bit is cleared (default), the data received on the A2B_DRX0 pin is associated with the lower half of the bus data slots, and the data received on the A2B_DRX1 pin is associated with the upper half of the bus data slots. When the A2B_I2SCFG.RX2PINTL bit is set, the data received on the A2B_DRX0 pin is associated with the even bus data slots (slot 0, slot 2, ..., slot 30), and the data received on the A2B_DRX1 pin is associated with the odd bus data slots (slot 1, slot 3, ..., slot 31). |
| | | 0 No Interleaving |
| | | 1 Interleaving |
| 5 (R/W) | RX1EN | I2S RX 1 Enable. The A2B_I2SCFG.RX1EN bit enables I ² S/TDM receive data on the A2B_DRX1 pin. This bit has no effect if the A2B_PDMCTL.PDM1EN bit is set. |
| | | 0 Disabled |
| | | 1 Enabled |
| 4 (R/W) | RX0EN | I2S RX 0 Enable. The A2B_I2SCFG.RX0EN bit enables I ² S/TDM receive data on the A2B_DRX0 pin. This bit has no effect if the A2B_PDMCTL.PDM0EN bit is set. |
| | | 0 Disabled |
| | | 1 Enabled |
| 3 (R/W) | TXBCLKINV | TX BCLK Invert. The A2B_I2SCFG.TXBCLKINV bit controls the BCLK edge that the A2B_DTX0 and A2B_DTX1 pins are driven on. For subordinate nodes only, this is also the driving edge for the A2B_SYNC pin. |
| | | 0 Drive on Rising Edge of BCLK |
| | | 1 Drive on Falling Edge of BCLK |

Table 7-44: A2B_I2SCFG Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 2 (R/W) | TX2PINTL | TX 2 Pin Interleave. The A2B_I2SCFG.TX2PINTL bit is only used when TDM data is transmitted on two pins simultaneously. When this bit is cleared (default), the data transmitted on the A2B_DTX0 pin is associated with the lower half of the bus data slots, and the data transmitted on the A2B_DTX1 pin is associated with the upper half of the bus data slots. When the A2B_I2SCFG.TX2PINTL bit is set, the even bus data slots (slot 0, slot 2, ..., slot 30) transmitted on the A2B_DTX0 pin, and the odd bus data slots (slot 1, slot 3, ..., slot 31) are transmitted on the A2B_DTX1 pin. |
| | | 0 Disabled |
| | | 1 Enabled |
| 1 (R/W) | TX1EN | I2S TX 1 Enable. The A2B_I2SCFG.TX1EN bit enables I ² S/TDM transmit data on the A2B_DTX1 pin. |
| | | 0 Disabled |
| | | 1 Enabled |
| 0 (R/W) | TX0EN | I2S TX 0 Enable. The A2B_I2SCFG.TX0EN bit enables I ² S/TDM transmit data on the A2B_DTX0 pin. |
| | | 0 Disabled |
| | | 1 Enabled |

I2S Rate Register (Target Only)

The `A2B_I2SRATE` register controls the I²S/TDM interfaces in target nodes, which may run at a multiple of the superframe rate.

Address: 0x43

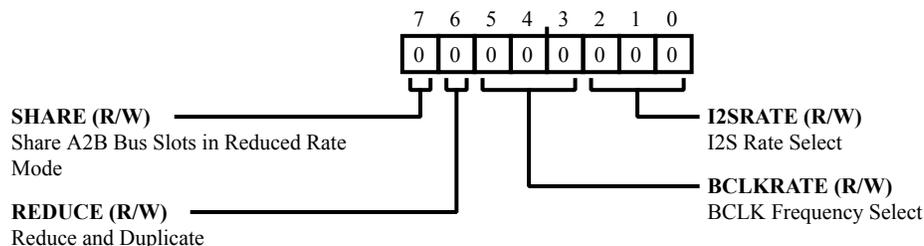


Figure 7-44: A2B_I2SRATE Register Diagram

Table 7-45: A2B_I2SRATE Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7 (R/W) | SHARE | Share A2B Bus Slots in Reduced Rate Mode. The <code>A2B_I2SRATE.SHARE</code> bit function applies only when the local sample rate is lower than the superframe rate. When the <code>A2B_I2SRATE.SHARE</code> bit is set, I ² S/TDM data for the local node is time multiplexed on the A ² B bus. Only <code>A2B_I2SRRSOFFS.RRSOFFSET</code> values of 0 or 1 are supported when the <code>A2B_I2SRATE.SHARE</code> bit is enabled. |
| | | 0 Disabled |
| | | 1 Enabled |
| 6 (R/W) | REDUCE | Reduce and Duplicate. The <code>A2B_I2SRATE.REDUCE</code> bit function applies only when the local sample rate is higher than the superframe rate. When the <code>A2B_I2SRATE.REDUCE</code> bit is set, the number of received samples is reduced so that only one sample is used per superframe, and transmitted samples are duplicated so that only one sample is needed per superframe. |
| | | 0 Disabled |
| | | 1 Enabled |
| 5:3 (R/W) | BCLKRATE | BCLK Frequency Select. The <code>A2B_I2SRATE.BCLKRATE</code> bit field is used to select an alternate BCLK frequency for a reduced rate target node. The nominal BCLK frequency is determined by the superframe frequency (SFF), settings, and reduced rate divide ratio (from <code>A2B_I2SRRATE.RRDIV</code> and <code>A2B_I2SRATE.I2SRATE</code>). |
| | | 0 BCLK frequency as configured in <code>A2B_I2SGCFG</code> |
| | | 1 SYNC frequency x 2048 |

Table 7-45: A2B_I2SRATE Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| | | 2 SYNC frequency x 4096 |
| | | 4 SFF frequency x 64 |
| | | 5 SFF frequency x 128 |
| | | 6 SFF frequency x 256 |
| 2:0 (R/W) | I2SRATE | <p>I2S Rate Select.</p> <p>The <code>A2B_I2SRATE.I2SRATE</code> bit sets the rate for I²S/TDM transmit and receive operations in the local target node. This sample rate is based on the superframe frequency (SFF is either 48 kHz or 44.1 kHz).</p> |
| | | 0 SFF x 1 |
| | | 1 SFF / 2 |
| | | 2 SFF / 4 |
| | | 3 SFF / <code>A2B_I2SRRATE.RRDIV</code> |
| | | 5 SFF x 2 |
| | | 6 SFF x 4 |

I2S Transmit Data Offset Register (Controller Only)

The `A2B_I2STXOFFSET` register controls the number of I²S transmit channels which are skipped before the node begins transmitting data.

Address: 0x44

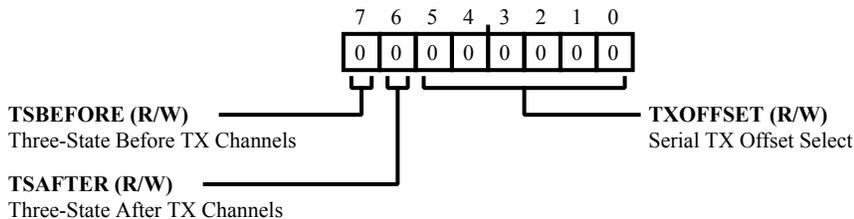


Figure 7-45: A2B_I2STXOFFSET Register Diagram

Table 7-46: A2B_I2STXOFFSET Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 7 (R/W) | TSBEFORE | Three-State Before TX Channels. When the <code>A2B_I2STXOFFSET.TSBEFORE</code> bit is cleared (default), the <code>A2B_DTX0</code> and <code>A2B_DTX1</code> pins are driven low at the beginning of each frame for the number of data channels defined in <code>A2B_I2STXOFFSET.TXOFFSET</code> . When this bit is set, the <code>A2B_DTX0</code> and <code>A2B_DTX1</code> pins are instead three-stated for the number of data channels defined in <code>A2B_I2STXOFFSET.TXOFFSET</code> . |
| | | 0 Disable |
| | | 1 Enable |
| 6 (R/W) | TSAFTER | Three-State After TX Channels. When the <code>A2B_I2STXOFFSET.TSAFTER</code> bit is cleared (default), the <code>A2B_DTX0</code> and <code>A2B_DTX1</code> pins are driven low after all valid channels have been transmitted. When the <code>A2B_I2STXOFFSET.TSAFTER</code> bit is set, the <code>A2B_DTX0</code> and <code>A2B_DTX1</code> pins are instead three-stated after all valid channels have been transmitted. |
| | | 0 Disable |
| | | 1 Enable |

Table 7-46: A2B_I2STXOFFSET Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 5:0 (R/W) | TXOFFSET | Serial TX Offset Select. The A2B_I2STXOFFSET.TXOFFSET bit field defines the number of I ² S/TDM channels that are skipped before the node begins transmitting data. The valid values for this field are 0-63. |
| | | 0 No TX offset |
| | | 1 1 TDM channel |
| | | 62 62 TDM channels |
| | | 63 63 TDM channels |

I2S Receive Data Offset Register (Controller Only)

The `A2B_I2SRXOFFSET` register controls the number of I²S receive channels which are skipped before the node begins receiving data.

Address: 0x45

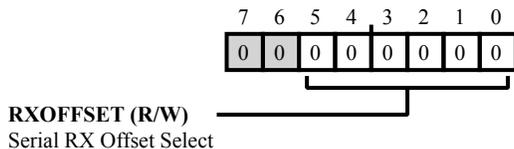


Figure 7-46: A2B_I2SRXOFFSET Register Diagram

Table 7-47: A2B_I2SRXOFFSET Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 5:0 (R/W) | RXOFFSET | Serial RX Offset Select. The <code>A2B_I2SRXOFFSET.RXOFFSET</code> bit field defines the number of I ² S/TDM channels that are skipped before the node begins receiving data. The valid values for this field are 0-63. |
| | | 0 No RX offset |
| | | 62 62 TDM channels |
| | | 63 63 TDM channels |

SYNC Offset Register (Subordinate Only)

The `A2B_SYNCOFFSET` register adjusts the A²B bus clock (f_{SYSBCLK}) cycle count on which the `A2B_SYNC` pin indicates the start of an audio frame. A²B subordinate nodes can all sample exactly at the same time by individually compensating for their propagation delay with this register setting.

The `A2B_SYNCOFFSET` register must be programmed before any of the data pin enable bits are set (`A2B_I2SCFG.TX0EN`, `A2B_I2SCFG.TX1EN`, `A2B_I2SCFG.RX0EN`, `A2B_I2SCFG.RX1EN`, `A2B_PDMCTL.PDM0EN`, or `A2B_PDMCTL.PDM1EN`).

Address: 0x46

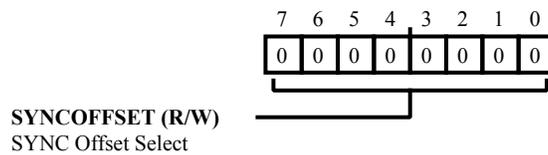


Figure 7-47: A2B_SYNCOFFSET Register Diagram

Table 7-48: A2B_SYNCOFFSET Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|------------|---|
| 7:0 (R/W) | SYNCOFFSET | SYNC Offset Select. The <code>A2B_SYNCOFFSET.SYNCOFFSET</code> bit field adjusts the system clock cycle where the <code>A2B_SYNC</code> pin indicates the start of an audio frame. The value programmed to the <code>A2B_SYNCOFFSET.SYNCOFFSET</code> field is the 8-bit signed, two's complement representation of the integer value defining the number of <code>SYSBCLK</code> cycles that lag the SYNC signal before the superframe begins. Valid values for this field range from no SYNC offset (0x00) to the SYNC occurring 127 cycles before the start of the superframe (0x81). |
| | 0 | No offset |
| | 1-128 | Reserved |
| | 129 | 127 <code>SYSBCLK</code> cycles |
| | 130-254 | 126 to 2 <code>SYSBCLK</code> cycles (respectively) |
| | 255 | 1 <code>SYSBCLK</code> cycle |

PDM Control Register

The `A2B_PDMCTL` register provides enable bits for the pulse density modulators.

Address: 0x47

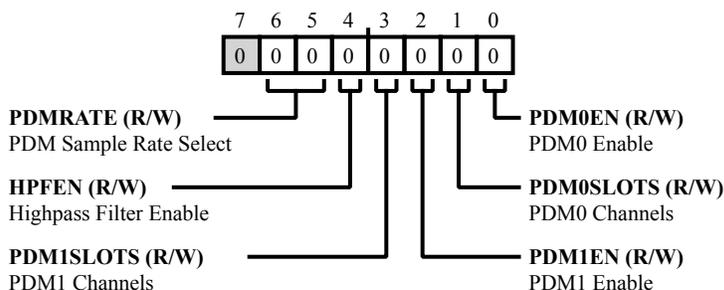


Figure 7-48: A2B_PDMCTL Register Diagram

Table 7-49: A2B_PDMCTL Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|-----------|---|
| 6:5 (R/W) | PDMRATE | PDM Sample Rate Select. The <code>A2B_PDMCTL.PDMRATE</code> bit field controls the output rate of the PDM demodulators, which is based off of the superframe rate (SFF). Changes to the <code>A2B_PDMCTL.PDMRATE</code> field do not change the PDM clock frequency. For a subordinate node, setting the node to a reduced rate changes the SYNC and PDM clock frequencies. Setting the subordinate node to an increased rate changes only the SYNC. The PDM clock frequency stays at 3.07MHz. |
| | | 0 SFF |
| | | 1 SFF/2 |
| | | 2 SFF/4 |
| 4 (R/W) | HPFEN | Highpass Filter Enable. The <code>A2B_PDMCTL.HPFEN</code> bit controls whether or not the high pass filter is used on received PDM data. |
| | | 0 Disabled |
| | | 1 Enabled |
| 3 (R/W) | PDM1SLOTS | PDM1 Channels. The <code>A2B_PDMCTL.PDM1SLOTS</code> bit controls whether the PDM signal on the <code>A2B_DRX1</code> pin is one channel (mono) or two channels (stereo). |
| | | 0 Mono |
| | | 1 Stereo |

Table 7-49: A2B_PDMCTL Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|-----------|---|
| 2 (R/W) | PDM1EN | PDM1 Enable. The A2B_PDMCTL.PDM1EN bit enables PDM reception on the A2B_DRX1/ A2B_IO6 pin. |
| | | 0 Disabled |
| | | 1 Enabled |
| 1 (R/W) | PDM0SLOTS | PDM0 Channels. The A2B_PDMCTL.PDM0SLOTS bit controls whether the PDM signal on the A2B_DRX0 pin is one channel (mono) or two channels (stereo). |
| | | 0 Mono |
| | | 1 Stereo |
| 0 (R/W) | PDM0EN | PDM0 Enable. The A2B_PDMCTL.PDM0EN bit enables PDM reception on the A2B_DRX0/ A2B_IO5 pin. |
| | | 0 Disabled |
| | | 1 Enabled |

Error Management Register

The `A2B_ERRMGMT` register provides options for reporting communication errors over the I²S/TDM interface.

Address: 0x48

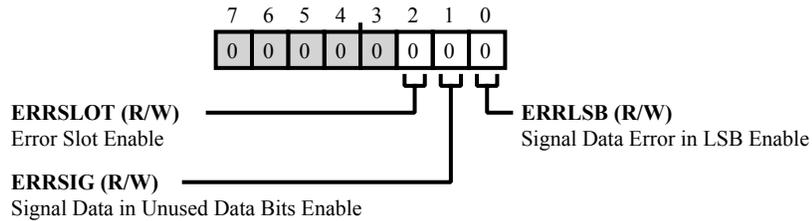


Figure 7-49: A2B_ERRMGMT Register Diagram

Table 7-50: A2B_ERRMGMT Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 2 (R/W) | ERRSLOT | Error Slot Enable. Setting the <code>A2B_ERRMGMT.ERRSLOT</code> bit causes the transceiver to append an extra I ² S/TDM data channel to the TDM stream to indicate A ² B errors in the received data slots. |
| | | 0 Disabled |
| | | 1 Enabled |
| 1 (R/W) | ERRSIG | Signal Data in Unused Data Bits Enable. When the <code>A2B_ERRMGMT.ERRSIG</code> is set, any unused data bits in each I ² S/TDM channel indicate data errors. |
| | | 0 Disabled |
| | | 1 Enabled |
| 0 (R/W) | ERRLSB | Signal Data Error in LSB Enable. When the <code>A2B_ERRMGMT.ERRLSB</code> bit is set, the LSB of each I ² S/TDM sample is replaced with an active-high status bit indicating that there is an error in the data slot (1 = error, 0 = no error). |
| | | 0 Disabled |
| | | 1 Enabled |

GPIO Output Data Register

The `A2B_GPIODAT` register controls output data for general-purpose I/O pins.

Address: 0x4A

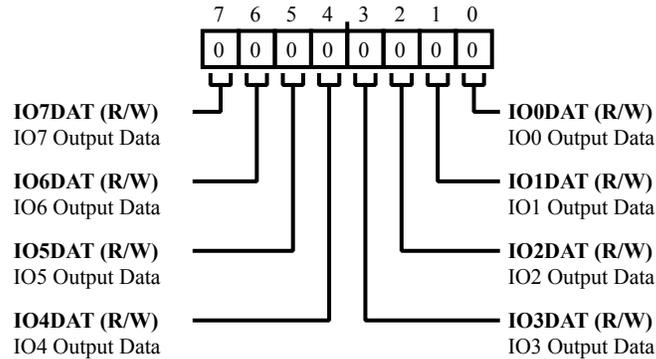


Figure 7-50: A2B_GPIODAT Register Diagram

Table 7-51: A2B_GPIODAT Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7 (R/W) | IO7DAT | IO7 Output Data. The value of the <code>A2B_GPIODAT.IO7DAT</code> bit is driven onto the IO7 pin when it is in GPIO mode with its output driver enabled (<code>A2B_GPIOOEN.IO7OEN=1</code>). |
| | | 0 Output Low |
| | | 1 Output High |
| 6 (R/W) | IO6DAT | IO6 Output Data. The value of the <code>A2B_GPIODAT.IO6DAT</code> bit is driven onto the IO6 pin when it is in GPIO mode with its output driver enabled (<code>A2B_GPIOOEN.IO6OEN=1</code>). |
| | | 0 Output Low |
| | | 1 Output High |
| 5 (R/W) | IO5DAT | IO5 Output Data. The value of the <code>A2B_GPIODAT.IO5DAT</code> bit is driven onto the IO5 pin when it is in GPIO mode with its output driver enabled (<code>A2B_GPIOOEN.IO5OEN=1</code>). |
| | | 0 Output Low |
| | | 1 Output High |

Table 7-51: A2B_GPIODAT Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 4 (R/W) | IO4DAT | IO4 Output Data. The value of the A2B_GPIODAT.IO4DAT bit is driven onto the IO4 pin when it is in GPIO mode with its output driver enabled (A2B_GPIOOEN.IO4OEN=1). |
| | | 0 Output Low |
| | | 1 Output High |
| 3 (R/W) | IO3DAT | IO3 Output Data. The value of the A2B_GPIODAT.IO3DAT bit is driven onto the IO3 pin when it is in GPIO mode with its output driver enabled (A2B_GPIOOEN.IO3OEN=1). |
| | | 0 Output Low |
| | | 1 Output High |
| 2 (R/W) | IO2DAT | IO2 Output Data. The value of the A2B_GPIODAT.IO2DAT bit is driven onto the IO2 pin when it is in GPIO mode with its output driver enabled (A2B_GPIOOEN.IO2OEN=1). |
| | | 0 Output Low |
| | | 1 Output High |
| 1 (R/W) | IO1DAT | IO1 Output Data. The value of the A2B_GPIODAT.IO1DAT bit is driven onto the IO1 pin when it is in GPIO mode with its output driver enabled (A2B_GPIOOEN.IO1OEN=1). |
| | | 0 Output Low |
| | | 1 Output High |
| 0 (R/W) | IO0DAT | IO0 Output Data. The value of the A2B_GPIODAT.IO0DAT bit is driven onto the IO0 pin when it is in GPIO mode with its output driver enabled (A2B_GPIOOEN.IO0OEN=1). |
| | | 0 Output Low |
| | | 1 Output High |

GPIO Output Data Set Register

The `A2B_GPIODATSET` register allows setting of individual GPIO output register bits (write 1 to set) without influencing the states of the other GPIO output register bits. Reads from this address return the value in the GPIO output data (`A2B_GPIODAT`) register.

Address: 0x4B

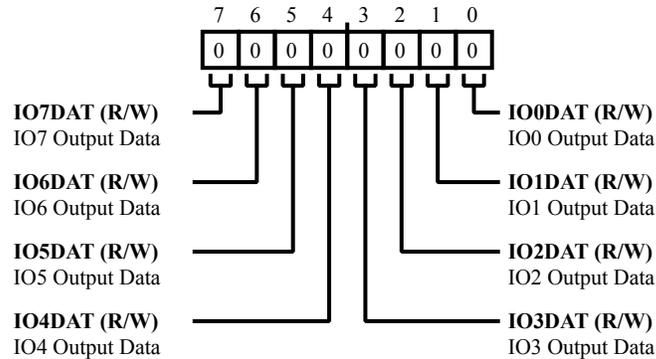


Figure 7-51: A2B_GPIODATSET Register Diagram

Table 7-52: A2B_GPIODATSET Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 7 (R/W1S) | IO7DSET | IO7 Data Set. The <code>A2B_GPIODATSET.IO7DSET</code> bit executes a write-1-to-set action for the <code>A2B_GPIODAT.IO7DAT</code> bit. |
| | | 0 No Action |
| | | 1 Set Bit |
| 6 (R/W1S) | IO6DSET | IO6 Data Set. The <code>A2B_GPIODATSET.IO6DSET</code> bit executes a write-1-to-set action for the <code>A2B_GPIODAT.IO6DAT</code> bit. |
| | | 0 No Action |
| | | 1 Set Bit |
| 5 (R/W1S) | IO5DSET | IO5 Data Set. The <code>A2B_GPIODATSET.IO5DSET</code> bit executes a write-1-to-set action for the <code>A2B_GPIODAT.IO5DAT</code> bit. |
| | | 0 No Action |
| | | 1 Set Bit |

Table 7-52: A2B_GPIODATSET Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 4 (R/W1S) | IO4DSET | IO4 Data Set. The A2B_GPIODATSET.IO4DSET bit executes a write-1-to-set action for the A2B_GPIODAT.IO4DAT bit. |
| | | 0 No Action |
| | | 1 Set Bit |
| 3 (R/W1S) | IO3DSET | IO3 Data Set. The A2B_GPIODATSET.IO3DSET bit executes a write-1-to-set action for the A2B_GPIODAT.IO3DAT bit. |
| | | 0 No Action |
| | | 1 Set Bit |
| 2 (R/W1S) | IO2DSET | IO2 Data Set. The A2B_GPIODATSET.IO2DSET bit executes a write-1-to-set action for the A2B_GPIODAT.IO2DAT bit. |
| | | 0 No Action |
| | | 1 Set Bit |
| 1 (R/W1S) | IO1DSET | IO1 Data Set. The A2B_GPIODATSET.IO1DSET bit executes a write-1-to-set action for the A2B_GPIODAT.IO1DAT bit. |
| | | 0 No Action |
| | | 1 Set Bit |
| 0 (R/W1S) | IO0DSET | IO0 Data Set. The A2B_GPIODATSET.IO0DSET bit executes a write-1-to-set action for the A2B_GPIODAT.IO0DAT bit. |
| | | 0 No Action |
| | | 1 Set Bit |

GPIO Output Data Clear Register

The `A2B_GPIODATCLR` register allows clearing of individual GPIO output register bits to 0 (write 1 to clear) without influencing the states of the other GPIO output register bits. Reads from this address return the value in the GPIO output data (`A2B_GPIODAT`) register.

Address: 0x4C

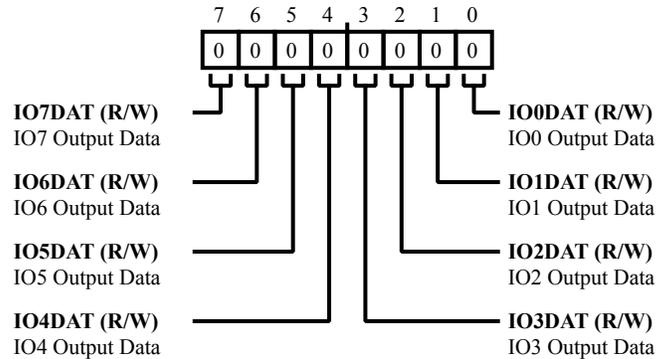


Figure 7-52: A2B_GPIODATCLR Register Diagram

Table 7-53: A2B_GPIODATCLR Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 7 (R/W1C) | IO7DCLR | IO7 Data Clear. The <code>A2B_GPIODATCLR</code> . <code>IO7DCLR</code> bit executes a write-1-to-clear action for the <code>A2B_GPIODAT</code> . <code>IO7DAT</code> bit. |
| | | 0 No Action |
| | | 1 Clear Bit |
| 6 (R/W1C) | IO6DCLR | IO6 Data Clear. The <code>A2B_GPIODATCLR</code> . <code>IO6DCLR</code> bit executes a write-1-to-clear action for the <code>A2B_GPIODAT</code> . <code>IO6DAT</code> bit. |
| | | 0 No Action |
| | | 1 Clear Bit |
| 5 (R/W1C) | IO5DCLR | IO5 Data Clear. The <code>A2B_GPIODATCLR</code> . <code>IO5DCLR</code> bit executes a write-1-to-clear action for the <code>A2B_GPIODAT</code> . <code>IO5DAT</code> bit. |
| | | 0 No Action |
| | | 1 Clear Bit |

Table 7-53: A2B_GPIODATCLR Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 4 (R/W1C) | IO4DCLR | IO4 Data Clear. The A2B_GPIODATCLR.IO4DCLR bit executes a write-1-to-clear action for the A2B_GPIODAT.IO4DAT bit. |
| | | 0 No Action |
| | | 1 Clear Bit |
| 3 (R/W1C) | IO3DCLR | IO3 Data Clear. The A2B_GPIODATCLR.IO3DCLR bit executes a write-1-to-clear action for the A2B_GPIODAT.IO3DAT bit. |
| | | 0 No Action |
| | | 1 Clear Bit |
| 2 (R/W1C) | IO2DCLR | IO2 Data Clear. The A2B_GPIODATCLR.IO2DCLR bit executes a write-1-to-clear action for the A2B_GPIODAT.IO2DAT bit. |
| | | 0 No Action |
| | | 1 Clear Bit |
| 1 (R/W1C) | IO1DCLR | IO1 Data Clear. The A2B_GPIODATCLR.IO1DCLR bit executes a write-1-to-clear action for the A2B_GPIODAT.IO1DAT bit. |
| | | 0 No Action |
| | | 1 Clear Bit |
| 0 (R/W1C) | IO0DCLR | IO0 Data Clear. The A2B_GPIODATCLR.IO0DCLR bit executes a write-1-to-clear action for the A2B_GPIODAT.IO0DAT bit. Only subordinate nodes can output data on this pin. |
| | | 0 No Action |
| | | 1 Clear Bit |

GPIO Output Enable Register

The `A2B_GPIOOEN` register controls the output enables of the general-purpose I/O pins.

Address: 0x4D

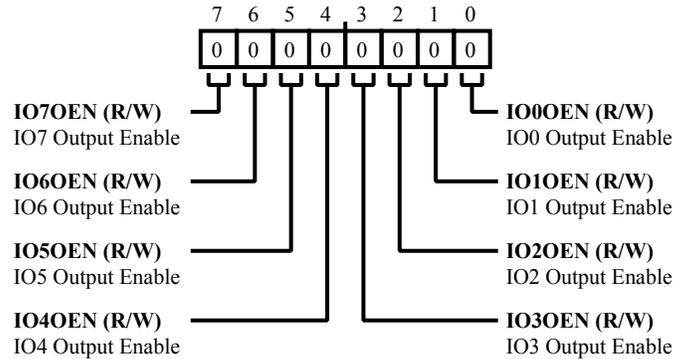


Figure 7-53: A2B_GPIOOEN Register Diagram

Table 7-54: A2B_GPIOOEN Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 7 (R/W) | IO7OEN | IO7 Output Enable. The <code>A2B_GPIOOEN.IO7OEN</code> bit configures the IO7 pin as an output when the pin is in GPIO mode. |
| | | 0 Disable |
| | | 1 Enable |
| 6 (R/W) | IO6OEN | IO6 Output Enable. The <code>A2B_GPIOOEN.IO6OEN</code> bit configures the DRX1/IO6 pin as an output when the pin is in GPIO mode. |
| | | 0 Disable |
| | | 1 Enable |
| 5 (R/W) | IO5OEN | IO5 Output Enable. The <code>A2B_GPIOOEN.IO5OEN</code> bit configures the DRX0/IO5 pin as an output when the pin is in GPIO mode. |
| | | 0 Disable |
| | | 1 Enable |

Table 7-54: A2B_GPIOOEN Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 4 (R/W) | IO4OEN | IO4 Output Enable. The A2B_GPIOOEN.IO4OEN bit configures the DTX1/IO4 pin as an output when the pin is in GPIO mode. |
| | | 0 Disable |
| | | 1 Enable |
| 3 (R/W) | IO3OEN | IO3 Output Enable. The A2B_GPIOOEN.IO3OEN bit configures the DTX0/IO3 pin as an output when the pin is in GPIO mode. |
| | | 0 Disable |
| | | 1 Enable |
| 2 (R/W) | IO2OEN | IO2 Output Enable. The A2B_GPIOOEN.IO2OEN bit configures the ADR2/IO2 pin as an output when the pin is in GPIO mode. |
| | | 0 Disable |
| | | 1 Enable |
| 1 (R/W) | IO1OEN | IO1 Output Enable. The A2B_GPIOOEN.IO1OEN bit configures the ADR1/IO1 pin as an output when the pin is in GPIO mode. |
| | | 0 Disable |
| | | 1 Enable |
| 0 (R/W) | IO0OEN | IO0 Output Enable. The A2B_GPIOOEN.IO0OEN bit configures the IRQ/IO0 pin as an output when the pin is in GPIO mode. The A2B_GPIOOEN.IO0OEN bit has no effect in a main node. Only subordinate nodes can output data on this pin. |
| | | 0 Disable |
| | | 1 Enable |

GPIO Input Enable Register

The `A2B_GPIOIEN` register controls the input enables of the general purpose I/O pins.

Address: 0x4E

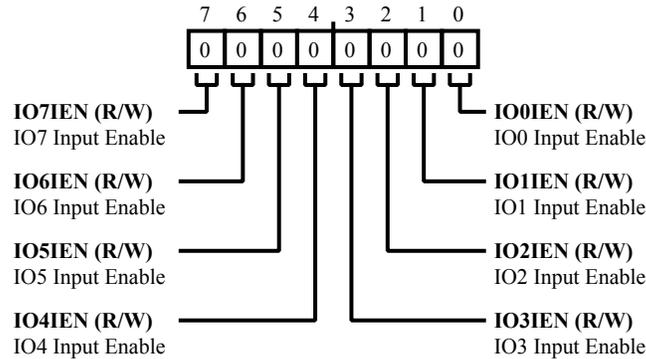


Figure 7-54: A2B_GPIOIEN Register Diagram

Table 7-55: A2B_GPIOIEN Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 7 (R/W) | IO7IEN | IO7 Input Enable. The <code>A2B_GPIOIEN.IO7IEN</code> bit is the input enable for the IO7 pin. |
| | | 0 Disable |
| | | 1 Enable |
| 6 (R/W) | IO6IEN | IO6 Input Enable. The <code>A2B_GPIOIEN.IO6IEN</code> bit is the input enable for the DRX1/IO6 pin. |
| | | 0 Disable |
| | | 1 Enable |
| 5 (R/W) | IO5IEN | IO5 Input Enable. The <code>A2B_GPIOIEN.IO5IEN</code> bit is the input enable for the DRX0/IO5 pin. |
| | | 0 Disable |
| | | 1 Enable |
| 4 (R/W) | IO4IEN | IO4 Input Enable. The <code>A2B_GPIOIEN.IO4IEN</code> bit is the input enable for the DTX1/IO4 pin. |
| | | 0 Disable |
| | | 1 Enable |

Table 7-55: A2B_GPIOIEN Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 3 (R/W) | IO3IEN | IO3 Input Enable. The A2B_GPIOIEN . IO3IEN bit is the input enable for the DTX0/IO3 pin. |
| | | 0 Disable |
| | | 1 Enable |
| 2 (R/W) | IO2IEN | IO2 Input Enable. The A2B_GPIOIEN . IO2IEN bit is the input enable for the ADR2/IO2 pin. |
| | | 0 Disable |
| | | 1 Enable |
| 1 (R/W) | IO1IEN | IO1 Input Enable. The A2B_GPIOIEN . IO1IEN bit is the input enable for the ADR1/IO1 pin. |
| | | 0 Disable |
| | | 1 Enable |
| 0 (R/W) | IO0IEN | IO0 Input Enable. The A2B_GPIOIEN . IO0IEN bit is the input enable for the IRQ/IO0 pin. This bit has no effect in a main node. |
| | | 0 Disable |
| | | 1 Enable |

GPIO Input Value Register

The `A2B_GPIOIN` register returns the value of enabled general-purpose I/O input pins.

Address: 0x4F

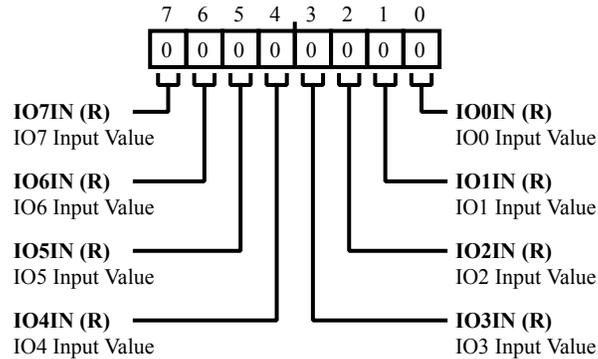


Figure 7-55: A2B_GPIOIN Register Diagram

Table 7-56: A2B_GPIOIN Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 7 (R/NW) | IO7IN | IO7 Input Value. The <code>A2B_GPIOIN.IO7IN</code> bit contains the value of the IO7 pin when in input GPIO mode (<code>A2B_GPIOIEN.IO7IEN=1</code>). Otherwise, the bit is low. |
| 6 (R/NW) | IO6IN | IO6 Input Value. The <code>A2B_GPIOIN.IO6IN</code> bit contains the value of the DRX1/IO6 pin when in input GPIO mode (<code>A2B_GPIOIEN.IO6IEN=1</code>). Otherwise, the bit is low. |
| 5 (R/NW) | IO5IN | IO5 Input Value. The <code>A2B_GPIOIN.IO5IN</code> bit contains the value of the DRX0/IO5 pin when in input GPIO mode (<code>A2B_GPIOIEN.IO5IEN=1</code>). Otherwise, the bit is low. |
| 4 (R/NW) | IO4IN | IO4 Input Value. The <code>A2B_GPIOIN.IO4IN</code> bit contains the value of the DTX1/IO4 pin when in input GPIO mode (<code>A2B_GPIOIEN.IO4IEN=1</code>). Otherwise, the bit is low. |
| 3 (R/NW) | IO3IN | IO3 Input Value. The <code>A2B_GPIOIN.IO3IN</code> bit contains the value of the DTX0/IO3 pin when in input GPIO mode (<code>A2B_GPIOIEN.IO3IEN=1</code>). Otherwise, the bit is low. |
| 2 (R/NW) | IO2IN | IO2 Input Value. The <code>A2B_GPIOIN.IO2IN</code> bit contains the value of the ADR2/IO2 pin when in input GPIO mode (<code>A2B_GPIOIEN.IO2IEN=1</code>). Otherwise, the bit is low. |

Table 7-56: A2B_GPIOIN Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 1 (R/NW) | IO1IN | IO1 Input Value. The A2B_GPIOIN.IO1IN bit contains the value of the ADR1/IO1 pin when the A2B_GPIOIEN.IO1IEN bit is high. Otherwise the bit is low. |
| 0 (R/NW) | IO0IN | IO0 Input Value. The A2B_GPIOIN.IO0IN bit contains the value of the IRQ/IO0 pin when in input GPIO mode (A2B_GPIOIEN.IO0IEN=1). Otherwise, the bit is low. This bit is only relevant in subordinate nodes and always reads 0 in a main node. |

Pin Interrupt Enable Register

The `A2B_PINTEN` register enables input-enabled GPIO pins to generate an interrupt.

Address: 0x50

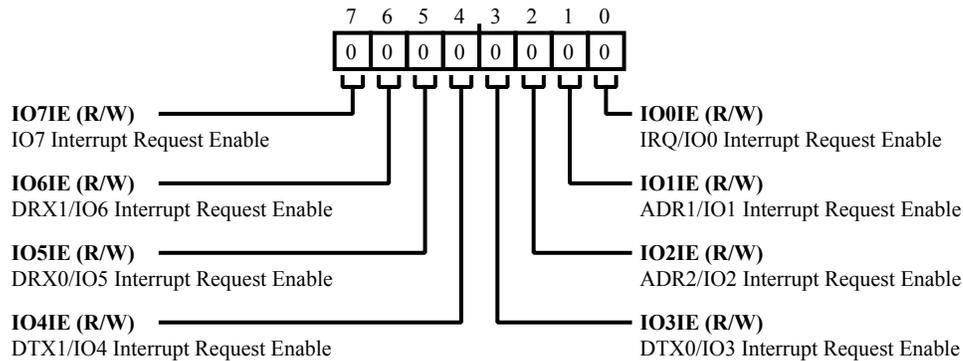


Figure 7-56: A2B_PINTEN Register Diagram

Table 7-57: A2B_PINTEN Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7 (R/W) | IO7IE | IO7 Interrupt Request Enable. The <code>A2B_PINTEN.IO7IE</code> bit enables the IO7 input to generate an interrupt request when a rising edge is sensed. |
| | | 0 Disabled |
| | | 1 Enabled |
| 6 (R/W) | IO6IE | DRX1/IO6 Interrupt Request Enable. The <code>A2B_PINTEN.IO6IE</code> bit enables the DRX1/IO6 input to generate an interrupt request when a rising edge is sensed. |
| | | 0 Disabled |
| | | 1 Enabled |
| 5 (R/W) | IO5IE | DRX0/IO5 Interrupt Request Enable. The <code>A2B_PINTEN.IO5IE</code> bit enables the DRX0/IO5 input to generate an interrupt request when a rising edge is sensed. |
| | | 0 Disabled |
| | | 1 Enabled |

Table 7-57: A2B_PINTEN Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 4 (R/W) | IO4IE | DTX1/IO4 Interrupt Request Enable. The A2B_PINTEN.IO4IE bit enables the DTX1/IO4 input to generate an interrupt request when a rising edge is sensed. |
| | | 0 Disabled |
| | | 1 Enabled |
| 3 (R/W) | IO3IE | DTX0/IO3 Interrupt Request Enable. The A2B_PINTEN.IO3IE bit enables the DTX0/IO3 input to generate an interrupt request when a rising edge is sensed. |
| | | 0 Disabled |
| | | 1 Enabled |
| 2 (R/W) | IO2IE | ADR2/IO2 Interrupt Request Enable. The A2B_PINTEN.IO2IE bit enables the ADR2/IO2 input to generate an interrupt request when a rising edge is sensed. |
| | | 0 Disabled |
| | | 1 Enabled |
| 1 (R/W) | IO1IE | ADR1/IO1 Interrupt Request Enable. The A2B_PINTEN.IO1IE bit enables bit enables the IO1 input to generate an interrupt request when a rising edge is sensed. |
| | | 0 Disabled |
| | | 1 Enabled |
| 0 (R/W) | IO0IE | IRQ/IO0 Interrupt Request Enable. The A2B_PINTEN.IO0IE bit enables the IO0 input to generate an interrupt request when a rising edge is sensed. This bit has no effect in a main node. |
| | | 0 Disabled |
| | | 1 Enabled |

Pin Interrupt Invert Register

The `A2B_PINTINV` register is used to invert pin inputs in the path to interrupt generation.

Address: 0x51

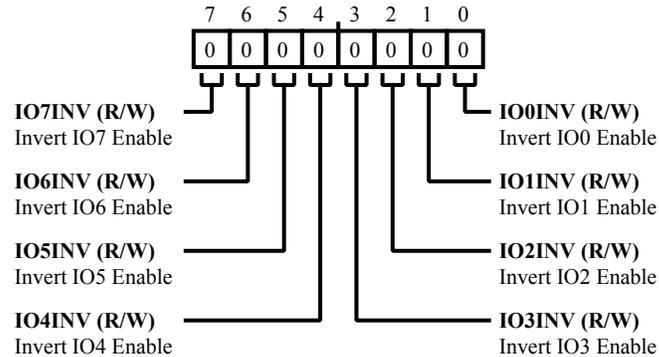


Figure 7-57: A2B_PINTINV Register Diagram

Table 7-58: A2B_PINTINV Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 7 (R/W) | IO7INV | Invert IO7 Enable. Setting the <code>A2B_PINTINV.IO7INV</code> bit inverts the polarity of the IO7 pin interrupt request input such that a falling edge sensed on the pin generates the interrupt rather than the rising edge (default). |
| | | 0 Disabled |
| | | 1 Enabled |
| 6 (R/W) | IO6INV | Invert IO6 Enable. Setting the <code>A2B_PINTINV.IO6INV</code> bit inverts the polarity of the DRX1/IO6 pin interrupt request input such that a falling edge sensed on the pin generates the interrupt rather than the rising edge (default). |
| | | 0 Disabled |
| | | 1 Enabled |
| 5 (R/W) | IO5INV | Invert IO5 Enable. Setting the <code>A2B_PINTINV.IO5INV</code> bit inverts the polarity of the DRX0/IO5 pin interrupt request input such that a falling edge sensed on the pin generates the interrupt rather than the rising edge (default). |
| | | 0 Disabled |
| | | 1 Enabled |

Table 7-58: A2B_PINTINV Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 4 (R/W) | IO4INV | Invert IO4 Enable. Setting the A2B_PINTINV.IO4INV bit inverts the polarity of the DTX1/IO4 pin interrupt request input such that a falling edge sensed on the pin generates the interrupt rather than the rising edge (default). |
| | | 0 Disabled |
| | | 1 Enabled |
| 3 (R/W) | IO3INV | Invert IO3 Enable. Setting the A2B_PINTINV.IO3INV bit inverts the polarity of the DTX0/IO3 pin interrupt request input such that a falling edge sensed on the pin generates the interrupt rather than the rising edge (default). |
| | | 0 Disabled |
| | | 1 Enabled |
| 2 (R/W) | IO2INV | Invert IO2 Enable. Setting the A2B_PINTINV.IO2INV bit inverts the polarity of the ADR2/IO2 pin interrupt request input such that a falling edge sensed on the pin generates the interrupt rather than the rising edge (default). |
| | | 0 Disabled |
| | | 1 Enabled |
| 1 (R/W) | IO1INV | Invert IO1 Enable. Setting the A2B_PINTINV.IO1INV bit inverts the polarity of the ADR1/IO1 pin interrupt request input such that a falling edge sensed on the pin generates the interrupt rather than the rising edge (default). |
| | | 0 Disabled |
| | | 1 Enabled |
| 0 (R/W) | IO0INV | Invert IO0 Enable. Setting the A2B_PINTINV.IO0INV bit inverts the polarity of the IRQ/IO0 pin interrupt request input such that a falling edge sensed on the pin generates the interrupt rather than the rising edge (default). This bit has no effect in a main node. |
| | | 0 Disabled |
| | | 1 Enabled |

Pin Configuration Register

The `A2B_PINCFG` register configures various digital pin characteristics.

Address: 0x52

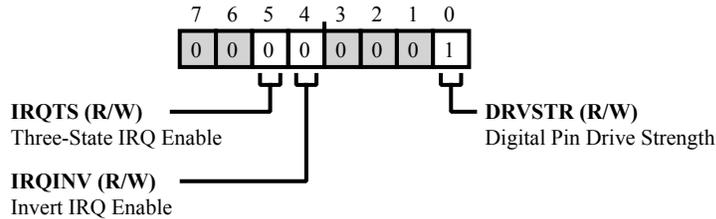


Figure 7-58: A2B_PINCFG Register Diagram

Table 7-59: A2B_PINCFG Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 5 (R/W) | IRQTS | Three-State IRQ Enable. When the <code>A2B_PINCFG.IRQTS</code> bit is cleared (default), the IRQ pin is always actively driven. Setting the <code>A2B_PINCFG.IRQTS</code> bit causes the transceiver to drive the IRQ pin when the interrupt is active and to three-state the IRQ pin when inactive. |
| | | 0 Disabled |
| | | 1 Enabled |
| 4 (R/W) | IRQINV | Invert IRQ Enable. When the <code>A2B_PINCFG.IRQINV</code> bit is cleared (default), the IRQ pin is active high. Setting the <code>A2B_PINCFG.IRQINV</code> bit makes the IRQ pin active low. |
| | | 0 Disabled |
| | | 1 Enabled |
| 0 (R/W) | DRVSTR | Digital Pin Drive Strength. The <code>A2B_PINCFG.DRVSTR</code> bit controls the drive strength of non-I ² C digital output pins. The <code>A2B_SCL</code> and <code>A2B_SDA</code> pins always have a high drive strength. |
| | | 0 Low Drive Strength |
| | | 1 High Drive Strength |

I2S Test Register

The `A2B_I2STEST` register enables a test mode to verify and debug the I²S/TDM interface.

Address: 0x53

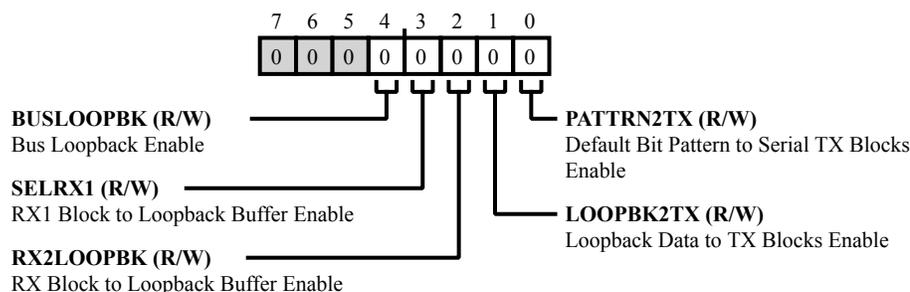


Figure 7-59: A2B_I2STEST Register Diagram

Table 7-60: A2B_I2STEST Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|-----------|---|
| 4 (R/W) | BUSLOOPBK | Bus Loopback Enable. The <code>A2B_I2STEST.BUSLOOPBK</code> bit enables data loop back from the <code>A2B_DTX0</code> pin to the <code>A2B_DRX0</code> pin and from the <code>A2B_DTX1</code> pin to the <code>A2B_DRX1</code> pin. The <code>A2B_I2STEST.LOOPBK2TX</code> , <code>A2B_I2STEST.RX2LOOPBK</code> , and <code>A2B_I2STEST.SELRX1</code> are ignored when this bit is set. |
| | | 0 Disabled |
| | | 1 Enabled |
| 3 (R/W) | SELRX1 | RX1 Block to Loopback Buffer Enable. When the <code>A2B_I2STEST.SELRX1</code> bit is cleared (default), the <code>RX0</code> block is used for the loopback test when the <code>A2B_I2STEST.RX2LOOPBK</code> bit is set. When the <code>A2B_I2STEST.SELRX1</code> bit is set, data from the <code>DRX1</code> block is used instead. |
| | | 0 Disabled |
| | | 1 Enabled |
| 2 (R/W) | RX2LOOPBK | RX Block to Loopback Buffer Enable. When the <code>A2B_I2STEST.RX2LOOPBK</code> bit is set, the receive bit pattern on either the <code>A2B_DRX0</code> or <code>A2B_DRX1</code> pins (as controlled by the <code>A2B_I2STEST.SELRX1</code> bit) is stored in the TX frame buffer. The <code>A2B_I2SCFG.RX0EN</code> , <code>A2B_I2SCFG.RX1EN</code> , and <code>A2B_I2SCFG.RX2PINTL</code> bits are ignored when this bit is set. |
| | | 0 Disabled |
| | | 1 Enabled |

Table 7-60: A2B_I2STEST Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|-----------|---|
| 1 (R/W) | LOOPBK2TX | <p>Loopback Data to TX Blocks Enable.</p> <p>When the A2B_I2STEST.LOOPBK2TX bit is set, data received on the A2B_DRX0 or A2B_DRX1 pin (as controlled by the A2B_I2STEST.SELRX1 bit) is sent out on the A2B_DTX0 and A2B_DTX1 pins.</p> <p>If the A2B_I2STEST.RX2LOOPBK bit is not set when the A2B_I2STEST.LOOPBK2TX bit is set, the default bit pattern is sent in all channels. If the A2B_I2STEST.RX2LOOPBK bit is cleared while the A2B_I2STEST.LOOPBK2TX bit is set, the last received frame is repeated. The A2B_I2SCFG.TX0EN, A2B_I2SCFG.TX1EN, and A2B_I2SCFG.TX2PINTL bits are ignored when this bit is set.</p> |
| | | 0 Disabled |
| | | 1 Enabled |
| 0 (R/W) | PATTRN2TX | <p>Default Bit Pattern to Serial TX Blocks Enable.</p> <p>When the A2B_I2STEST.PATTRN2TX bit is set, a default bit pattern (up to 32 bits) is sent in all channels on the A2B_DTX0 and A2B_DTX1 pins. The A2B_I2SCFG.TX0EN, A2B_I2SCFG.TX1EN, and A2B_I2SCFG.TX2PINTL bits are ignored when this bit is set. This bit is ignored when the A2B_I2STEST.LOOPBK2TX bit is set.</p> |
| | | 0 Disabled |
| | | 1 Enabled |

Raise Interrupt Register

The `A2B_RAISE` register allows the host to generate an interrupt in any node in the system through software. This register must be written over the A²B bus, as writes to this register from the local I²C port have no effect.

Address: 0x54

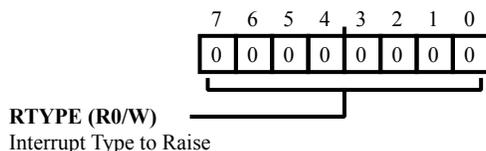


Figure 7-60: A2B_RAISE Register Diagram

Table 7-61: A2B_RAISE Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7:0 (R0/W) | RTYPE | <p>Interrupt Type to Raise.</p> <p>The <code>A2B_RAISE.RTYPE</code> bit field is written with the type of interrupt to raise. Any valid interrupt type may be generated in any node in the system. If the <code>RTYPE</code> field does not match a valid interrupt type for the node being written, no action will be taken.</p> |
| | | 0 HDCNTERR |
| | | 1 DDERR |
| | | 2 CRCERR |
| | | 3 DPERR |
| | | 4 BECOVF |
| | | 5 SRFERR |
| | | 6 SRFCRCERR |
| | | 9 PWRERR - Positive Terminal Shorted to Ground |
| | | 10 PWRERR - Negative Terminal Shorted to VBat |
| | | 11 PWRERR - Short of Wires |
| | | 12 PWRERR - Cable Disconnected or Open Circuit or Wrong Port |
| | | 13 PWRERR - Cable is Reverse Connected or Wrong Port |
| | | 15 PWRERR - Indeterminate Fault |
| | | 16 IO0PND - Subordinate node Only |
| | | 17 IO1PND |
| | | 18 IO2PND |

Table 7-61: A2B_RAISE Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| | | 19 IO3PND |
| | | 20 IO4PND |
| | | 21 IO5PND |
| | | 22 IO6PND |
| | | 23 IO7PND |
| | | 24 DSCDONE - Main Node Only |
| | | 25 I2CERR - Main Node Only |
| | | 26 ICRCERR - Main Node Only |
| | | 41 PWRERR - Non-Localized Short to Ground |
| | | 42 PWRERR - Non-Localized Short to VBat |
| | | 253 Subordinate INTTYPE Read Error - Main Node Only |
| | | 254 Standby Done - Main Node Only |
| | | 255 MSTR_RUNNING - Main Node Only |

Generate Bus Error

The `A2B_GENERR` register allows the host to generate bus errors from any node in the system through software. This register must be written over the A²B bus, as writes to this register from the local I²C port have no effect.

Address: 0x55

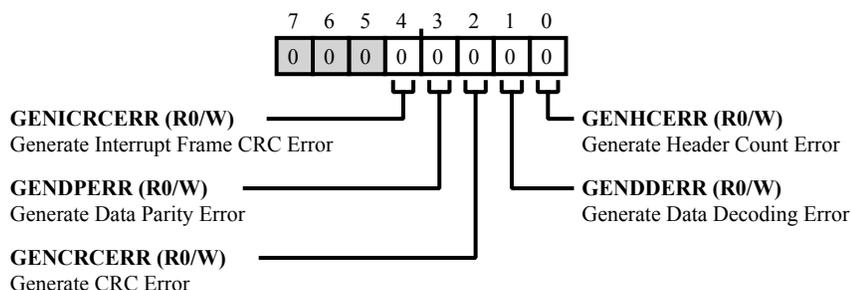


Figure 7-61: A2B_GENERR Register Diagram

Table 7-62: A2B_GENERR Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|------------|---|
| 4 (R0/W) | GENICRCERR | Generate Interrupt Frame CRC Error. A write of 1 to the <code>A2B_GENERR.GENICRCERR</code> bit instructs a subordinate node to generate an interrupt frame CRC error on the A ² B bus. A write of 1 to this bit in a main node has no effect. |
| | | 0 No Action |
| | | 1 Generate Error |
| 3 (R0/W) | GENDPERR | Generate Data Parity Error. A write of 1 to the <code>A2B_GENERR.GENDPERR</code> bit instructs the node to generate a data parity error on the A ² B bus. |
| | | 0 No Action |
| | | 1 Generate Error |
| 2 (R0/W) | GENCRCERR | Generate CRC Error. A write of 1 to the <code>A2B_GENERR.GENCRCERR</code> bit instructs the node to generate a CRC error on the A ² B bus. |
| | | 0 No Action |
| | | 1 Generate Error |

Table 7-62: A2B_GENERR Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 1 (R0/W) | GENDDERR | Generate Data Decoding Error. A write of 1 to the A2B_GENERR.GENDDERR bit instructs the node to generate a Data Decode Error on the A ² B bus. |
| | | 0 No Action |
| | | 1 Generate Error |
| 0 (R0/W) | GENHCERR | Generate Header Count Error. A write of 1 to the A2B_GENERR.GENHCERR bit instructs the node to generate a header count error on the A ² B bus. |
| | | 0 No Action |
| | | 1 Generate Error |

I2S Reduced Rate Register (Controller Node Only, Auto-Broadcast)

The `A2B_I2SRRATE` register provides a means of reducing the A²B bus data rate by delivering data on a subset of superframes rather than on each superframe, thereby reducing the overall bus power.

When the `A2B_I2SRRATE` register is written in the controller node, the new setting is automatically broadcast over the A²B bus to all discovered target nodes. Local host writes to this register in a target node have no effect.

Address: 0x56

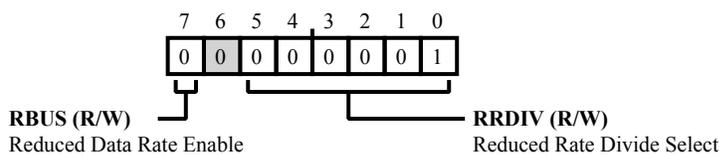


Figure 7-62: A2B_I2SRRATE Register Diagram

Table 7-63: A2B_I2SRRATE Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 7 (R/W) | RBUS | Reduced Data Rate Enable. When the <code>A2B_I2SRRATE.RBUS</code> bit is set, the bus is configured for reduced-rate data slots where downstream data and upstream data are only delivered once every <code>A2B_I2SRRATE.RRDIV</code> superframes. |
| | | 0 Disabled |
| | | 1 Enabled |
| 5:0 (R/W) | RRDIV | Reduced Rate Divide Select. The <code>A2B_I2SRRATE.RRDIV</code> bit field configures the superframe rate at which the I ² S/TDM data is active on the bus. For example, when <code>A2B_I2SRRATE.RRDIV</code> =16, I ² S/TDM data is active every 16th superframe rather than every superframe. Valid settings for this field are only those listed. |
| | | 1 Superframe frequency (SFF) |
| | | 2 SFF/2 |
| | | 4 SFF/4 |
| | | 8 SFF/8 |
| | | 12 SFF/12 |
| | | 16 SFF/16 |
| | | 20 SFF/20 |
| | | 24 SFF/24 |

Table 7-63: A2B_I2SRRATE Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration | |
|---------------------|----------|-------------------------|--------|
| | | 28 | SFF/28 |
| | | 32 | SFF/32 |

I2S Reduced Rate Control Register

The `A2B_I2SRCTL` register provides bits for controlling the I²S reduced rate strobe.

Address: 0x57

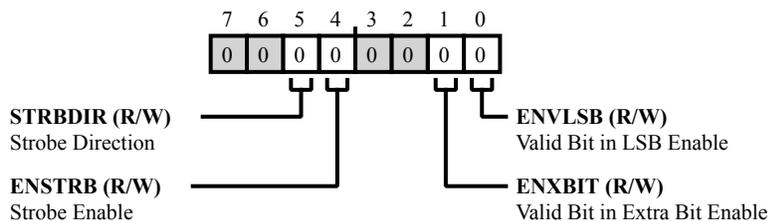


Figure 7-63: A2B_I2SRCTL Register Diagram

Table 7-64: A2B_I2SRCTL Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 5 (R/W) | STRBDIR | Strobe Direction. When the strobe signal is configured as an input, the <code>A2B_I2SRCTL.STRBDIR</code> bit influences the timing of frames on the bus. For a divide-by-N reduced rate, the strobe must be high once every N frames. |
| | | 0 Input |
| | | 1 Output |
| 4 (R/W) | ENSTRB | Strobe Enable. When the <code>A2B_I2SRCTL.ENSTRB</code> bit is set, the IO7 pin is used as a strobe, indicating the audio frame where the reduced-rate data is updated. |
| 1 (R/W) | ENXBIT | Valid Bit in Extra Bit Enable. The <code>A2B_I2SRCTL.ENXBIT</code> bit is only meaningful in a full-rate target that is receiving reduced rate data from the bus. It does not affect data going over the bus. When the <code>A2B_I2SRCTL.ENXBIT</code> bit is set, the bit after the LSB in each I ² S/TDM channel is high for the superframe with new data and low for the other superframes. |
| 0 (R/W) | ENVLSB | Valid Bit in LSB Enable. If the <code>A2B_I2SRCTL.ENVLSB</code> bit is set in a reduced-rate subordinate, the LSB of the data field is high for a new piece of data and low for a repeated piece of data. The <code>A2B_I2SRCTL.ENVLSB</code> bit is applicable in the subordinate node only. If the reduced-rate target node sets <code>A2B_I2SRCTL.ENVLSB</code> and the receiving controller node's <code>A2B_I2SRCTL.ENXBIT</code> bit is set, the output of the TDM data channel is xxxx11 for the first sampled word and xxxx00 for any repeated samples. Additionally, if the <code>A2B_I2SRATE.SHARE</code> bit is set in the reduced-rate target node, the LSB (additional bit) is high for the first data sample and low for the other samples. |

I2S Reduced Rate SYNC Offset Register (Target Only)

The `A2B_I2SRRSOFFS` register controls the SYNC offset for target nodes.

Address: 0x58

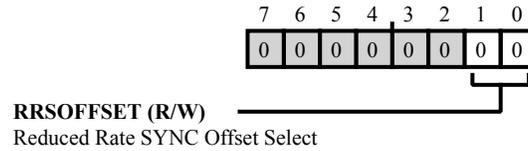


Figure 7-64: A2B_I2SRRSOFFS Register Diagram

Table 7-65: A2B_I2SRRSOFFS Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|-----------|---|
| 1:0 (R/W) | RRSOFFSET | <p>Reduced Rate SYNC Offset Select.</p> <p>A write of N to the <code>A2B_I2SRRSOFFS.RRSOFFSET</code> bit field instructs a target node, using a reduced I²S/TDM rate, to offset the SYNC edge to the left by N superframes. When the reduced-rate target's <code>A2B_I2SRATE.SHARE</code> bit is set, this field can only be programmed to 0 or 1.</p> |
| | | 0 No Offset |
| | | 1 1 Superframe Earlier |
| | | 2 2 Superframes Earlier |
| | | 3 3 Superframes Earlier |

CLKOUT1 Configuration Register

The `A2B_CLK1CFG` register enables an output clock on the `A2B_ADR1/A2B_IO1` pin and sets its frequency.

Address: `0x59`

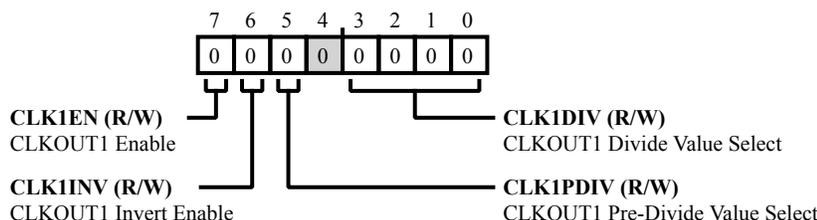


Figure 7-65: `A2B_CLK1CFG` Register Diagram

Table 7-66: `A2B_CLK1CFG` Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7 (R/W) | CLK1EN | CLKOUT1 Enable. When the <code>A2B_CLK1CFG.CLK1EN</code> bit is set, the <code>ADR1/IO1</code> pin is configured as a clock output, and GPIO functionality for the <code>ADR1/IO1</code> pin is disabled. |
| | | 0 Disabled |
| | | 1 Enabled |
| 6 (R/W) | CLK1INV | CLKOUT1 Invert Enable. When the <code>A2B_CLK1CFG.CLK1INV</code> bit is set, the clock output to the <code>ADR1/IO1</code> pin is inverted (moved 180 degrees out of phase). |
| | | 0 Disabled |
| | | 1 Enabled |
| 5 (R/W) | CLK1PDIV | CLKOUT1 Pre-Divide Value Select. The <code>A2B_CLK1CFG.CLK1PDIV</code> bit selects a pre-divide of either 2 or 32 from the PLL clock. At a 48 kHz sample frequency, the PLL clock has a frequency of 98.304 MHz. The PLL clock is 2048 times the sample frequency. |
| | | 0 Pre-divide is 2 |
| | | 1 Pre-divide is 32 |
| 3:0 (R/W) | CLK1DIV | CLKOUT1 Divide Value Select. The <code>A2B_CLK1CFG.CLK1DIV</code> bit field selects a divisor between 2 and 32 that is applied to the pre-divided clock before going to the pin. The divide ratio is $2 * (\text{CLK1DIV} + 1)$. |

CLKOUT2 Configuration Register

The `A2B_CLK2CFG` register enables an output clock on the `A2B_ADR2/A2B_IO2` pin and sets its frequency.

Address: `0x5A`

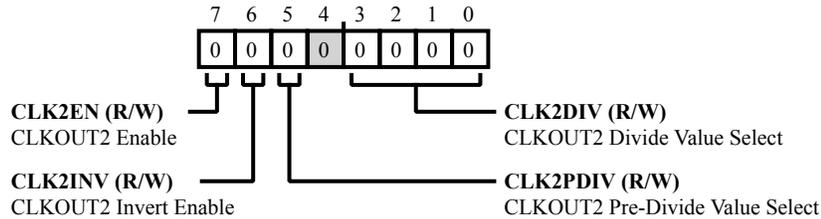


Figure 7-66: `A2B_CLK2CFG` Register Diagram

Table 7-67: `A2B_CLK2CFG` Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 7 (R/W) | CLK2EN | CLKOUT2 Enable. When the <code>A2B_CLK2CFG.CLK2EN</code> bit is set, the <code>ADR2/IO2</code> pin is configured as a clock output, and GPIO functionality for the <code>ADR2/IO2</code> pin is disabled. |
| | | 0 Disabled |
| | | 1 Enabled |
| 6 (R/W) | CLK2INV | CLKOUT2 Invert Enable. When the <code>A2B_CLK2CFG.CLK2INV</code> bit is set, the clock output to the <code>ADR2/IO2</code> pin is inverted (moved 180 degrees out of phase). |
| | | 0 Disabled |
| | | 1 Enabled |
| 5 (R/W) | CLK2PDIV | CLKOUT2 Pre-Divide Value Select. The <code>A2B_CLK2CFG.CLK2PDIV</code> bit selects a pre-divide of either 2 or 32 from the PLL clock. |
| | | 0 Pre-Divide is 2 |
| | | 1 Pre-Divide is 32 |
| 3:0 (R/W) | CLK2DIV | CLKOUT2 Divide Value Select. The <code>A2B_CLK2CFG.CLK2DIV</code> bit field selects a divisor between 2 and 32 that is applied to the pre-divided clock before going to the pin. The divide ratio is $2 * (\text{CLK1DIV} + 1)$. |

Bus Monitor Mode Configuration Register

The `A2B_BMMCFG` register configures settings for bus monitor mode.

Address: 0x5B

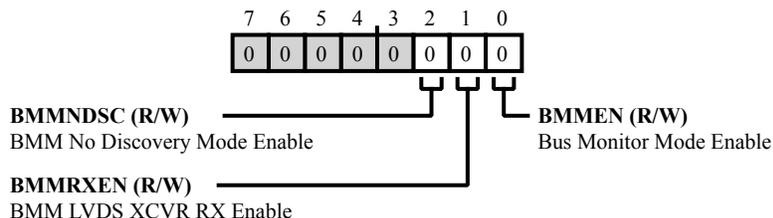


Figure 7-67: A2B_BMMCFG Register Diagram

Table 7-68: A2B_BMMCFG Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 2 (R/W) | BMMNDSC | BMM No Discovery Mode Enable. The <code>A2B_BMMCFG.BMMNDSC</code> bit is used to enable No Discovery Mode in Bus Monitor Mode. |
| | | 0 Disabled |
| | | 1 Enabled |
| 1 (R/W) | BMMRXEN | BMM LVDS XCVR RX Enable. The <code>A2B_BMMCFG.BMMRXEN</code> bit is used to enable LVDS RX in Bus Monitor Mode. |
| | | 0 Disabled |
| | | 1 Enabled |
| 0 (R/W) | BMMEN | Bus Monitor Mode Enable. The <code>A2B_BMMCFG.BMMEN</code> bit is used to enable Bus Monitor Mode. |
| | | 0 Disabled |
| | | 1 Enabled |

Sustain Configuration Register (Subordinate Only)

The `A2B_SUSCFG` register is used to configure sustain functionality in a subordinate node.

Address: 0x5C

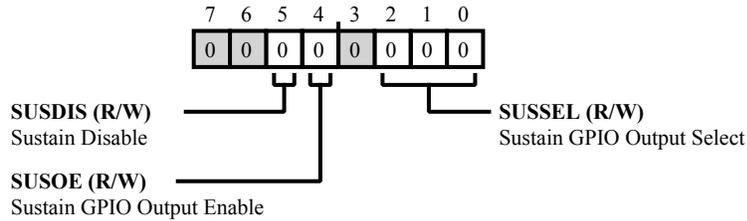


Figure 7-68: A2B_SUSCFG Register Diagram

Table 7-69: A2B_SUSCFG Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------------|----------|---------------------------------|
| 5 (R/W) | SUSDIS | Sustain Disable. |
| | | 0 Enable sustain mode |
| | | 1 Disable sustain mode |
| 4 (R/W) | SUSOE | Sustain GPIO Output Enable. |
| | | 0 Disable sustain mode output |
| | | 1 Enable sustain mode output |
| 2:0 (R/W) | SUSSEL | Sustain GPIO Output Select. |
| | | 0 Sustain output on IO0 |
| | | 1 Sustain output on IO1 |
| | | 2 Sustain output on IO2 |
| | | 3 Sustain output on IO3 |
| | | 4 Sustain output on IO4 |
| | | 5 Sustain output on IO5 |
| | | 6 Sustain output on IO6 |
| 7 Sustain output on IO7 | | |

PDM Control 2 Register

The `A2B_PDMCTL2` register provides a means of routing and handling PDM clock and data signals differently to accommodate various PDM configurations.

Address: 0x5D

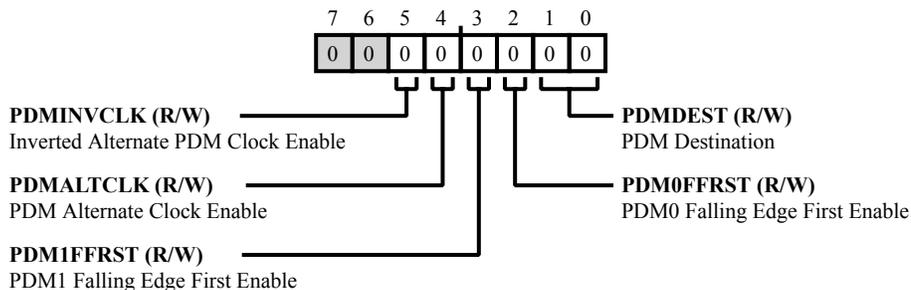


Figure 7-69: A2B_PDMCTL2 Register Diagram

Table 7-70: A2B_PDMCTL2 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|-----------|---|
| 5 (R/W) | PDMINVCLK | Inverted Alternate PDM Clock Enable. When the <code>A2B_PDMCTL2.PDMINVCLK</code> bit is set in a subordinate node, and the <code>A2B_PDMCTL2.PDMALTCLK</code> bit is set, an inverted version of the PDMCLK on the IO7 pin is driven on the BCLK pin. I ² S/TDM is still supported in this mode, but the BCLK frequency is constrained to 64x the SYNC frequency. |
| 4 (R/W) | PDMALTCLK | PDM Alternate Clock Enable. When the <code>A2B_PDMCTL2.PDMALTCLK</code> bit is set and at least one PDM input pin is enabled, the IO7 pin is used as the PDMCLK clock output pin. For a target node, this allows the BCLK frequency to be set from the I ² S/TDM configuration even when PDM functions are enabled. For a controller node, this allows the PDM clock to be a different frequency than the input BCLK. The frequency of the PDM clock on IO7 is 64x the SYNC frequency. If both PDM input pins are disabled (<code>A2B_PDMCTL.PDM0EN = A2B_PDMCTL.PDM1EN = 0</code>), the <code>A2B_PDMCTL2.PDMALTCLK</code> bit is ignored. |
| 3 (R/W) | PDM1FFRST | PDM1 Falling Edge First Enable. When the <code>A2B_PDMCTL2.PDM1FFRST</code> bit is cleared (default), PDM1 data on the DRX1 pin is sampled rising edge first. When the <code>A2B_PDMCTL2.PDM1FFRST</code> bit is set, the DRX1 pin is sampled falling edge first. |
| 2 (R/W) | PDM0FFRST | PDM0 Falling Edge First Enable. When the <code>A2B_PDMCTL2.PDM0FFRST</code> bit is cleared (default), PDM0 data on the DRX0 pin is sampled rising edge first. When the <code>A2B_PDMCTL2.PDM0FFRST</code> bit is set, the DRX0 pin is sampled falling edge first. |

Table 7-70: A2B_PDMCTL2 Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 1:0 (R/W) | PDMDEST | <p>PDM Destination.</p> <p>The A2B_PDMCTL2 . PDMDEST bit field selects how PDM data is routed. By default, PDM data received by the DRX0 and DRX1 pins goes to the A²B bus after demodulation. The demodulated data can instead or also be routed over the I²S/TDM port to the local node using one or more of the DTXn pins.</p> |
| | | 0 (Default) A ² B bus only |
| | | 1 DTXn pin(s) only |
| | | 2 A ² B bus and DTXn pin(s) |
| | | 3 Reserved |

Upstream Data RX Mask 0 Register (Target Only)

The `A2B_UPMASK0` register identifies which upstream data slots (from 0 to 7) are received from the A²B bus. These data slots may be transmitted via I²S/TDM and follow any downstream slots which were received by the target node during the downstream portion of the superframe (defined by the `A2B_LDNSLOTS` register). Changes to this register only take effect after setting the `A2B_CONTROL.NEWSTRCT` bit of the controller node.

Address: 0x60

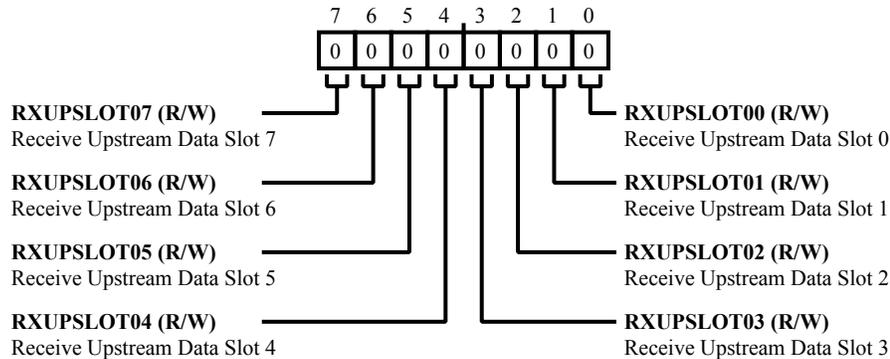


Figure 7-70: A2B_UPMASK0 Register Diagram

Table 7-71: A2B_UPMASK0 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|------------|--|
| 7 (R/W) | RXUPSLOT07 | Receive Upstream Data Slot 7. The <code>A2B_UPMASK0.RXUPSLOT07</code> bit defines whether or not upstream data slot 7 is received by the local target node. |
| | | 0 Upstream Data Slot 7 RX Disabled |
| | | 1 Upstream Data Slot 7 RX Enabled |
| 6 (R/W) | RXUPSLOT06 | Receive Upstream Data Slot 6. The <code>A2B_UPMASK0.RXUPSLOT06</code> bit defines whether or not upstream data slot 6 is received by the local target node. |
| | | 0 Upstream Data Slot 6 RX Disabled |
| | | 1 Upstream Data Slot 6 RX Enabled |
| 5 (R/W) | RXUPSLOT05 | Receive Upstream Data Slot 5. The <code>A2B_UPMASK0.RXUPSLOT05</code> bit defines whether or not upstream data slot 5 is received by the local target node. |
| | | 0 Upstream Data Slot 5 RX Disabled |
| | | 1 Upstream Data Slot 5 RX Enabled |

Table 7-71: A2B_UPMASK0 Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|------------|---|
| 4 (R/W) | RXUPSLOT04 | Receive Upstream Data Slot 4. The A2B_UPMASK0.RXUPSLOT04 bit defines whether or not upstream data slot 4 is received by the local target node. |
| | | 0 Upstream Data Slot 4 RX Disabled |
| | | 1 Upstream Data Slot 4 RX Enabled |
| 3 (R/W) | RXUPSLOT03 | Receive Upstream Data Slot 3. The A2B_UPMASK0.RXUPSLOT03 bit defines whether or not upstream data slot 3 is received by the local target node. |
| | | 0 Upstream Data Slot 3 RX Disabled |
| | | 1 Upstream Data Slot 3 RX Enabled |
| 2 (R/W) | RXUPSLOT02 | Receive Upstream Data Slot 2. The A2B_UPMASK0.RXUPSLOT02 bit defines whether or not upstream data slot 2 is received by the local target node. |
| | | 0 Upstream Data Slot 2 RX Disabled |
| | | 1 Upstream Data Slot 2 RX Enabled |
| 1 (R/W) | RXUPSLOT01 | Receive Upstream Data Slot 1. The A2B_UPMASK0.RXUPSLOT01 bit defines whether or not upstream data slot 1 is received by the local target node. |
| | | 0 Upstream Data Slot 1 RX Disabled |
| | | 1 Upstream Data Slot 1 RX Enabled |
| 0 (R/W) | RXUPSLOT00 | Receive Upstream Data Slot 0. The A2B_UPMASK0.RXUPSLOT00 bit defines whether or not upstream data slot 0 is received by the local target node. |
| | | 0 Upstream Data Slot 0 RX Disabled |
| | | 1 Upstream Data Slot 0 RX Enabled |

Upstream Data RX Mask 1 Register (Target Only)

The `A2B_UPMASK1` register identifies which upstream data slots (from 8 to 15) are received from the A²B bus. These data slots may be transmitted via I²S/TDM and follow any downstream slots which were received by the target node during the downstream portion of the superframe (defined by the `A2B_LDNSLOTS` register). Changes to this register only take effect after setting the `A2B_CONTROL.NEWSTRCT` bit of the controller node.

Address: 0x61

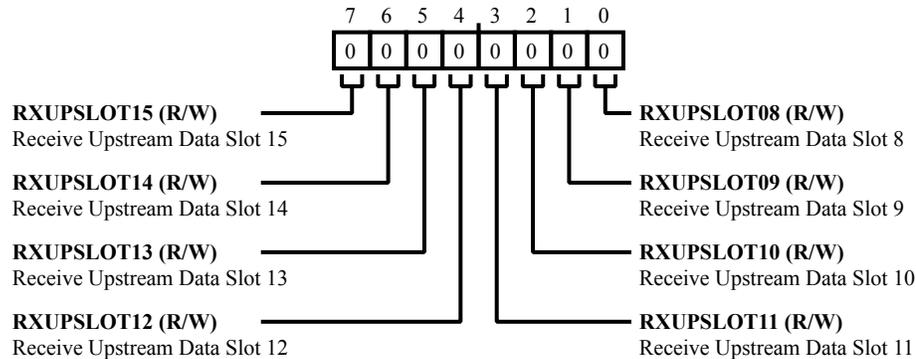


Figure 7-71: A2B_UPMASK1 Register Diagram

Table 7-72: A2B_UPMASK1 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|------------|--|
| 7 (R/W) | RXUPSLOT15 | Receive Upstream Data Slot 15. The <code>A2B_UPMASK1.RXUPSLOT15</code> bit defines whether or not upstream data slot 15 is received by the local target node. |
| | | 0 Upstream Data Slot 15 RX Disabled |
| | | 1 Upstream Data Slot 15 RX Enabled |
| 6 (R/W) | RXUPSLOT14 | Receive Upstream Data Slot 14. The <code>A2B_UPMASK1.RXUPSLOT14</code> bit defines whether or not upstream data slot 14 is received by the local target node. |
| | | 0 Upstream Data Slot 14 RX Disabled |
| | | 1 Upstream Data Slot 14 RX Enabled |
| 5 (R/W) | RXUPSLOT13 | Receive Upstream Data Slot 13. The <code>A2B_UPMASK1.RXUPSLOT13</code> bit defines whether or not upstream data slot 13 is received by the local target node. |
| | | 0 Upstream Data Slot 13 RX Disabled |
| | | 1 Upstream Data Slot 13 RX Enabled |

Table 7-72: A2B_UPMASK1 Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|------------|---|
| 4 (R/W) | RXUPSLOT12 | Receive Upstream Data Slot 12. The A2B_UPMASK1.RXUPSLOT12 bit defines whether or not upstream data slot 12 is received by the local target node. |
| | | 0 Upstream Data Slot 12 RX Disabled |
| | | 1 Upstream Data Slot 12 RX Enabled |
| 3 (R/W) | RXUPSLOT11 | Receive Upstream Data Slot 11. The A2B_UPMASK1.RXUPSLOT11 bit defines whether or not upstream data slot 11 is received by the local target node. |
| | | 0 Upstream Data Slot 11 RX Disabled |
| | | 1 Upstream Data Slot 11 RX Enabled |
| 2 (R/W) | RXUPSLOT10 | Receive Upstream Data Slot 10. The A2B_UPMASK1.RXUPSLOT10 bit defines whether or not upstream data slot 10 is received by the local target node. |
| | | 0 Upstream Data Slot 10 RX Disabled |
| | | 1 Upstream Data Slot 10 RX Enabled |
| 1 (R/W) | RXUPSLOT09 | Receive Upstream Data Slot 9. The A2B_UPMASK1.RXUPSLOT09 bit defines whether or not upstream data slot 9 is received by the local target node. |
| | | 0 Upstream Data Slot 9 RX Disabled |
| | | 1 Upstream Data Slot 9 RX Enabled |
| 0 (R/W) | RXUPSLOT08 | Receive Upstream Data Slot 8. The A2B_UPMASK1.RXUPSLOT08 bit defines whether or not upstream data slot 8 is received by the local target node. |
| | | 0 Upstream Data Slot 8 RX Disabled |
| | | 1 Upstream Data Slot 8 RX Enabled |

Upstream Data RX Mask 2 Register (Target Only)

The `A2B_UPMASK2` register identifies which upstream data slots (from 16 to 23) are received from the A²B bus. These data slots may be transmitted via I²S/TDM and follow any downstream slots which were received by the target node during the downstream portion of the superframe (defined by the `A2B_LDNSLOTS` register). Changes to this register only take effect after setting the `A2B_CONTROL.NEWSTRCT` bit of the controller node.

Address: 0x62

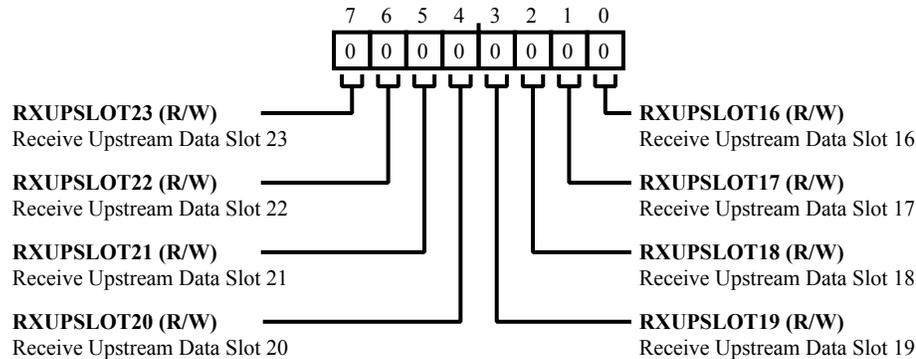


Figure 7-72: A2B_UPMASK2 Register Diagram

Table 7-73: A2B_UPMASK2 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|------------|--|
| 7 (R/W) | RXUPSLOT23 | Receive Upstream Data Slot 23. The <code>A2B_UPMASK2.RXUPSLOT23</code> bit defines whether or not upstream data slot 23 is received by the local target node. |
| | | 0 Upstream Data Slot 23 RX Disabled |
| | | 1 Upstream Data Slot 23 RX Enabled |
| 6 (R/W) | RXUPSLOT22 | Receive Upstream Data Slot 22. The <code>A2B_UPMASK2.RXUPSLOT22</code> bit defines whether or not upstream data slot 22 is received by the local target node. |
| | | 0 Upstream Data Slot 22 RX Disabled |
| | | 1 Upstream Data Slot 22 RX Enabled |
| 5 (R/W) | RXUPSLOT21 | Receive Upstream Data Slot 21. The <code>A2B_UPMASK2.RXUPSLOT21</code> bit defines whether or not upstream data slot 21 is received by the local target node. |
| | | 0 Upstream Data Slot 21 RX Disabled |
| | | 1 Upstream Data Slot 21 RX Enabled |

Table 7-73: A2B_UPMASK2 Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|------------|---|
| 4 (R/W) | RXUPSLOT20 | Receive Upstream Data Slot 20. The A2B_UPMASK2.RXUPSLOT20 bit defines whether or not upstream data slot 20 is received by the local target node. |
| | | 0 Upstream Data Slot 20 RX Disabled |
| | | 1 Upstream Data Slot 20 RX Enabled |
| 3 (R/W) | RXUPSLOT19 | Receive Upstream Data Slot 19. The A2B_UPMASK2.RXUPSLOT19 bit defines whether or not upstream data slot 19 is received by the local target node. |
| | | 0 Upstream Data Slot 19 RX Disabled |
| | | 1 Upstream Data Slot 19 RX Enabled |
| 2 (R/W) | RXUPSLOT18 | Receive Upstream Data Slot 18. The A2B_UPMASK2.RXUPSLOT18 bit defines whether or not upstream data slot 18 is received by the local target node. |
| | | 0 Upstream Data Slot 18 RX Disabled |
| | | 1 Upstream Data Slot 18 RX Enabled |
| 1 (R/W) | RXUPSLOT17 | Receive Upstream Data Slot 17. The A2B_UPMASK2.RXUPSLOT17 bit defines whether or not upstream data slot 17 is received by the local target node. |
| | | 0 Upstream Data Slot 17 RX Disabled |
| | | 1 Upstream Data Slot 17 RX Enabled |
| 0 (R/W) | RXUPSLOT16 | Receive Upstream Data Slot 16. The A2B_UPMASK2.RXUPSLOT16 bit defines whether or not upstream data slot 16 is received by the local target node. |
| | | 0 Upstream Data Slot 16 RX Disabled |
| | | 1 Upstream Data Slot 16 RX Enabled |

Upstream Data RX Mask 3 Register (Target Only)

The `A2B_UPMASK3` register identifies which upstream data slots (from 24 to 31) are received from the A²B bus. These data slots may be transmitted via I²S/TDM and follow any downstream slots which were received by the target node during the downstream portion of the superframe (defined by the `A2B_LDNSLOTS` register). Changes to this register only take effect after setting the `A2B_CONTROL.NEWSTRCT` bit of the controller node.

Address: 0x63

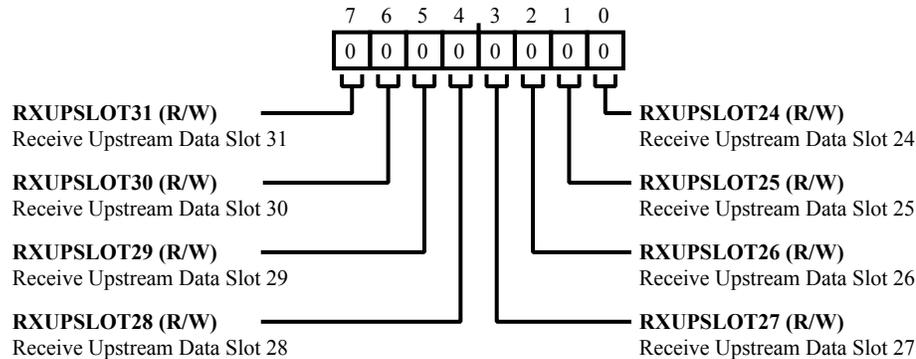


Figure 7-73: A2B_UPMASK3 Register Diagram

Table 7-74: A2B_UPMASK3 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|------------|--|
| 7 (R/W) | RXUPSLOT31 | Receive Upstream Data Slot 31. The <code>A2B_UPMASK3.RXUPSLOT31</code> bit defines whether or not upstream data slot 31 is received by the local target node. |
| | | 0 Upstream Data Slot 31 RX Disabled |
| | | 1 Upstream Data Slot 31 RX Enabled |
| 6 (R/W) | RXUPSLOT30 | Receive Upstream Data Slot 30. The <code>A2B_UPMASK3.RXUPSLOT30</code> bit defines whether or not upstream data slot 30 is received by the local target node. |
| | | 0 Upstream Data Slot 30 RX Disabled |
| | | 1 Upstream Data Slot 30 RX Enabled |
| 5 (R/W) | RXUPSLOT29 | Receive Upstream Data Slot 29. The <code>A2B_UPMASK3.RXUPSLOT29</code> bit defines whether or not upstream data slot 29 is received by the local target node. |
| | | 0 Upstream Data Slot 29 RX Disabled |
| | | 1 Upstream Data Slot 29 RX Enabled |

Table 7-74: A2B_UPMASK3 Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|------------|---|
| 4 (R/W) | RXUPSLOT28 | Receive Upstream Data Slot 28. The A2B_UPMASK3.RXUPSLOT28 bit defines whether or not upstream data slot 28 is received by the local target node. |
| | | 0 Upstream Data Slot 28 RX Disabled |
| | | 1 Upstream Data Slot 28 RX Enabled |
| 3 (R/W) | RXUPSLOT27 | Receive Upstream Data Slot 27. The A2B_UPMASK3.RXUPSLOT27 bit defines whether or not upstream data slot 27 is received by the local target node. |
| | | 0 Upstream Data Slot 27 RX Disabled |
| | | 1 Upstream Data Slot 27 RX Enabled |
| 2 (R/W) | RXUPSLOT26 | Receive Upstream Data Slot 26. The A2B_UPMASK3.RXUPSLOT26 bit defines whether or not upstream data slot 26 is received by the local target node. |
| | | 0 Upstream Data Slot 26 RX Disabled |
| | | 1 Upstream Data Slot 26 RX Enabled |
| 1 (R/W) | RXUPSLOT25 | Receive Upstream Data Slot 25. The A2B_UPMASK3.RXUPSLOT25 bit defines whether or not upstream data slot 25 is received by the local target node. |
| | | 0 Upstream Data Slot 25 RX Disabled |
| | | 1 Upstream Data Slot 25 RX Enabled |
| 0 (R/W) | RXUPSLOT24 | Receive Upstream Data Slot 24. The A2B_UPMASK3.RXUPSLOT24 bit defines whether or not upstream data slot 24 is received by the local target node. |
| | | 0 Disabled |
| | | 1 Enabled |

Local Upstream Channel Offset Register (Target Only)

In a target node, the `A2B_UPOFFSET` register defines the number of data channels received via I²S/TDM/PDM that are skipped before data slots are transmitted upstream on the A²B bus.

Address: 0x64

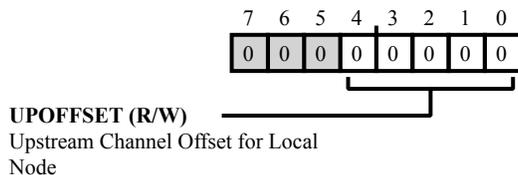


Figure 7-74: A2B_UPOFFSET Register Diagram

Table 7-75: A2B_UPOFFSET Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 4:0 (R/W) | UPOFFSET | Upstream Channel Offset for Local Node. The <code>A2B_UPOFFSET.UPOFFSET</code> bit field defines the number of data channels received via I ² S/TDM/PDM that are skipped before data slots are transmitted upstream on the A ² B bus. |

Downstream Data RX Mask 0 Register (Target Only)

The `A2B_DNMASK0` register identifies the downstream data slots (from 0 to 7) that are received from the A²B bus. These data slots may be transmitted via I²S/TDM. If none of the bits in this register are set, the `A2B_LDNSLOTS` register defines the number of downstream data slots taken by the local node, as in the AD2410. Changes to this register only take effect after setting the `A2B_CONTROL.NEWSTRCT` bit of the controller node.

Address: 0x65

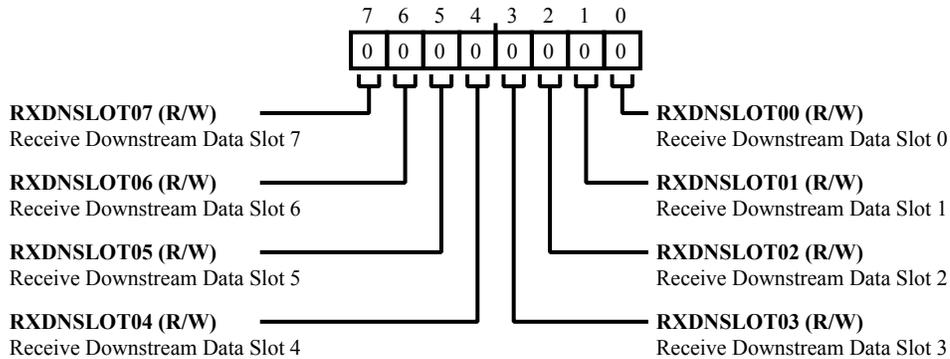


Figure 7-75: A2B_DNMASK0 Register Diagram

Table 7-76: A2B_DNMASK0 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|------------|--|
| 7 (R/W) | RXDNSLOT07 | Receive Downstream Data Slot 7. The <code>A2B_DNMASK0.RXDNSLOT07</code> bit defines whether or not downstream data slot 7 is received by the local target node. |
| | | 0 Downstream Data Slot 7 RX Disabled |
| | | 1 Downstream Data Slot 7 RX Enabled |
| 6 (R/W) | RXDNSLOT06 | Receive Downstream Data Slot 6. The <code>A2B_DNMASK0.RXDNSLOT06</code> bit defines whether or not downstream data slot 6 is received by the local target node. |
| | | 0 Downstream Data Slot 6 RX Disabled |
| | | 1 Downstream Data Slot 6 RX Enabled |
| 5 (R/W) | RXDNSLOT05 | Receive Downstream Data Slot 5. The <code>A2B_DNMASK0.RXDNSLOT05</code> bit defines whether or not downstream data slot 5 is received by the local target node. |
| | | 0 Downstream Data Slot 5 RX Disabled |
| | | 1 Downstream Data Slot 5 RX Enabled |

Table 7-76: A2B_DNMask0 Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|------------|---|
| 4 (R/W) | RXDNSLOT04 | Receive Downstream Data Slot 4. The A2B_DNMask0.RXDNSLOT04 bit defines whether or not downstream data slot 4 is received by the local target node. |
| | | 0 Downstream Data Slot 4 RX Disabled |
| | | 1 Downstream Data Slot 4 RX Enabled |
| 3 (R/W) | RXDNSLOT03 | Receive Downstream Data Slot 3. The A2B_DNMask0.RXDNSLOT03 bit defines whether or not downstream data slot 3 is received by the local target node. |
| | | 0 Downstream Data Slot 3 RX Disabled |
| | | 1 Downstream Data Slot 3 RX Enabled |
| 2 (R/W) | RXDNSLOT02 | Receive Downstream Data Slot 2. The A2B_DNMask0.RXDNSLOT02 bit defines whether or not downstream data slot 2 is received by the local target node. |
| | | 0 Downstream Data Slot 2 RX Disabled |
| | | 1 Downstream Data Slot 2 RX Enabled |
| 1 (R/W) | RXDNSLOT01 | Receive Downstream Data Slot 1. The A2B_DNMask0.RXDNSLOT01 bit defines whether or not downstream data slot 1 is received by the local target node. |
| | | 0 Downstream Data Slot 1 RX Disabled |
| | | 1 Downstream Data Slot 1 RX Enabled |
| 0 (R/W) | RXDNSLOT00 | Receive Downstream Data Slot 0. The A2B_DNMask0.RXDNSLOT00 bit defines whether or not downstream data slot 0 is received by the local target node. |
| | | 0 Downstream Data Slot 0 RX Disabled |
| | | 1 Downstream Data Slot 0 RX Enabled |

Downstream Data RX Mask 1 Register (Target Only)

The `A2B_DNMASK1` register identifies the downstream data slots (from 8 to 15) that are received from the A²B bus. These data slots may be transmitted via I²S/TDM. If none of the bits in this register are set, the `A2B_LDNSLOTS` register defines the number of downstream data slots taken by the local node, as in the AD2410. Changes to this register only take effect after setting the `A2B_CONTROL.NEWSTRCT` bit of the controller node.

Address: 0x66

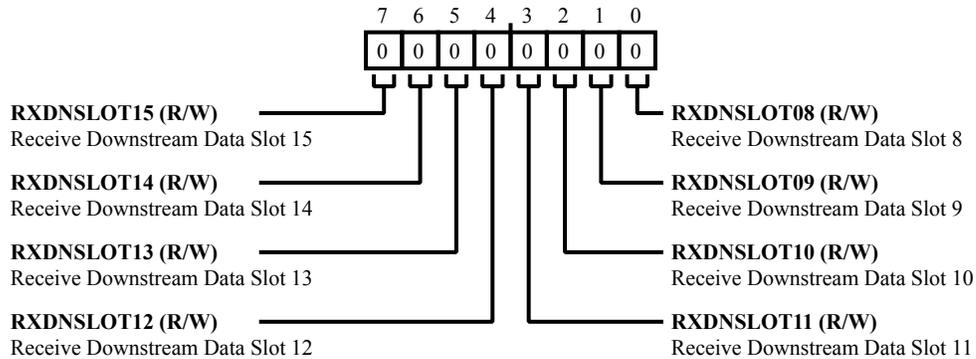


Figure 7-76: A2B_DNMASK1 Register Diagram

Table 7-77: A2B_DNMASK1 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|------------|--|
| 7 (R/W) | RXDNSLOT15 | Receive Downstream Data Slot 15. The <code>A2B_DNMASK1.RXDNSLOT15</code> bit defines whether or not downstream data slot 15 is received by the local target node. |
| | | 0 Downstream Data Slot 15 RX Disabled |
| | | 1 Downstream Data Slot 15 RX Enabled |
| 6 (R/W) | RXDNSLOT14 | Receive Downstream Data Slot 14. The <code>A2B_DNMASK1.RXDNSLOT14</code> bit defines whether or not downstream data slot 14 is received by the local target node. |
| | | 0 Downstream Data Slot 14 RX Disabled |
| | | 1 Downstream Data Slot 14 RX Enabled |
| 5 (R/W) | RXDNSLOT13 | Receive Downstream Data Slot 13. The <code>A2B_DNMASK1.RXDNSLOT13</code> bit defines whether or not downstream data slot 13 is received by the local target node. |
| | | 0 Downstream Data Slot 13 RX Disabled |
| | | 1 Downstream Data Slot 13 RX Enabled |

Table 7-77: A2B_DNMask1 Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|------------|---|
| 4 (R/W) | RXDNSLOT12 | Receive Downstream Data Slot 12. The A2B_DNMask1.RXDNSLOT12 bit defines whether or not downstream data slot 12 is received by the local target node. |
| | | 0 Downstream Data Slot 12 RX Disabled |
| | | 1 Downstream Data Slot 12 RX Enabled |
| 3 (R/W) | RXDNSLOT11 | Receive Downstream Data Slot 11. The A2B_DNMask1.RXDNSLOT11 bit defines whether or not downstream data slot 11 is received by the local target node. |
| | | 0 Downstream Data Slot 11 RX Disabled |
| | | 1 Downstream Data Slot 11 RX Enabled |
| 2 (R/W) | RXDNSLOT10 | Receive Downstream Data Slot 10. The A2B_DNMask1.RXDNSLOT10 bit defines whether or not downstream data slot 10 is received by the local target node. |
| | | 0 Downstream Data Slot 10 RX Disabled |
| | | 1 Downstream Data Slot 10 RX Enabled |
| 1 (R/W) | RXDNSLOT09 | Receive Downstream Data Slot 9. The A2B_DNMask1.RXDNSLOT09 bit defines whether or not downstream data slot 9 is received by the local target node. |
| | | 0 Downstream Data Slot 9 RX Disabled |
| | | 1 Downstream Data Slot 9 RX Enabled |
| 0 (R/W) | RXDNSLOT08 | Receive Downstream Data Slot 8. The A2B_DNMask1.RXDNSLOT08 bit defines whether or not downstream data slot 8 is received by the local target node. |
| | | 0 Downstream Data Slot 8 RX Disabled |
| | | 1 Downstream Data Slot 8 RX Enabled |

Downstream Data RX Mask 2 Register (Target Only)

The `A2B_DNMASK2` register identifies the downstream data slots (from 16 to 23) that are received from the A²B bus. These data slots may be transmitted via I²S/TDM. If none of the bits in this register are set, the `A2B_LDNSLOTS` register defines the number of downstream data slots taken by the local node. Changes to this register only take effect after setting the `A2B_CONTROL.NEWSTRCT` bit of the controller node.

Address: 0x67

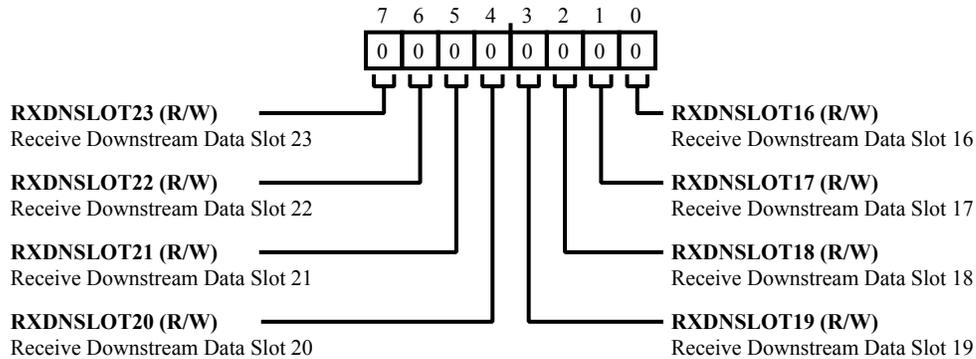


Figure 7-77: A2B_DNMASK2 Register Diagram

Table 7-78: A2B_DNMASK2 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|------------|--|
| 7 (R/W) | RXDNSLOT23 | Receive Downstream Data Slot 23. The <code>A2B_DNMASK2.RXDNSLOT23</code> bit defines whether or not downstream data slot 23 is received by the local target node. |
| | | 0 Downstream Data Slot 23 RX Disabled |
| | | 1 Downstream Data Slot 23 RX Enabled |
| 6 (R/W) | RXDNSLOT22 | Receive Downstream Data Slot 22. The <code>A2B_DNMASK2.RXDNSLOT22</code> bit defines whether or not downstream data slot 22 is received by the local target node. |
| | | 0 Downstream Data Slot 22 RX Disabled |
| | | 1 Downstream Data Slot 22 RX Enabled |
| 5 (R/W) | RXDNSLOT21 | Receive Downstream Data Slot 21. The <code>A2B_DNMASK2.RXDNSLOT21</code> bit defines whether or not downstream data slot 21 is received by the local target node. |
| | | 0 Downstream Data Slot 21 RX Disabled |
| | | 1 Downstream Data Slot 21 RX Enabled |

Table 7-78: A2B_DNMASK2 Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|------------|---|
| 4 (R/W) | RXDNSLOT20 | Receive Downstream Data Slot 20. The A2B_DNMASK2.RXDNSLOT20 bit defines whether or not downstream data slot 20 is received by the local target node. |
| | | 0 Downstream Data Slot 20 RX Disabled |
| | | 1 Downstream Data Slot 20 RX Enabled |
| 3 (R/W) | RXDNSLOT19 | Receive Downstream Data Slot 19. The A2B_DNMASK2.RXDNSLOT19 bit defines whether or not downstream data slot 19 is received by the local target node. |
| | | 0 Downstream Data Slot 19 RX Disabled |
| | | 1 Downstream Data Slot 19 RX Enabled |
| 2 (R/W) | RXDNSLOT18 | Receive Downstream Data Slot 18. The A2B_DNMASK2.RXDNSLOT18 bit defines whether or not downstream data slot 18 is received by the local target node. |
| | | 0 Downstream Data Slot 18 RX Disabled |
| | | 1 Downstream Data Slot 18 RX Enabled |
| 1 (R/W) | RXDNSLOT17 | Receive Downstream Data Slot 17. The A2B_DNMASK2.RXDNSLOT17 bit defines whether or not downstream data slot 17 is received by the local target node. |
| | | 0 Downstream Data Slot 17 RX Disabled |
| | | 1 Downstream Data Slot 17 RX Enabled |
| 0 (R/W) | RXDNSLOT16 | Receive Downstream Data Slot 16. The A2B_DNMASK2.RXDNSLOT16 bit defines whether or not downstream data slot 16 is received by the local target node. |
| | | 0 Downstream Data Slot 16 RX Disabled |
| | | 1 Downstream Data Slot 16 RX Enabled |

Downstream Data RX Mask 3 Register (Target Only)

The `A2B_DNMASK3` register identifies the downstream data slots (from 24 to 31) that are received from the A²B bus. These data slots may be transmitted via I²S/TDM. If none of the bits in this register are set, the `A2B_LDNSLOTS` register defines the number of downstream data slots taken by the local node, as in the AD2410. Changes to this register only take effect after setting the `A2B_CONTROL.NEWSTRCT` bit of the controller node.

Address: 0x68

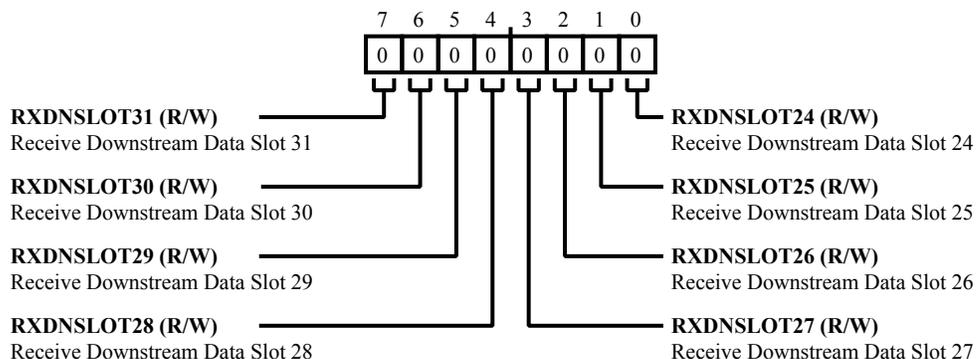


Figure 7-78: A2B_DNMASK3 Register Diagram

Table 7-79: A2B_DNMASK3 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|------------|--|
| 7 (R/W) | RXDNSLOT31 | Receive Downstream Data Slot 31. The <code>A2B_DNMASK3.RXDNSLOT31</code> bit defines whether or not downstream data slot 31 is received by the local target node. |
| | | 0 Downstream Data Slot 31 RX Disabled |
| | | 1 Downstream Data Slot 31 RX Enabled |
| 6 (R/W) | RXDNSLOT30 | Receive Downstream Data Slot 30. The <code>A2B_DNMASK3.RXDNSLOT30</code> bit defines whether or not downstream data slot 30 is received by the local target node. |
| | | 0 Downstream Data Slot 30 RX Disabled |
| | | 1 Downstream Data Slot 30 RX Enabled |
| 5 (R/W) | RXDNSLOT29 | Receive Downstream Data Slot 29. The <code>A2B_DNMASK3.RXDNSLOT29</code> bit defines whether or not downstream data slot 29 is received by the local target node. |
| | | 0 Downstream Data Slot 29 RX Disabled |
| | | 1 Downstream Data Slot 29 RX Enabled |

Table 7-79: A2B_DNMask3 Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|------------|---|
| 4 (R/W) | RXDNSLOT28 | Receive Downstream Data Slot 28. The A2B_DNMask3.RXDNSLOT28 bit defines whether or not downstream data slot 28 is received by the local target node. |
| | | 0 Downstream Data Slot 28 RX Disabled |
| | | 1 Downstream Data Slot 28 RX Enabled |
| 3 (R/W) | RXDNSLOT27 | Receive Downstream Data Slot 27. The A2B_DNMask3.RXDNSLOT27 bit defines whether or not downstream data slot 27 is received by the local target node. |
| | | 0 Downstream Data Slot 27 RX Disabled |
| | | 1 Downstream Data Slot 27 RX Enabled |
| 2 (R/W) | RXDNSLOT26 | Receive Downstream Data Slot 26. The A2B_DNMask3.RXDNSLOT26 bit defines whether or not downstream data slot 26 is received by the local target node. |
| | | 0 Downstream Data Slot 26 RX Disabled |
| | | 1 Downstream Data Slot 26 RX Enabled |
| 1 (R/W) | RXDNSLOT25 | Receive Downstream Data Slot 25. The A2B_DNMask3.RXDNSLOT25 bit defines whether or not downstream data slot 25 is received by the local target node. |
| | | 0 Downstream Data Slot 25 RX Disabled |
| | | 1 Downstream Data Slot 25 RX Enabled |
| 0 (R/W) | RXDNSLOT24 | Receive Downstream Data Slot 24. The A2B_DNMask3.RXDNSLOT24 bit defines whether or not downstream data slot 24 is received by the local target node. |
| | | 0 Downstream Data Slot 24 RX Disabled |
| | | 1 Downstream Data Slot 24 RX Enabled |

Local Downstream Channel Offset Register (Target Only)

In a target node, the `A2B_DNOFFSET` register defines the number of data channels received via I²S/TDM/PDM that are skipped before data slots are transmitted downstream on the A²B bus. The value in the `A2B_DNOFFSET` register is used only if any of the bits in the `A2B_DNMASK0` through `A2B_DNMASK3` registers are set and the `A2B_LDNSLOTS` register is non-zero.

Address: 0x69

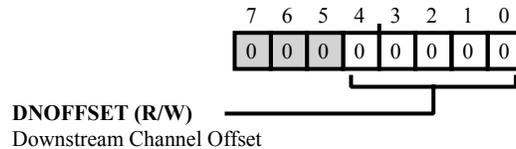


Figure 7-79: A2B_DNOFFSET Register Diagram

Table 7-80: A2B_DNOFFSET Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 4:0 (R/W) | DNOFFSET | Downstream Channel Offset. The <code>A2B_DNOFFSET.DNOFFSET</code> bit field defines the number of data channels received via I ² S/TDM/PDM that are skipped before data slots are transmitted downstream on the A ² B bus. |

Chip ID Register 0

The `A2B_CHIPID0` through `A2B_CHIPID5` registers concatenate to form a unique 48-bit ID for the transceiver, where `A2B_CHIPID0` contains the LSB (bits 7:0) and `A2B_CHIPID5` contains the MSB (bits 47:40).

Address: 0x6A

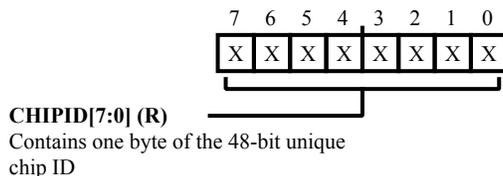


Figure 7-80: A2B_CHIPID0 Register Diagram

Table 7-81: A2B_CHIPID0 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7:0 (R/NW) | CHIPID | Contains one byte of the 48-bit unique chip ID. |

Chip ID Register 1

The A2B_CHIPID0 through A2B_CHIPID5 registers concatenate to form a unique 48-bit ID for the transceiver, where A2B_CHIPID0 contains the LSB (bits 7:0) and A2B_CHIPID5 contains the MSB (bits 47:40).

Address: 0x6B

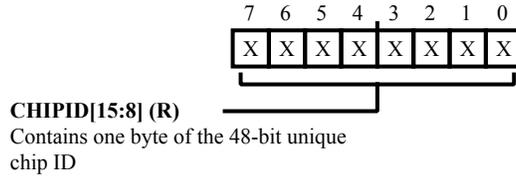


Figure 7-81: A2B_CHIPID1 Register Diagram

Table 7-82: A2B_CHIPID1 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7:0 (R/NW) | CHIPID | Contains one byte of the 48-bit unique chip ID. |

Chip ID Register 2

The A2B_CHIPID0 through A2B_CHIPID5 registers concatenate to form a unique 48-bit ID for the transceiver, where A2B_CHIPID0 contains the LSB (bits 7:0) and A2B_CHIPID5 contains the MSB (bits 47:40).

Address: 0x6C

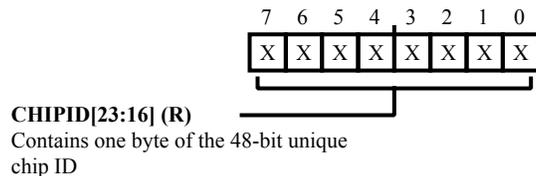


Figure 7-82: A2B_CHIPID2 Register Diagram

Table 7-83: A2B_CHIPID2 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7:0 (R/NW) | CHIPID | Contains one byte of the 48-bit unique chip ID. |

Chip ID Register 3

The A2B_CHIPID0 through A2B_CHIPID5 registers concatenate to form a unique 48-bit ID for the transceiver, where A2B_CHIPID0 contains the LSB (bits 7:0) and A2B_CHIPID5 contains the MSB (bits 47:40).

Address: 0x6D

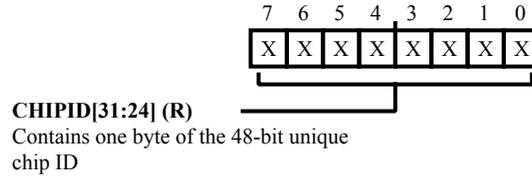


Figure 7-83: A2B_CHIPID3 Register Diagram

Table 7-84: A2B_CHIPID3 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7:0 (R/NW) | CHIPID | Contains one byte of the 48-bit unique chip ID. |

Chip ID Register 4

The A2B_CHIPID0 through A2B_CHIPID5 registers concatenate to form a unique 48-bit ID for the transceiver, where A2B_CHIPID0 contains the LSB (bits 7:0) and A2B_CHIPID5 contains the MSB (bits 47:40).

Address: 0x6E

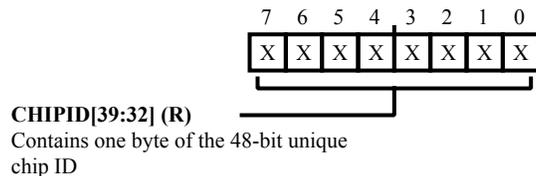


Figure 7-84: A2B_CHIPID4 Register Diagram

Table 7-85: A2B_CHIPID4 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7:0 (R/NW) | CHIPID | Contains one byte of the 48-bit unique chip ID. |

Chip ID Register 5

The A2B_CHIPID0 through A2B_CHIPID5 registers concatenate to form a unique 48-bit ID for the transceiver, where A2B_CHIPID0 contains the LSB (bits 7:0) and A2B_CHIPID5 contains the MSB (bits 47:40).

Address: 0x6F

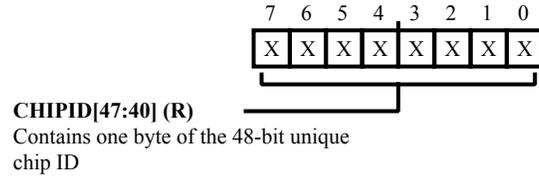


Figure 7-85: A2B_CHIPID5 Register Diagram

Table 7-86: A2B_CHIPID5 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7:0 (R/NW) | CHIPID | Contains one byte of the 48-bit unique chip ID. |

GPIO Over Distance Enable Register

The `A2B_GPIODEN` register controls the general-purpose I/O pins for use in GPIO Over Distance.

Address: 0x80

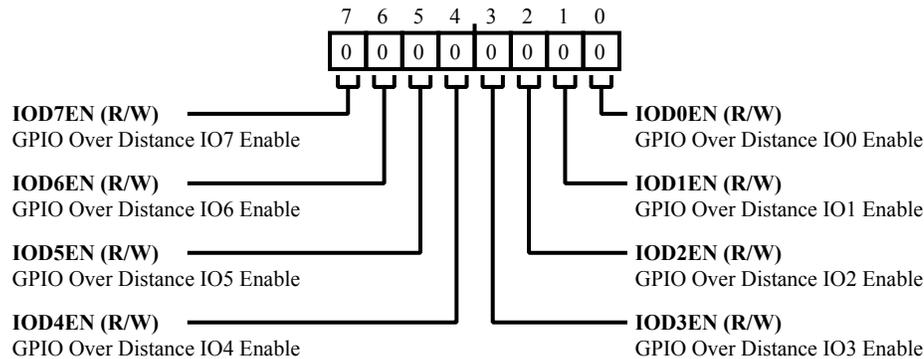


Figure 7-86: A2B_GPIODEN Register Diagram

Table 7-87: A2B_GPIODEN Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|---|
| 7 (R/W) | IOD7EN | GPIO Over Distance IO7 Enable. The <code>A2B_GPIODEN.IOD7EN</code> bit enables GPIO Over Distance for IO7. |
| | | 0 GPIO Over Distance for IO7 Disabled |
| | | 1 GPIO Over Distance for IO7 Enabled |
| 6 (R/W) | IOD6EN | GPIO Over Distance IO6 Enable. The <code>A2B_GPIODEN.IOD6EN</code> bit enables GPIO Over Distance for IO6. |
| | | 0 GPIO Over Distance for IO6 Disabled |
| | | 1 GPIO Over Distance for IO6 Enabled |
| 5 (R/W) | IOD5EN | GPIO Over Distance IO5 Enable. The <code>A2B_GPIODEN.IOD5EN</code> bit enables GPIO Over Distance for IO5. |
| | | 0 GPIO Over Distance for IO5 Disabled |
| | | 1 GPIO Over Distance for IO5 Enabled |
| 4 (R/W) | IOD4EN | GPIO Over Distance IO4 Enable. The <code>A2B_GPIODEN.IOD4EN</code> bit enables GPIO Over Distance for IO4. |
| | | 0 GPIO Over Distance for IO4 Disabled |
| | | 1 GPIO Over Distance for IO4 Enabled |

Table 7-87: A2B_GPIODEN Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 3 (R/W) | IOD3EN | GPIO Over Distance IO3 Enable. The A2B_GPIODEN.IOD3EN bit enables GPIO Over Distance for IO3. |
| | | 0 GPIO Over Distance for IO3 Disabled |
| | | 1 GPIO Over Distance for IO3 Enabled |
| 2 (R/W) | IOD2EN | GPIO Over Distance IO2 Enable. The A2B_GPIODEN.IOD2EN bit enables GPIO Over Distance for IO2. |
| | | 0 GPIO Over Distance for IO2 Disabled |
| | | 1 GPIO Over Distance for IO2 Enabled |
| 1 (R/W) | IOD1EN | GPIO Over Distance IO1 Enable. The A2B_GPIODEN.IOD1EN bit enables GPIO Over Distance for IO1. |
| | | 0 GPIO Over Distance for IO1 Disabled |
| | | 1 GPIO Over Distance for IO1 Enabled |
| 0 (R/W) | IOD0EN | GPIO Over Distance IO0 Enable. The A2B_GPIODEN.IOD0EN bit enables GPIO Over Distance for IO0. |
| | | 0 GPIO Over Distance for IO0 Disabled |
| | | 1 GPIO Over Distance for IO0 Enabled |

GPIO Over Distance Mask 0 Register

Address: 0x81

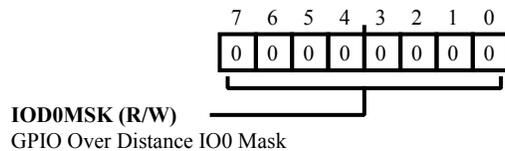


Figure 7-87: A2B_GPIOD0MSK Register Diagram

Table 7-88: A2B_GPIOD0MSK Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|------------------------------|
| 7:0 (R/W) | IOD0MSK | GPIO Over Distance IO0 Mask. |

GPIO Over Distance Mask 1 Register

Address: 0x82

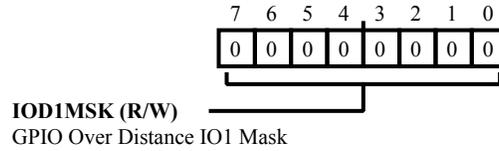


Figure 7-88: A2B_GPIOD1MSK Register Diagram

Table 7-89: A2B_GPIOD1MSK Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|------------------------------|
| 7:0 (R/W) | IOD1MSK | GPIO Over Distance IO1 Mask. |

GPIO Over Distance Mask 2 Register

Address: 0x83

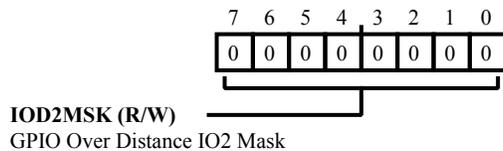


Figure 7-89: A2B_GPIOD2MSK Register Diagram

Table 7-90: A2B_GPIOD2MSK Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|------------------------------|
| 7:0 (R/W) | IOD2MSK | GPIO Over Distance IO2 Mask. |

GPIO Over Distance Mask 3 Register

Address: 0x84

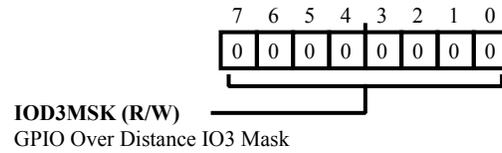


Figure 7-90: A2B_GPIOD3MSK Register Diagram

Table 7-91: A2B_GPIOD3MSK Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|------------------------------|
| 7:0 (R/W) | IOD3MSK | GPIO Over Distance IO3 Mask. |

GPIO Over Distance Mask 4 Register

Address: 0x85

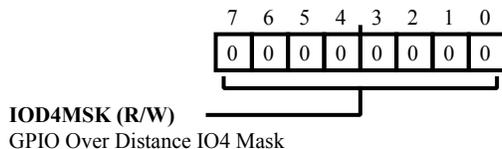


Figure 7-91: A2B_GPIOD4MSK Register Diagram

Table 7-92: A2B_GPIOD4MSK Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|------------------------------|
| 7:0 (R/W) | IOD4MSK | GPIO Over Distance IO4 Mask. |

GPIO Over Distance Mask 5 Register

Address: 0x86

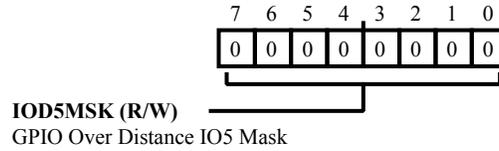


Figure 7-92: A2B_GPIOD5MSK Register Diagram

Table 7-93: A2B_GPIOD5MSK Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|------------------------------|
| 7:0 (R/W) | IOD5MSK | GPIO Over Distance IO5 Mask. |

GPIO Over Distance Mask 6 Register

Address: 0x87

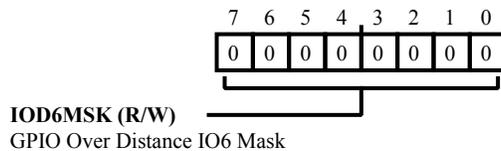


Figure 7-93: A2B_GPIOD6MSK Register Diagram

Table 7-94: A2B_GPIOD6MSK Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|------------------------------|
| 7:0 (R/W) | IOD6MSK | GPIO Over Distance IO6 Mask. |

GPIO Over Distance Mask 7 Register

Address: 0x88

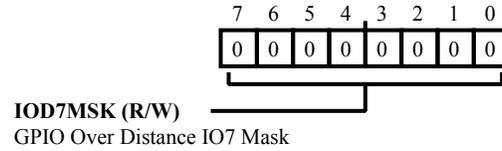


Figure 7-94: A2B_GPIOD7MSK Register Diagram

Table 7-95: A2B_GPIOD7MSK Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|------------------------------|
| 7:0 (R/W) | IOD7MSK | GPIO Over Distance IO7 Mask. |

GPIO Over Distance Data Register

Address: 0x89

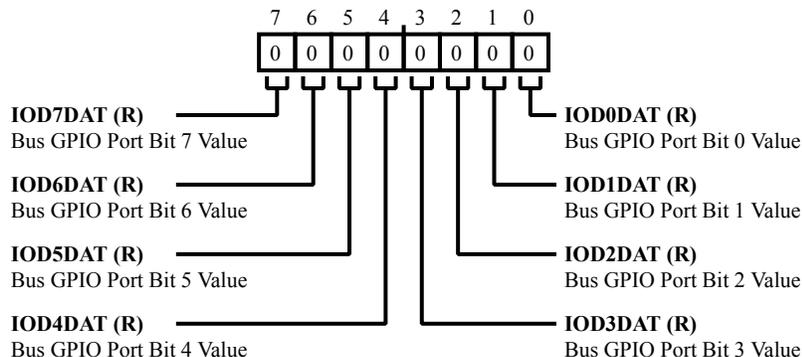


Figure 7-95: A2B_GPIODDAT Register Diagram

Table 7-96: A2B_GPIODDAT Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|----------------------------|
| 7 (R/NW) | IOD7DAT | Bus GPIO Port Bit 7 Value. |
| 6 (R/NW) | IOD6DAT | Bus GPIO Port Bit 6 Value. |
| 5 (R/NW) | IOD5DAT | Bus GPIO Port Bit 5 Value. |
| 4 (R/NW) | IOD4DAT | Bus GPIO Port Bit 4 Value. |
| 3 (R/NW) | IOD3DAT | Bus GPIO Port Bit 3 Value. |
| 2 (R/NW) | IOD2DAT | Bus GPIO Port Bit 2 Value. |
| 1 (R/NW) | IOD1DAT | Bus GPIO Port Bit 1 Value. |
| 0 (R/NW) | IOD0DAT | Bus GPIO Port Bit 0 Value. |

GPIO Over Distance Invert Register

Address: 0x8A

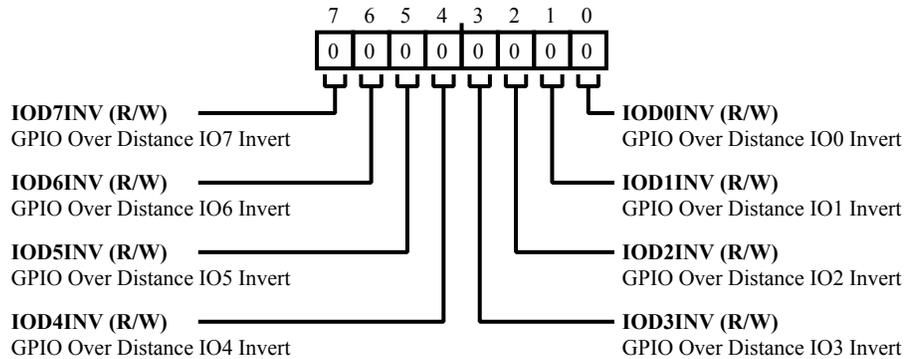


Figure 7-96: A2B_GPIODINV Register Diagram

Table 7-97: A2B_GPIODINV Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--------------------------------|
| 7 (R/W) | IOD7INV | GPIO Over Distance IO7 Invert. |
| 6 (R/W) | IOD6INV | GPIO Over Distance IO6 Invert. |
| 5 (R/W) | IOD5INV | GPIO Over Distance IO5 Invert. |
| 4 (R/W) | IOD4INV | GPIO Over Distance IO4 Invert. |
| 3 (R/W) | IOD3INV | GPIO Over Distance IO3 Invert. |
| 2 (R/W) | IOD2INV | GPIO Over Distance IO2 Invert. |
| 1 (R/W) | IOD1INV | GPIO Over Distance IO1 Invert. |
| 0 (R/W) | IOD0INV | GPIO Over Distance IO0 Invert. |

Mailbox 0 Control Register (Subordinate Only)

The `A2B_MBOX0CTL` register contains bits that control direction, message length and interrupts.

Address: 0x90

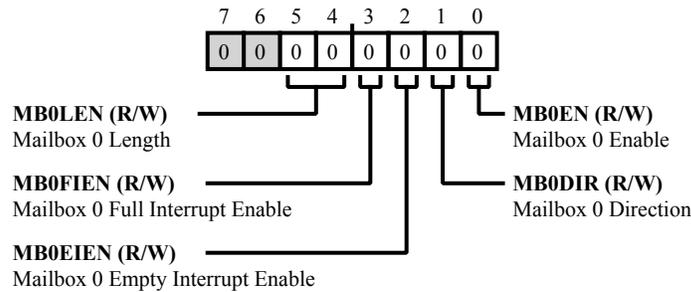


Figure 7-97: A2B_MBOX0CTL Register Diagram

Table 7-98: A2B_MBOX0CTL Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 5:4 (R/W) | MB0LEN | Mailbox 0 Length. The <code>A2B_MBOX0CTL.MB0LEN</code> bit field controls the length of Mailbox 0. |
| | | 0 1 Byte |
| | | 1 2 Bytes |
| | | 2 3 Bytes |
| | | 3 4 Bytes |
| 3 (R/W) | MB0FIEN | Mailbox 0 Full Interrupt Enable. The <code>A2B_MBOX0CTL.MB0FIEN</code> bit enables an interrupt which is generated when Mailbox 0 becomes full. |
| | | 0 Mailbox 0 Interrupt on Full Disabled |
| | | 1 Mailbox 0 Interrupt on Full Enabled |
| 2 (R/W) | MB0EIEN | Mailbox 0 Empty Interrupt Enable. The <code>A2B_MBOX0CTL.MB0EIEN</code> bit enables an interrupt which is generated when Mailbox 0 becomes empty. |
| | | 0 Mailbox 0 Interrupt on Empty Disabled |
| | | 1 Mailbox 0 Interrupt on Empty Enabled |
| 1 (R/W) | MB0DIR | Mailbox 0 Direction. The <code>A2B_MBOX0CTL.MB0DIR</code> bit controls the direction of Mailbox 0. |
| | | 0 Mailbox 0 is Receive Mailbox |
| | | 1 Mailbox 0 is Transmit Mailbox |

Table 7-98: A2B_MBOX0CTL Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration | |
|---------------------|----------|--|--------------------|
| 0 (R/W) | MBOEN | Mailbox 0 Enable. Setting the A2B_MBOX0CTL.MBOEN bit enables Mailbox 0. | |
| | | 0 | Mailbox 0 Disabled |
| | | 1 | Mailbox 0 Enabled |

Mailbox 0 Status Register (Subordinate Only)

The `A2B_MBOX0STAT` register reports the status of the configured mailbox interrupts.

Address: 0x91

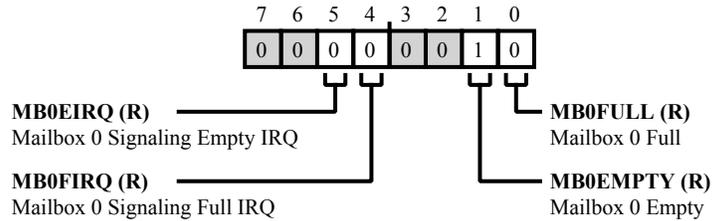


Figure 7-98: A2B_MBOX0STAT Register Diagram

Table 7-99: A2B_MBOX0STAT Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 5 (R/NW) | MB0EIRQ | Mailbox 0 Signaling Empty IRQ. The <code>A2B_MBOX0STAT.MB0EIRQ</code> bit indicates whether or not the Mailbox 0 empty interrupt is active. |
| | | 0 Mailbox 0 Empty Interrupt Inactive |
| | | 1 Mailbox 0 Empty Interrupt Active |
| 4 (R/NW) | MB0FIRQ | Mailbox 0 Signaling Full IRQ. The <code>A2B_MBOX0STAT.MB0FIRQ</code> bit indicates whether or not the Mailbox 0 full interrupt is active. |
| | | 0 Mailbox 0 Full Interrupt Inactive |
| | | 1 Mailbox 0 Full Interrupt Active |
| 1 (R/NW) | MB0EMPTY | Mailbox 0 Empty. The <code>A2B_MBOX0STAT.MB0EMPTY</code> bit indicates whether or not Mailbox 0 is empty. |
| | | 0 Mailbox 0 Currently Not Empty |
| | | 1 Mailbox 0 Currently Empty |
| 0 (R/NW) | MB0FULL | Mailbox 0 Full. The <code>A2B_MBOX0STAT.MB0FULL</code> bit indicates whether or not Mailbox 0 is full. |
| | | 0 Mailbox 0 Currently Not Full |
| | | 1 Mailbox 0 Currently Full |

Mailbox 0 Byte 0 Register (Subordinate Only)

Address: 0x92

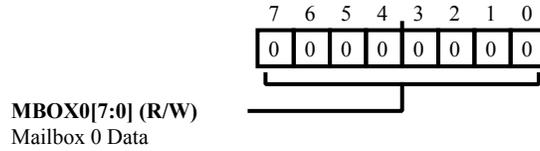


Figure 7-99: A2B_MBOX0B0 Register Diagram

Table 7-100: A2B_MBOX0B0 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|-------------------------|
| 7:0 (R/W) | MBOX0 | Mailbox 0 Data. |

Mailbox 0 Byte 1 Register (Subordinate Only)

Address: 0x93

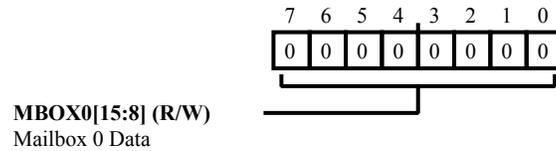


Figure 7-100: A2B_MBOX0B1 Register Diagram

Table 7-101: A2B_MBOX0B1 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|-------------------------|
| 7:0 (R/W) | MBOX0 | Mailbox 0 Data. |

Mailbox 0 Byte 2 Register (Subordinate Only)

Address: 0x94

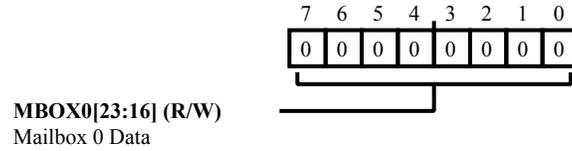


Figure 7-101: A2B_MBOX0B2 Register Diagram

Table 7-102: A2B_MBOX0B2 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|-------------------------|
| 7:0 (R/W) | MBOX0 | Mailbox 0 Data. |

Mailbox 0 Byte 3 Register (Subordinate Only)

Address: 0x95

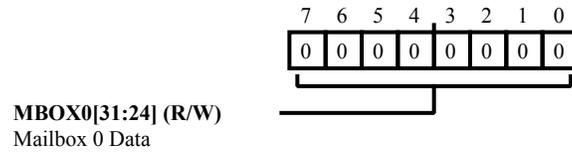


Figure 7-102: A2B_MBOX0B3 Register Diagram

Table 7-103: A2B_MBOX0B3 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|-------------------------|
| 7:0 (R/W) | MBOX0 | Mailbox 0 Data. |

Mailbox 1 Control Register (Subordinate Only)

The `A2B_MBOX1CTL` register contains bits that control direction, message length and interrupts.

Address: 0x96

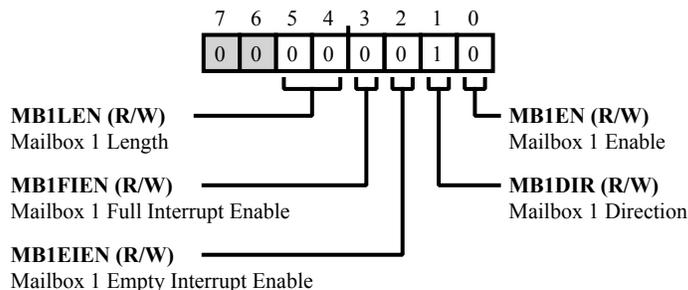


Figure 7-103: A2B_MBOX1CTL Register Diagram

Table 7-104: A2B_MBOX1CTL Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 5:4 (R/W) | MB1LEN | Mailbox 1 Length. The <code>A2B_MBOX1CTL.MB1LEN</code> bit field controls the length of Mailbox 1. |
| | | 0 1 Byte |
| | | 1 2 Bytes |
| | | 2 3 Bytes |
| | | 3 4 Bytes |
| 3 (R/W) | MB1FIEN | Mailbox 1 Full Interrupt Enable. The <code>A2B_MBOX1CTL.MB1FIEN</code> bit enables an interrupt which is generated when Mailbox 1 becomes full. |
| | | 0 Mailbox 1 Interrupt on Full Disabled |
| | | 1 Mailbox 1 Interrupt on Full Enabled |
| 2 (R/W) | MB1EIEN | Mailbox 1 Empty Interrupt Enable. The <code>A2B_MBOX1CTL.MB1EIEN</code> bit enables an interrupt which is generated when Mailbox 1 becomes empty. |
| | | 0 Mailbox 1 Interrupt on Empty Disabled |
| | | 1 Mailbox 1 Interrupt on Empty Enabled |
| 1 (R/W) | MB1DIR | Mailbox 1 Direction. The <code>A2B_MBOX1CTL.MB1DIR</code> bit controls the direction of Mailbox 1. |
| | | 0 Mailbox 1 is Receive Mailbox |
| | | 1 Mailbox 1 is Transmit Mailbox |

Table 7-104: A2B_MBOX1CTL Register Fields (Continued)

| Bit No. (Access) | Bit Name | Description/Enumeration | |
|---------------------|----------|--|--------------------|
| 0 (R/W) | MB1EN | Mailbox 1 Enable. Setting the A2B_MBOX1CTL.MB1EN bit enables Mailbox 1. | |
| | | 0 | Mailbox 1 Disabled |
| | | 1 | Mailbox 1 Enabled |

Mailbox 1 Status Register (Subordinate Only)

The `A2B_MBOX1STAT` register reports the status of the configured mailbox interrupts.

Address: 0x97

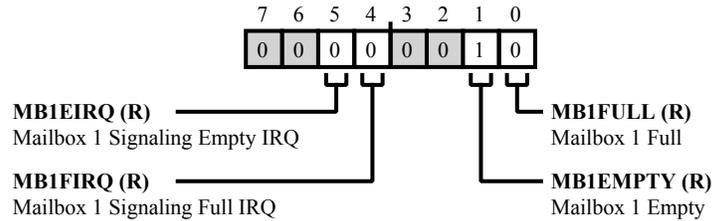


Figure 7-104: A2B_MBOX1STAT Register Diagram

Table 7-105: A2B_MBOX1STAT Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|--|
| 5 (R/NW) | MB1EIRQ | Mailbox 1 Signaling Empty IRQ. The <code>A2B_MBOX1STAT.MB1EIRQ</code> bit indicates whether or not the Mailbox 1 empty interrupt is active. |
| | | 0 Mailbox 1 Empty Interrupt Inactive |
| | | 1 Mailbox 1 Empty Interrupt Active |
| 4 (R/NW) | MB1FIRQ | Mailbox 1 Signaling Full IRQ. The <code>A2B_MBOX1STAT.MB1FIRQ</code> bit indicates whether or not the Mailbox 1 full interrupt is active. |
| | | 0 Mailbox 1 Full Interrupt Inactive |
| | | 1 Mailbox 1 Full Interrupt Active |
| 1 (R/NW) | MB1EMPTY | Mailbox 1 Empty. The <code>A2B_MBOX1STAT.MB1EMPTY</code> bit indicates whether or not Mailbox 1 is empty. |
| | | 0 Mailbox 1 Currently Not Empty |
| | | 1 Mailbox 1 Currently Empty |
| 0 (R/NW) | MB1FULL | Mailbox 1 Full. The <code>A2B_MBOX1STAT.MB1FULL</code> bit indicates whether or not Mailbox 1 is full. |
| | | 0 Mailbox 1 Currently Not Full |
| | | 1 Mailbox 1 Currently Full |

Mailbox 1 Byte 0 Register (Subordinate Only)

Address: 0x98

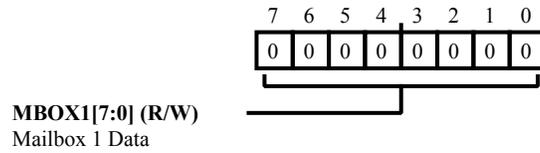


Figure 7-105: A2B_MBOX1B0 Register Diagram

Table 7-106: A2B_MBOX1B0 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|-------------------------|
| 7:0 (R/W) | MBOX1 | Mailbox 1 Data. |

Mailbox 1 Byte 1 Register (Subordinate Only)

Address: 0x99

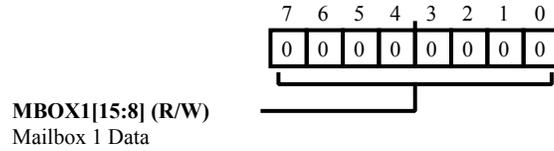


Figure 7-106: A2B_MBOX1B1 Register Diagram

Table 7-107: A2B_MBOX1B1 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|-------------------------|
| 7:0 (R/W) | MBOX1 | Mailbox 1 Data. |

Mailbox 1 Byte 2 Register (Subordinate Only)

Address: 0x9A

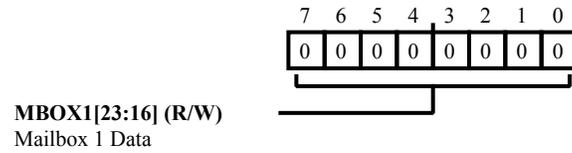


Figure 7-107: A2B_MBOX1B2 Register Diagram

Table 7-108: A2B_MBOX1B2 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|-------------------------|
| 7:0 (R/W) | MBOX1 | Mailbox 1 Data. |

Mailbox 1 Byte 3 Register (Subordinate Only)

Address: 0x9B

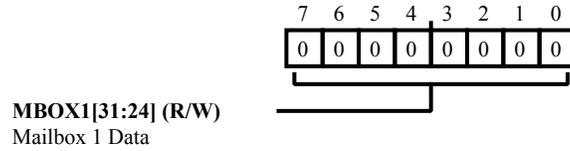


Figure 7-108: A2B_MBOX1B3 Register Diagram

Table 7-109: A2B_MBOX1B3 Register Fields

| Bit No. (Access) | Bit Name | Description/Enumeration |
|---------------------|----------|-------------------------|
| 7:0 (R/W) | MBOX1 | Mailbox 1 Data. |

8 Appendix A: Additional Discovery Flow Examples

The following sections provide additional information on modified, optimized, and advanced discovery flows. Any of the software flow diagrams can be used as a guide for discovery and initialization.

Modified Discovery Flow

In the *Modified Discovery Flow* figure, all of the subordinate nodes are discovered and immediately initialized, sequentially, from subordinate 0 to the last available subordinate in the system.

There is no further need for bus management after all nodes are discovered and programmed. But interrupt service routines may be used to react to special events (for example, an IRQ event from diagnosis). The IRQ pin can be used to signal such an event. Alternatively, the `A2B_INTTYPE` register can be polled to monitor interrupt events.

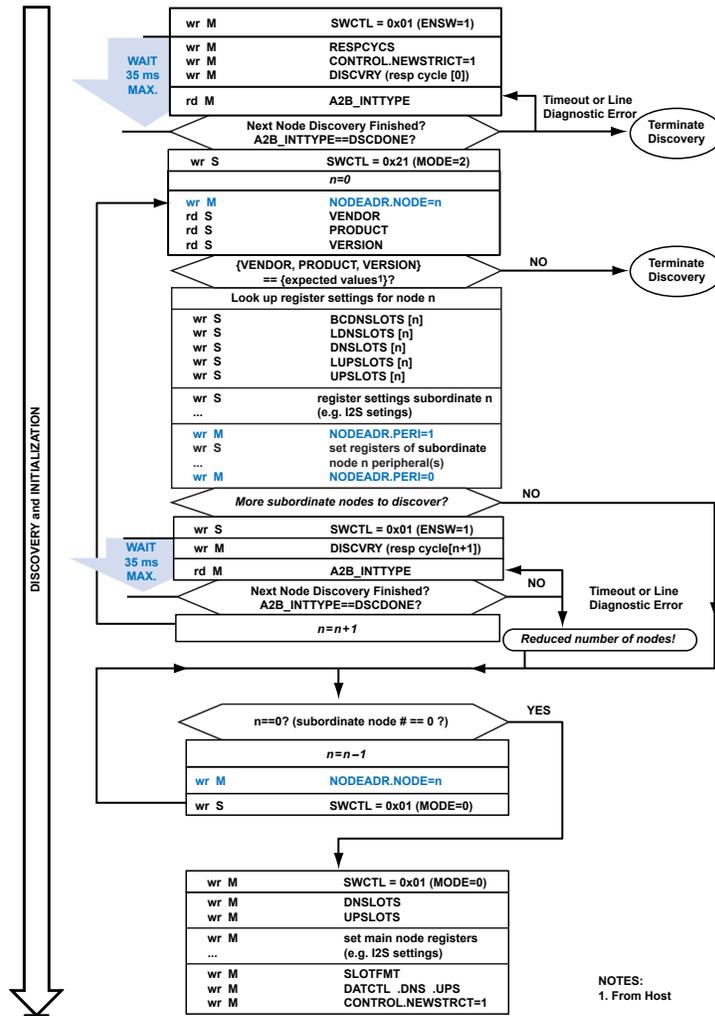


Figure 8-1: Modified Discovery Flow

Optimized Discovery Flow

A more optimized, fast discovery and initialization are shown in the *Optimized Discovery Flow* figure. Even before a node is initialized, the host tries to discover the next node. The time for the next node to be discovered is used to initialize the current node. This reduces the discovery and initialization time almost completely to the time it takes for the PLLs to find lock. Interrupt service routines are used to avoid repeated polling of registers, reducing the burden on the host processor.

There is no further need for bus management after all nodes are discovered and initialized. Interrupt service routines can be used to react to special events (for example, an IRQ event from diagnosis).

An advanced feature in the flow diagram is the use of node IDs. Node IDs allow the host to look up register settings based on IDs stored in each subordinate node’s EEPROM.

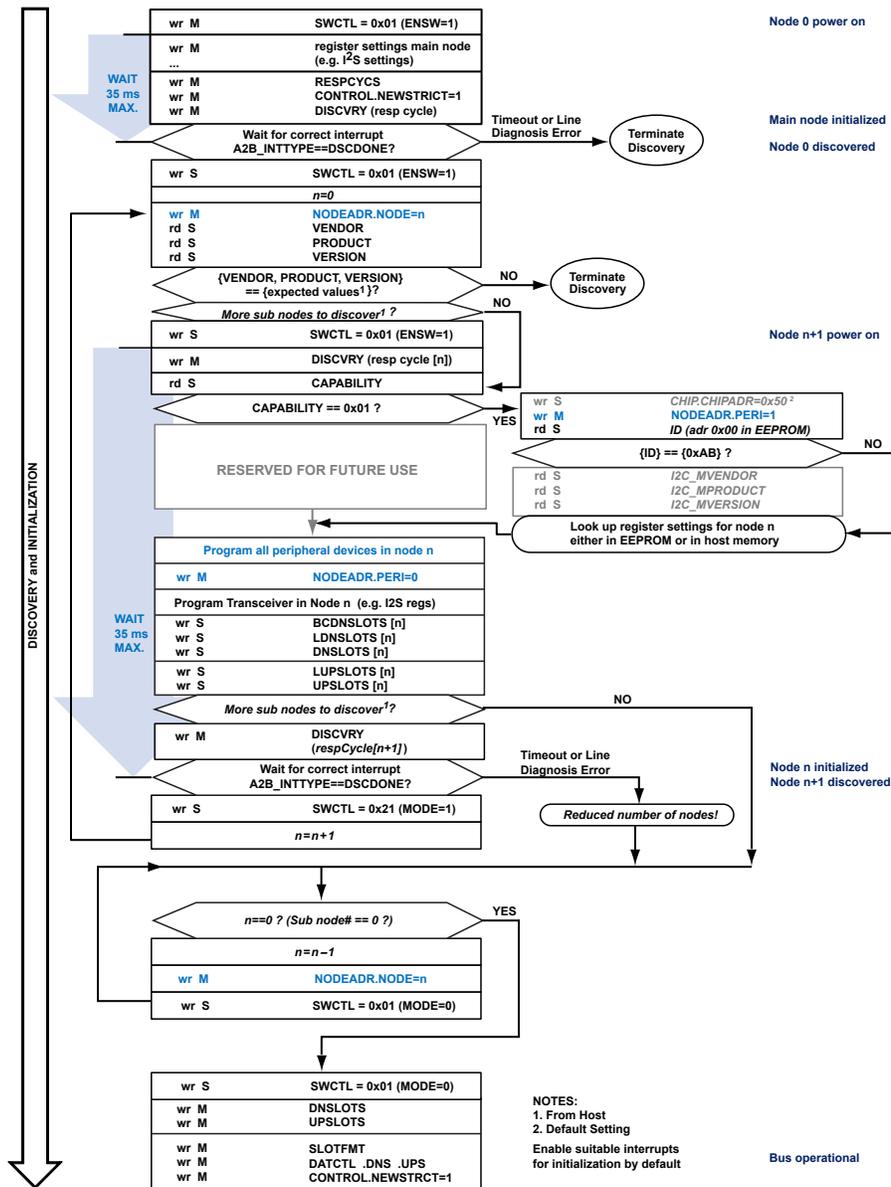


Figure 8-2: Optimized Discovery Flow

Advanced Discovery Flow

An advanced, fast flow of discovery and initialization is shown in the *Advanced Discovery Flow* figure. Even before a node is initialized, the host tries to discover the next node. The time for the next node to be discovered is used to initialize the current node. This reduces the discovery and initialization time almost completely to the time it takes for the PLLs to find lock. Synchronous exchange of data can start as soon as a main node and subordinate 0 node are initialized, while the next nodes that are not discovered and initialized can start up gradually. Use interrupt service routines to avoid repeated polling of registers, which reduces the burden on the host processor.

Another advanced feature in this flow diagram is the use of node IDs. Node IDs allow the host to look up register settings based on IDs stored in EEPROM of each subordinate node.

The subordinate nodes are reconfigured with the addition of every new node to adjust the amount of payload and, therefore, optimize bandwidth and power consumption. The optimum bus activity level is achieved with every addition of a new node even when not of the all nodes can be discovered.

This is especially advantageous when a host tries to perform “auto-discovery” without prior knowledge of the number of nodes in the system. The `A2B_DNSLOTS` and `A2B_UPSLOTS` register values can be calculated based on the `A2B_BCDNSLOTS`, `A2B_LDNSLOTS`, and `A2B_LUPSLOTS` information in each node. This can be part of the node ID capability information (for example, in the EEPROM of each subordinate node) or can be looked up based on the capability information.

Changing `A2B_DNSLOTS` and `A2B_UPSLOTS` in all nodes, depending on the number of nodes discovered, has an impact on the main nodes's I²S/TDM interface. The channel allocation changes when a new node that provides or consumes synchronous data is added.

Allowing synchronous payload operation on early nodes before the bus is fully discovered may or may not be desirable. The advanced discovery flow can be modified so that synchronous audio operation only starts after discovery (see [Optimized Discovery Flow](#)).

There is no further need for bus management after all of the nodes are discovered and initialized. Interrupt service routines can be used to react to special events (for example, an IRQ event from diagnosis).

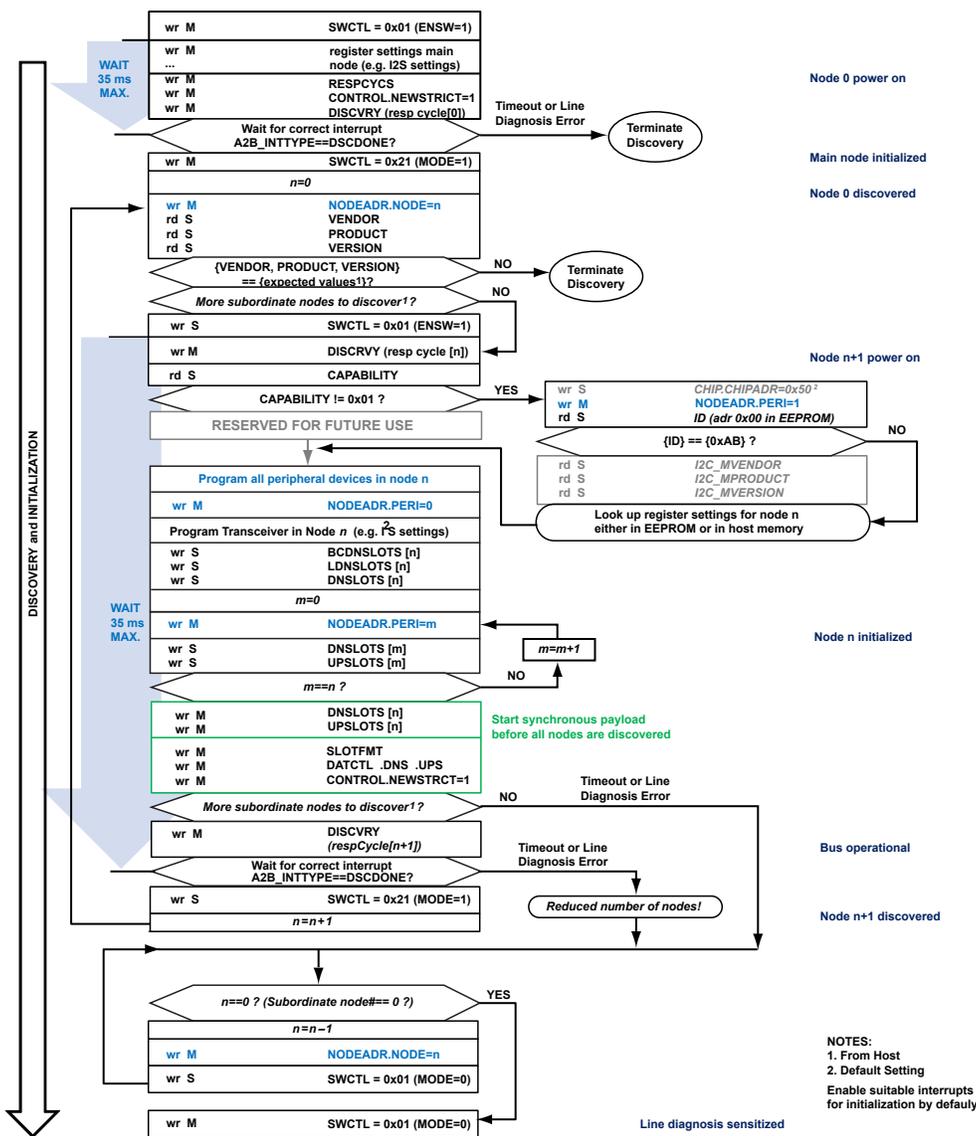


Figure 8-3: Advanced Discovery Flow

9 Appendix B: Response Cycle Formula

The `A2B_RESPCYCS` register is used to set the relative time, from the start of a control frame (SCF) to the moment the last subordinate node responds with a response frame (SRF). The register setting defines when earlier nodes in the A²B network should expect the response from the last subordinate node during the upstream portion of the superframe. If the last node fails to respond, the node immediately before the presumed last node does respond. The following sections provide information regarding how to program the main node and subordinate node `A2B_RESPCYCS` registers.

Configuring Main Node Response Cycles

The *Main Node Response Cycles* figure depicts how the main node response cycle value is determined.

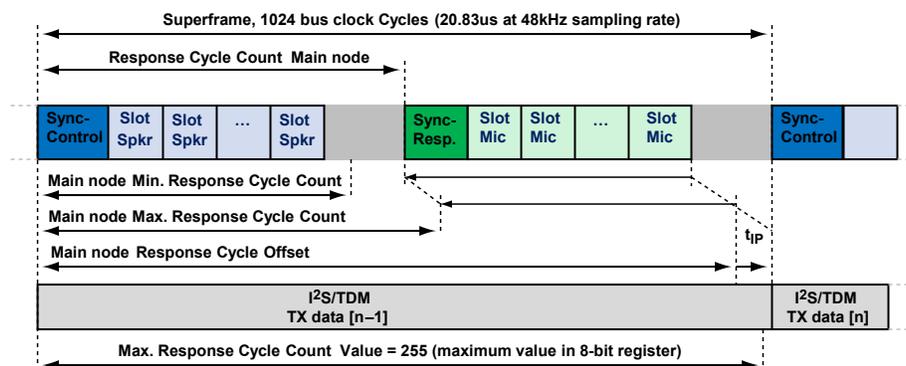


Figure 9-1: Main Node Response Cycles

In the *Main Node Response Cycles* figure:

- The *Main Node Minimum Response Cycle Count* is determined by the length of the downstream data, the minimum bus turn-around time, and the number of subordinate nodes.
- The *Main Node Maximum Response Cycle Count* is determined by the length of the upstream data and the *Main Node Response Cycle Offset*.
- The *Main Node Response Cycle Offset* ensures that sufficient internal processing time (t_{IP}) is available from the reception of the last upstream data bit in the receive buffer to the point at which this I²S/TDM data is output,

which starts synchronous to the next SCF and SYNC pin transition. The *A²B Main Node Response Offset (RESPOFFS)* table defines this constant *Main Node Response Cycle Offset*, which is a function of the A²B main node's TDM mode (`A2B_I2SGCFG.TDMMODE`) and I²S/TDM channel size (`A2B_I2SGCFG.TDMSS`).

Table 9-1: A²B Main Node Response Offset (RESPOFFS)

| TDM Mode (A ² B Main Node) | TDM Data Width (A ² B Main Node) | RESPOFFS |
|--|--|----------|
| TDM2/I ² S (<code>A2B_I2SGCFG.TDMMODE = 0</code>) | 16 bits (<code>A2B_I2SGCFG.TDMSS = 1</code>) | 238 |
| TDM2/I ² S (<code>A2B_I2SGCFG.TDMMODE = 0</code>) | 32 bits (<code>A2B_I2SGCFG.TDMSS = 0</code>) | 245 |
| TDM4 (<code>A2B_I2SGCFG.TDMMODE = 1</code>) | 16 bits (<code>A2B_I2SGCFG.TDMSS = 1</code>) | 245 |
| TDM4 (<code>A2B_I2SGCFG.TDMMODE = 1</code>) | 32 bits (<code>A2B_I2SGCFG.TDMSS = 0</code>) | 248 |
| TDM8 (<code>A2B_I2SGCFG.TDMMODE = 2</code>) | 16 bits (<code>A2B_I2SGCFG.TDMSS = 1</code>) | 248 |
| TDM8 (<code>A2B_I2SGCFG.TDMMODE = 2</code>) | 32 bits (<code>A2B_I2SGCFG.TDMSS = 0</code>) | 248 |
| TDM12 (<code>A2B_I2SGCFG.TDMMODE = 3</code>) | 16 bits (<code>A2B_I2SGCFG.TDMSS = 1</code>) | 248 |
| TDM12 (<code>A2B_I2SGCFG.TDMMODE = 3</code>) | 32 bits (<code>A2B_I2SGCFG.TDMSS = 0</code>) | 248 |
| TDM16 (<code>A2B_I2SGCFG.TDMMODE = 4</code>) | 16 bits (<code>A2B_I2SGCFG.TDMSS = 1</code>) | 248 |
| TDM16 (<code>A2B_I2SGCFG.TDMMODE = 4</code>) | 32 bits (<code>A2B_I2SGCFG.TDMSS = 0</code>) | 248 |
| TDM20 (<code>A2B_I2SGCFG.TDMMODE = 5</code>) | N/A | 248 |
| TDM24 (<code>A2B_I2SGCFG.TDMMODE = 6</code>) | N/A | 248 |
| TDM32 (<code>A2B_I2SGCFG.TDMMODE = 7</code>) | N/A | 248 |

Programming the main node `A2B_RESPCYCS` register is a function of the above *Main Node Response Cycle Offset* (RESPOFFS), as well as:

- the number of subordinate nodes in the system,
- the number of downstream A²B bus data slots received on the A-PORT at each subordinate (NUM_DNSLOTS),
- the width of the downstream A²B bus data slots (DNSLOT_SIZE),
- the number of upstream A²B bus data slots driven to the A-PORT by each subordinate (NUM_UPSLOTS), and
- the width of the upstream A²B bus data slots (UPSLOT_SIZE).

The upslot and downslot activity that is possible at any given node in the system is the first factor that contributes toward determining the value that must be programmed into the main node's `A2B_RESPCYCS` register. For each subordinate node *n* in the A²B topology, the following equations define the downstream (DNSLOT_ACTIVITY[n]) and upstream (UPSLOT_ACTIVITY[n]) activity for that node.

$$\begin{aligned} \text{DNSLOT_ACTIVITY}[n] &= \text{NUM_DNSLOTS} * (\text{DNSLOT_SIZE} + 1) \\ \text{UPSLOT_ACTIVITY}[n] &= \text{NUM_UPSLOTS} * (\text{UPSLOT_SIZE} + 1) \end{aligned}$$

NOTE: The DNSLOT_SIZE and UPSLOT_SIZE slot sizes are offset by 1 in the above calculations because the default slot format ([A2B_SLOTFMT](#)) appends a single parity bit to each data slot on the A²B bus, thereby increasing the number of bits on the A²B bus per slot by 1. For alternate slot formats, the number of bits that are appended for the chosen use case must be added instead of the 1 defined here, as presented in the A²B Bus Bits column in the Slot Format table in [A²B Slot Format](#).

Once the upslot and downslot activity for each subordinate node n is established, the equivalent upstream (RESPCYCS_UP[n]) and downstream (RESPCYCS_DN[n]) response cycle requirements can be calculated for each subordinate node, as governed by the following equations.

```
RESPCYCS_DN[n] = ((64 + DNSLOT_ACTIVITY[n])/4) + 4n + 2 // Round Up
RESPCYCS_UP[n] = RESPOFFS - (((64 + UPSLOT_ACTIVITY[n])/4) + 1) // Round Up
```

- RESPCYCS_DN[n] is the minimum response cycle register setting possible at the main node when considering the downstream activity at subordinate node n. The maximum value among those calculated for RESPCYCS_DN[n] is the minimum main node [A2B_RESPCYCS](#) setting (MAX(RESPCYCS_DN[n])).
- RESPCYCS_UP[n] is the maximum response cycle register setting possible at the main node when considering the upstream activity at subordinate node n. The minimum value among those calculated for RESPCYCS_UP[n] is the maximum main node [A2B_RESPCYCS](#) setting (MIN(RESPCYCS_UP[n])).

CAUTION: If MAX(RESPCYCS_DN[n]) > MIN(RESPCYCS_UP[n]), then the A²B bus bandwidth cannot accommodate the configuration.

The value that must be programmed into the main node's [A2B_RESPCYCS](#) register is the average of these minimum and maximum values:

```
A2B_RESPCYCS = (MAX(RESPCYCS_DN[n]) + MIN(RESPCYCS_UP[n])) / 2 // Round Down
```

Example Main Node [A2B_RESPCYCS](#) Calculation

A system with three nodes, the main node and two subordinate nodes (subordinate 0 and subordinate 1), is configured as shown in the *Three-Node A²B System Example* figure:

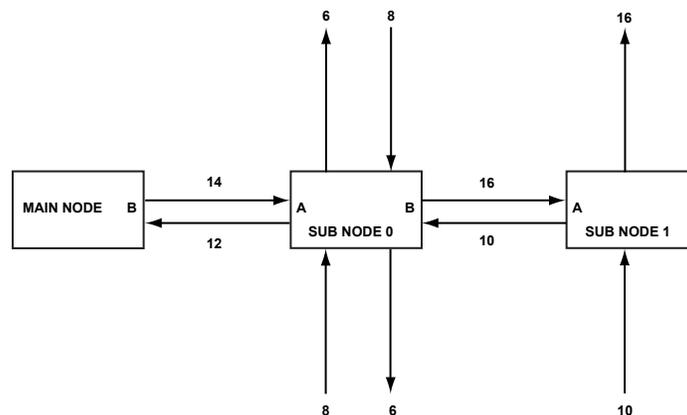


Figure 9-2: Three-Node A²B System Example

For the downstream portion of the superframe:

- Main node (configured for 32-bit TDM8 mode): sends 14 slots with a 24-bit slot size
- Subordinate 0: consumes six slots from the main node and passes the remaining eight slots to subordinate 1, then contributes eight additional slots to the downstream traffic (16 total slots sent from subordinate 0 to subordinate 1)
- Subordinate 1: consumes all 16 slots coming from subordinate-0

For the upstream portion of the superframe:

- Subordinate 1: sends ten slots with a 16-bit slot size
- Subordinate 0: consumes six slots from subordinate-1 and passes the remaining four slots to the main node, then contributes eight additional slots to the upstream traffic (12 total slots sent from subordinate 0 to the main node).
- Main node (configured for 32-bit TDM8 mode): consumes all 12 slots coming from subordinate 0.

The response cycle is determined using the following steps:

1. Calculate the upslot and downslot activity for each subordinate node:

```

DNSLOT_ACTIVITY[n] = NUM_DNSLOTS * (DNSLOT_SIZE + 1)
DNSLOT_ACTIVITY[0] = 14 * (24 + 1) = 350
DNSLOT_ACTIVITY[1] = 16 * (24 + 1) = 400

UPSLOT_ACTIVITY[n] = NUM_UPSLOTS * (UPSLOT_SIZE + 1)
UPSLOT_ACTIVITY[0] = 12 * (16 + 1) = 204
UPSLOT_ACTIVITY[1] = 10 * (16 + 1) = 170

```

2. With this information, calculate the response cycle requirements for each subordinate. From the *A²B Main Node Response Offset (RESPOFFS)* table, the TDM8 mode and 32-bit data combination yields RESPOFFS = 248.

```

RESPCYCS_DN[n] = ((64 + DNSLOT_ACTIVITY[n])/4) + 4n + 2 // Round Up
RESPCYCS_DN[0] = ((64 + 350)/4) + (4*0) + 2 = 103.5 + 0 + 2 = 105.5 = 106
RESPCYCS_DN[1] = ((64 + 400)/4) + (4*1) + 2 = 116.0 + 4 + 2 = 122.0 = 122

RESPCYCS_UP[n] = RESPOFFS - (((64 + UPSLOT_ACTIVITY[n])/4) + 1) // Round Up
RESPCYCS_UP[0] = 248 - (((64 + 204)/4) + 1) = 248 - (67.0 + 1) = 180.0 = 180
RESPCYCS_UP[1] = 248 - (((64 + 170)/4) + 1) = 248 - (58.5 + 1) = 188.5 = 189

```

The minimum main node `A2B_RESPCYCS` setting is the maximum value among the `RESPCYCS_DN[n]` calculations, which is 122, and the maximum setting is the minimum value among the `RESPCYCS_UP[n]` calculations, which is 180, and the average of the minimum and maximum values is:

```

(MAX(RESPCYCS_DN[n]) + MIN(RESPCYCS_UP[n])) / 2 // Round Down
(122 + 180) / 2 = 302 / 2 = 151.0 = 151

```

3. For this system configuration, program the main node `A2B_RESPCYCS` value to 151 (0x97).

Configuring Subordinate Node Response Cycles

Each subordinate node has its `A2B_RESPCYCS` register set during the system discovery process. The main transceiver programs its `A2B_DISCVRY` register with the response cycle value associated with the subordinate transceiver that it is attempting to discover. The appropriate value for each subordinate node (`SLV_RESPCYCS[n]`) is a function of the subordinate node's location in the A²B topology and the value programmed to the main node's `A2B_RESPCYCS` register (`MSTR_RESPCYCS`). The subordinate node nearest to the main node has a node number of 0, and the node number is incremented for each next-in-line subordinate node until the last-in-line subordinate node `n`. The `A2B_RESPCYCS` value to use for each subordinate node during discovery can be calculated using the following equation:

$$SLV_RESPCYCS[n] = MSTR_RESPCYCS - 4n$$

Using the Example Main Node `A2B_RESPCYCS` Calculation above (with `MSTR_RESPCYCS = 151`), the following equations determine the correct `A2B_RESPCYCS` value for the two subordinate nodes:

$$SLV_RESPCYCS[0] = MSTR_RESPCYCS - (4*0) = 151 - 0 = 151 \text{ (0x97)}$$

$$SLV_RESPCYCS[1] = MSTR_RESPCYCS - (4*1) = 151 - 4 = 147 \text{ (0x93)}$$

The following code sequence uses these values to proceed through the discovery process in the example system:

```
Write MSTR_RESPCYCS to the A2B_RESPCYCS register in the main node
Write 0x01 to the A2B_CONTROL register in the main node
Write 0x01 to the A2B_SWCTL register in the main node
Write 0x01 to the A2B_INTMSK2 register in the main node
Write SLV_RESPCYCS[0] to the A2B_DISCVRY register in the main node
    <Wait for Interrupt>

Write 0x00 to the A2B_NODEADR register in the main node
Write 0x01 to the A2B_SWCTL register in subordinate node 0
Write SLV_RESPCYCS[1] to the A2B_DISCVRY register in the main node
    <Wait for Interrupt>
```

Configuring Subordinate Node Response Cycles to Account for Cable Length

The transceiver is designed to adjust automatically to the time when responses are seen on the bus. This typically allows subordinate nodes to be discovered without requiring changes to the subordinate node response cycles, based on cable length. However, if the cable lengths in the system are known, use the following pseudocode to calculate the subordinate node response cycles:

```
if (n = 0)
    SLV_RESPCYCS[n] = MSTR_RESPCYCS
else
    if (cable_length > 12m)
        SLV_RESPCYCS[n] = SLV_RESPCYCS[n-1] - 6
    else if (cable_length > 5m)
        SLV_RESPCYCS[n] = SLV_RESPCYCS[n-1] - 5
    else
        SLV_RESPCYCS[n] = SLV_RESPCYCS[n-1] - 4
```

When discovering the A²B nodes without adjusting the response cycles for the cable length, it is possible for transient bit errors to be reported immediately, following the discovery of a new subordinate that is connected by a long cable (greater than 5m in length). These errors only persist for two to three superframes and can be safely cleared. Use the pseudocode above to adjust the response cycles for cable length and prevent bit errors during discovery.

Configuring Subordinate Node Response Cycles

Each subordinate node has its `A2B_RESPCYCS` register set during the system discovery process. The main transceiver programs its `A2B_DISCVRY` register with the response cycle value associated with the subordinate transceiver that it is attempting to discover. The appropriate value for each subordinate node (`SLV_RESPCYCS[n]`) is a function of the subordinate node's location in the A²B topology and the value programmed to the main node's `A2B_RESPCYCS` register (`MSTR_RESPCYCS`). The subordinate node nearest to the main node has a node number of 0, and the node number is incremented for each next-in-line subordinate node until the last-in-line subordinate node *n*. The `A2B_RESPCYCS` value to use for each subordinate node during discovery can be calculated using the following equation:

$$SLV_RESPCYCS[n] = MSTR_RESPCYCS - 4n$$

Using the Example Main Node `A2B_RESPCYCS` Calculation above (with `MSTR_RESPCYCS = 151`), the following equations determine the correct `A2B_RESPCYCS` value for the two subordinate nodes:

$$SLV_RESPCYCS[0] = MSTR_RESPCYCS - (4*0) = 151 - 0 = 151 \quad (0x97)$$

$$SLV_RESPCYCS[1] = MSTR_RESPCYCS - (4*1) = 151 - 4 = 147 \quad (0x93)$$

The following code sequence uses these values to proceed through the discovery process in the example system:

```
Write MSTR_RESPCYCS to the A2B_RESPCYCS register in the main node
Write 0x01 to the A2B_CONTROL register in the main node
Write 0x01 to the A2B_SWCTL register in the main node
Write 0x01 to the A2B_INTMSK2 register in the main node
Write SLV_RESPCYCS[0] to the A2B_DISCVRY register in the main node
    <Wait for Interrupt>

Write 0x00 to the A2B_NODEADR register in the main node
Write 0x01 to the A2B_SWCTL register in subordinate node 0
Write SLV_RESPCYCS[1] to the A2B_DISCVRY register in the main node
    <Wait for Interrupt>
```

10 Appendix C: Module ID and Module Configuration Memory

Module-specific descriptor information is saved in a storage device (EEPROM or similar), directly connected to an A²B transceiver via I²C and accessible over the A²B bus as a peripheral device. Such I²C-connected storage devices use the device address 0x50 (7-bit). This configuration memory contains module ID information and optional configuration blocks.

Configuration Memory

The contents of a configuration memory with no configuration blocks is shown in the *Memory Content with no Configuration Blocks* table.

| ADDRESS | CONTENTS |
|---------|---------------------------------------|
| 0x0000 | 0xAB (Indicates Configuration Memory) |
| 0x0001 | Module Vendor ID* |
| 0x0002 | Module Product ID |
| 0x0003 | Module Version ID |
| 0x0004 | Reserved - value should be ignored |
| 0x0005 | 0x00 (Number of Configuration Blocks) |
| 0x0006 | Reserved - value should be 0x00 |
| 0x0007 | CRC-8 |

*Assignment and management of Module Vendor IDs currently resides with Analog Devices Inc.

Figure 10-1: Memory Content with no Configuration Blocks

During and after discovery, the host can uniquely identify the subordinate node modules, based on the convention in the table. This information allows the host to look up all stored configuration settings and software drivers to automatically configure the A²B system, program A²B nodes, and initialize peripheral devices. A CRC byte is used to ensure data integrity.

Additionally, device specific configuration and setup information also can be stored in the configuration memory through the use of configuration blocks. The host can read this information and set up the subordinate node without any prior knowledge of the node. The contents of a configuration memory with configuration blocks is shown in the *Memory Content with Configuration Blocks* table.

| ADDRESS | CONTENTS |
|--|---------------------------------------|
| 0x0000 | 0xAB (Indicates Configuration Memory) |
| 0x0001 | Module Vendor ID |
| 0x0002 | Module Product ID |
| 0x0003 | Module Version ID |
| 0x0004 | Reserved - value should be ignored |
| 0x0005 | Number of Configuration Blocks |
| 0x0006 | Reserved - value should be 0x00 |
| 0x0007 | CRC-8 |
| 0x0008 to 7 + L ₁ | Configuration Block 1 |
| 8 + L ₁ to 7 + L ₁ + L ₂ | Configuration Block 2 |
| | • • • |
| | Configuration Block N |

Figure 10-2: Memory Content with Configuration Blocks

The contents of a configuration block are shown in the *Configuration Block Contents* figure.

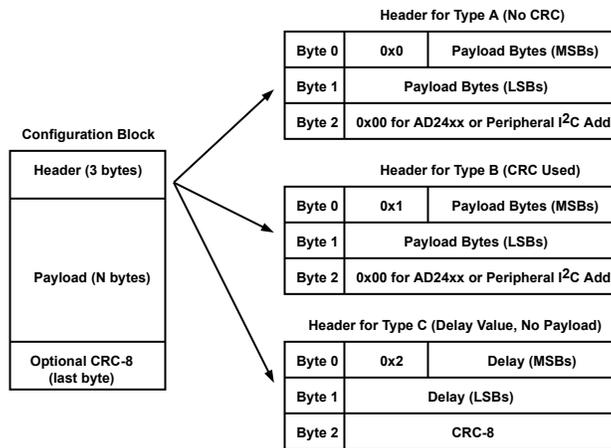


Figure 10-3: Configuration Block Contents

The first three bytes of a configuration block make up a header, which provides details about the configuration block. The first four bits of the header indicate the type of configuration block; see the *Configuration Block Header Types* table. Types A and B have a 12-bit field that gives the size (in bytes) of the payload. For a Type A configuration block, this field contains the number of bytes to be written during configuration. For a Type B configuration block, the value of this field is one more than the number of bytes to be written during configuration, because an 8-bit CRC is included at the end of the payload.

If the device to be programmed requires an address pointer, it is given at the start of the payload field. A Type C configuration block has a 12-bit field, which describes a delay to be inserted in the programming flow (in ms).

Table 10-1: Configuration Block Header Types

| Type Value | Meaning | Notes |
|------------|--|---|
| 0x0 | Type A config block, no CRC | All payload bytes written to target for configuration |
| 0x1 | Type B config block CRC-8 calculated on header + payload | Last payload byte not written to target |
| 0x2 | Type C config block delay value only (no payload) | CRC-8 calculated based on first 2 bytes in header |
| 0x3 - 0xF | Reserved | N/A |

The *Detailed View of Configuration Memory* figure shows a detailed view of configuration memory contents with N configuration blocks.

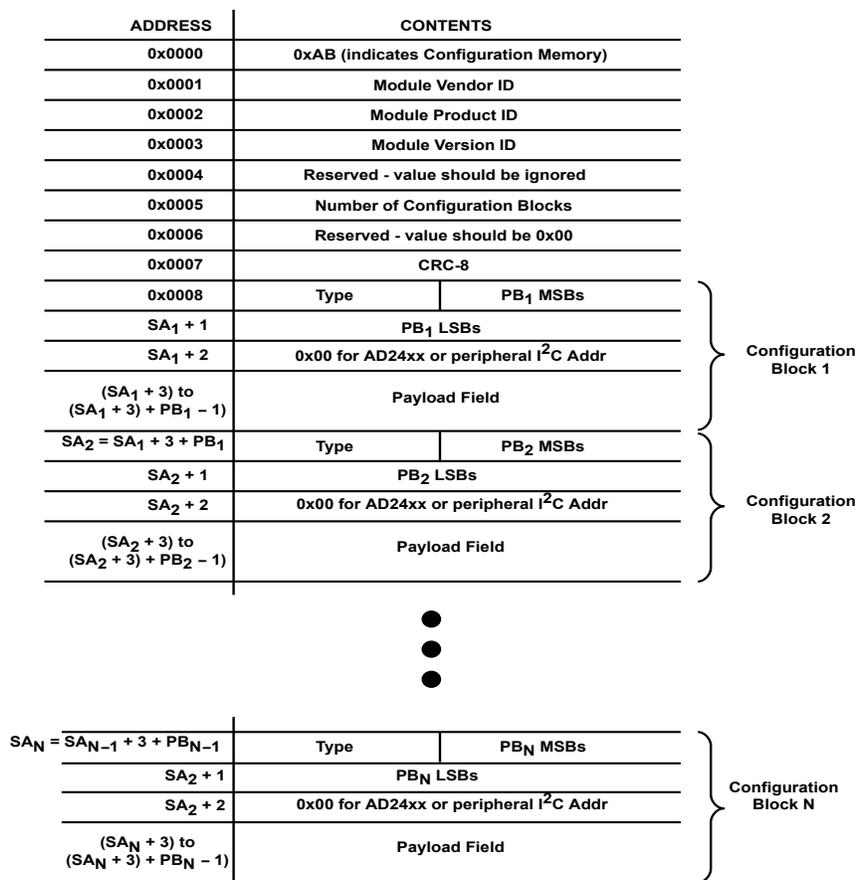


Figure 10-4: Detailed View of Configuration Memory

Notes:

- If address 0x0005 returns 0x00, no configuration blocks are present in the memory.
- PB_N is the number of bytes in the payload for a configuration block N (a 12-bit field).
- SA_N is the start address of a configuration block N. SA_N = SA_{N-1} + 3 + PB_{N-1}. SA₁ = 8.

The following tables show two examples of configuration memory containing programming information for an A²B subordinate node. The first byte of each payload field is a starting address for a burst write.

| ADDRESS | CONTENTS | | |
|------------------------|--|-----------|------------------------|
| 0x0000 | 0xAB | | |
| 0x0001 | Module Vendor ID | | |
| 0x0002 | Module Product ID | | |
| 0x0003 | Module Version ID | | |
| 0x0004 | Reserved | | |
| 0x0005 | 0x02 (Number of Configuration Blocks) | | |
| 0x0006 | Reserved - value should be 0x00 | | |
| 0x0007 | CRC-8 | | |
| SA ₁ 0x0008 | 0x00 (Type A, PB ₁ MSBs = 0) | } Payload | } Configuration Block |
| 0x0009 | 0x04 (PB ₁ LSBs = 4) | | |
| 0x000A | 0x00 (For AD24xx programming) | | |
| 0x000B | 0x0A (Address pointer for BCDNSLOTS) | | |
| 0x000C | Data for BCDNSLOTS | | |
| 0x000D | Data for LDNSLOTS | | |
| 0x000E | Data for LUPSLOTS | | |
| SA ₂ 0x000F | 0x00 (Type A, PB ₂ MSBs = 0) | } Payload | } Configuration Block2 |
| 0x0010 | 0x15 (PB ₂ LSBs = 21) | | |
| 0x0011 | 0x00 (For AD24xx programming) | | |
| 0x0012 | 0x3F (Address pointer for I2CCFG) | | |
| 0x0013 | Data for I2CCFG | | |
| 0x0014 | 0x00 (Data for PLLCTL) | | |
| 0x0015 | Data for I2SGCFG | | |
| 0x0016 | Data for I2SCFG | | |
| 0x0017 | Data for I2SRATE | | |
| 0x0018 | Reserved (ADDR 0x44 reserved for sub node) | | |
| 0x0019 | Reserved (ADDR 0x45 reserved for sub node) | | |
| 0x001A | Data for SYNCOFFSET | | |
| 0x001B | Data for PDMCTL | | |
| 0x001C | Data for ERRMGMT | | |
| 0x001D | 0x00 (addr 0x49 reserved) for AD242x | | |
| 0x001E | Data for GPIODAT | | |
| 0x001F | Data for GPIODATSET | | |
| 0x0020 | Data for GPIODATCLR | | |
| 0x0021 | Data for GPIOOEN | | |
| 0x0022 | Data for GPIOIEN | | |
| 0x0023 | 0x00 (Data for GPIOIN) | | |
| 0x0024 | Data for PINTEN | | |
| 0x0025 | Data for PINTINV | | |
| 0x0026 | Data for PINCFG | | |

Figure 10-5: Configuration Memory for AD242x Subordinate Node Configuration (Long)

| ADDRESS | CONTENTS | | |
|------------------------|--|---|-----------------------|
| 0x0000 | 0xAB | | |
| 0x0001 | Module Vendor ID | | |
| 0x0002 | Module Product ID | | |
| 0x0003 | Module Version ID | | |
| 0x0004 | Reserved | | |
| 0x0005 | 0x02 (Number of Configuration Blocks) | | |
| 0x0006 | Reserved - value should be 0x00 | | |
| 0x0007 | CRC-8 | | |
| SA ₁ 0x0008 | 0x10 (Type B, PB ₁ MSBs = 0) | } | Configuration Block 1 |
| 0x0009 | 0x04 (PB ₁ LSBs = 4) | | |
| 0x000A | 0x00 (For AD24xx programming) | | |
| 0x000B | 0x0B (Address pointer for LDNSLOTS) | | |
| 0x000C | Data for LDNSLOTS | | |
| 0x000D | Data for LDNSLOTS | | |
| 0x000E | CRC-8 | | |
| SA ₂ 0x000F | 0x10 (Type B, PB ₂ MSBs = 0) | } | Configuration Block 2 |
| 0x0010 | 0x0B (PB ₂ LSBs = 11) | | |
| 0x0011 | 0x00 (For AD24xx programming) | | |
| 0x0012 | 0x41 (Address pointer for I2SGCFG) | | |
| 0x0013 | Data for I2SGCFG | | |
| 0x0014 | Data for I2SCFG | | |
| 0x0015 | Data for I2SRATE | | |
| 0x0016 | Reserved (ADDR 0x44 reserved for sub node) | | |
| 0x0017 | Reserved (ADDR 0x45 reserved for sub node) | | |
| 0x0018 | Data for SYNCOFFSET | | |
| 0x0019 | Data for PDMCTL | | |
| 0x001A | Data for ERRMGMT | | |
| 0x001B | 0x00 (addr 0x49 reserved) for AD242x | | |
| 0x001C | CRC-8 | | |

Figure 10-6: Configuration Memory for AD242x Subordinate Node Configuration (Short)

The *Configuration Memory for ADAU1761* figure shows an example of configuration memory containing programming information for an ADAU1761 codec (which uses two address bytes per transaction).

| Address | Contents | | |
|-----------------|---------------------------------------|--|-----------------------|
| 0x0000 | 0xAB | | |
| 0x0001 | Module Vendor ID | | |
| 0x0002 | Module Product ID | | |
| 0x0003 | Module Version ID | | |
| 0x0004 | Reserved | | |
| 0x0005 | 0x08 (Number of Configuration Blocks) | | |
| 0x0006 | Reserved – value should be 0x00 | | |
| 0x0007 | CRC-8 | | |
| SA ₁ | 0x0008 | 0x00 (Type A, PB ₁ MSBs = 0) | Configuration Block 1 |
| | 0x0009 | 0x03 (PB ₁ LSBs = 3) | |
| | 0x000A | 0x39 (Peripheral I ² C Address) | |
| | 0x000B | 0x40 (Address MSB) | |
| | 0x000C | 0x00 (Address LSB) | |
| | 0x000D | 0x0f (Data for address 0x4000) | Configuration Block 2 |
| SA ₂ | 0x000E | 0x00 (Type A, PB ₂ MSBs = 0) | |
| | 0x000F | 0x08 (PB ₂ LSBs = 8) | |
| | 0x0010 | 0x39 (Peripheral I ² C Address) | |
| | 0x0011 | 0x40 (Address MSB) | |
| | 0x0012 | 0x02 (Address LSB) | Configuration Block 3 |
| | 0x0013 | 0x00 (Data for address 0x4002) | |
| | 0x0014 | 0x01 (Data for address 0x4003) | |
| | 0x0015 | 0x00 (Data for address 0x4004) | |
| | 0x0016 | 0x00 (Data for address 0x4005) | |
| | 0x0017 | 0x20 (Data for address 0x4006) | |
| | 0x0018 | 0x03 (Data for address 0x4007) | |
| SA ₃ | 0x0019 | 0x20 (Type C, Delay ₃ MSBs = 0) | Configuration Block 4 |
| | 0x001A | 0x64 (Delay ₃ LSBs = 100) | |
| | 0x001B | CRC-8 | |
| SA ₄ | 0x001C | 0x00 (Type A, PB ₄ MSBs = 0) | Configuration Block 4 |
| | 0x001D | 0x16 (PB ₄ LSBs = 22) | |
| | 0x001E | 0x39 (Peripheral I ² C Address) | |
| | 0x001F | 0x40 (Address MSB) | |
| | 0x0020 | 0x08 (Address LSB) | |
| | 0x0021 | Data for address 0x4008 | |
| | 0x0022 | Data for address 0x4009 | |
| | 0x0023 | Data for address 0x400A | |
| | | • | |
| | | • | |
| | | • | |
| | 0x0033 | Data for address 0x401A | |
| | 0x0034 | Data for address 0x401B | |
| SA ₅ | 0x0035 | 0x00 (Type A, PB ₅ MSBs = 0) | |
| | | • | |
| | | • | |
| | | • | |

Figure 10-7: Configuration Memory for ADAU1761

11 Appendix D: Interrupt Processing

The following sections describe the flow of interrupt processing by the host in the A²B system.

Main Node Running Interrupts

As shown in the *Main Node Running Interrupts* figure, the trigger (Main IRQ pin) is asserted after the main node locks its PLL to the SYNC signal or on a post discovery line fault.

NOTE: MSTR_RUNNING (A2B_INTTYPE= 0xFF) is a main node only interrupt.

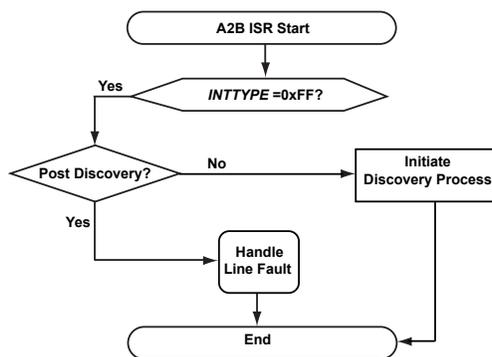


Figure 11-1: Main Node Running Interrupts

Action: read the A2B_INTSRC and A2B_INTTYPE registers and proceed to subordinate node discovery or handle a line fault. Note that a host read of the main node A2B_INTTYPE register clears the interrupt.

Discovery Done Interrupts

As shown in the *Discovery Done Interrupts* figure, the trigger (Main IRQ pin) is asserted after the main node sees a response from the discovery of a subordinate node. DSCDONE (A2B_INTTYPE= 0x18) is a main node only interrupt.

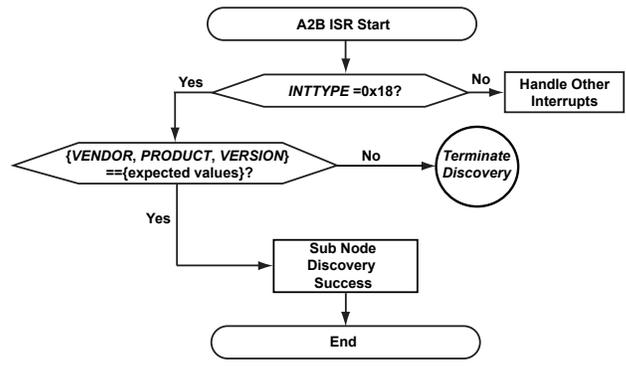


Figure 11-2: Discovery Done Interrupts

Action: read the `A2B_INTSRC` and `A2B_INTTYPE` registers and proceed to the node authentication and discovery process. Note that a host read of the main node `A2B_INTTYPE` register clears the interrupt.

Line Fault Interrupts

As shown in the *Line Fault Interrupts* figure, the trigger (main IRQ pin) is asserted after encountering a line fault during or post discovery.

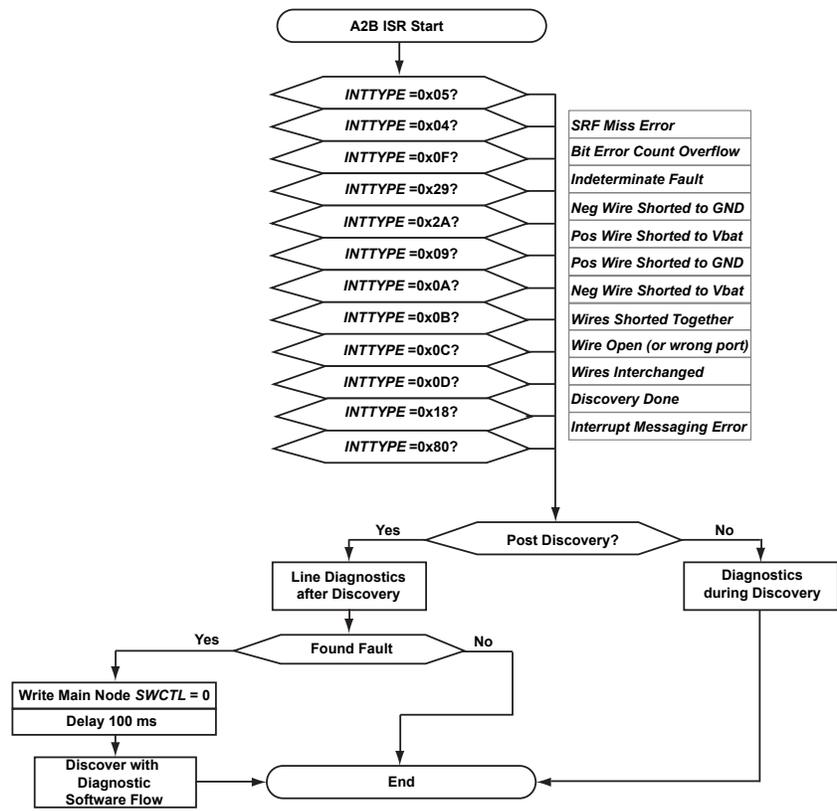


Figure 11-3: Line Fault Interrupts

Action: read the `A2B_INTTYPE` register and proceed with line diagnostics, as described in [A²B System Debug](#).

When the transceiver enters the RESET state due to critical faults such as BP short to GND, there is no indication to the host that this has occurred. If such functionality is desired in the system, designs can utilize termination resistors on the IRQ line as a function of its active polarity, as governed by the `A2B_PINCFG.IRQINV` bit. When `A2B_PINCFG.IRQINV = 0`, a pull-up resistor connected to the IRQ line pulls IRQ high when the transceiver threestates the IRQ pin while in the RESET state. The state can then be seen by the host controller as an active high edge pseudo-interrupt. The host reads the `A2B_INTSTAT` and `A2B_INTTYPE` registers as 0x00 (reset values), which can be interpreted as an event indicating that the transceiver has entered the RESET state. A pull-down resistor on the IRQ line has the same effect when `A2B_PINCFG.IRQINV=1` for negative edge interrupts at the host controller.

NOTE: The host controller must ignore the IRQ state prior to the `A2B_CONTROL.MSTR` bit being set, upon which the IRQ pin is driven to the inactive state.

Error Interrupts

As shown in the *Error Interrupts* figure, the trigger (Main IRQ pin) is asserted when any of the following errors is encountered.

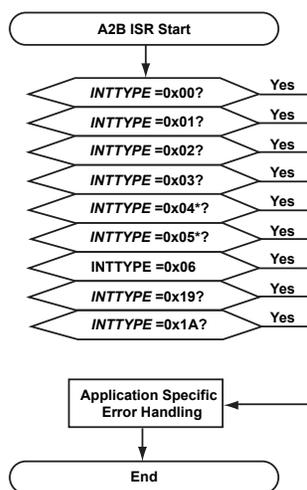


Figure 11-4: Error Interrupts

```

HDCNTERR=0x0
DDERR= 0x1
CRCERR= 0x2
DPERR= 0x3
BECOVF= 0x4* (Occurrence of Bit error count overflow interrupt, after
              resetting the error counter (BECNT) once every
              second, indicates bus issues )
SRFERR= 0x5* (10 time occurrence without interrupt status (INTSTAT)
              being cleared between pending interrupts shall be
              treated as bus lost condition/line fault)
SRFCRCERR=0x6 (Sub node Only)
  
```

I2CERR= 0x19 (Main node Only)
 ICRERR= 0x1A (Main node Only)

Action: read the [A2B_INTTYPE](#) register and proceed with line diagnostics, as described in [A²B System Debug](#).

General Purpose IO Pin Interrupts

As shown in the *General Purpose IO Pin Interrupts* figure, the trigger (Main IRQ pin) is asserted when any of the following errors is encountered.

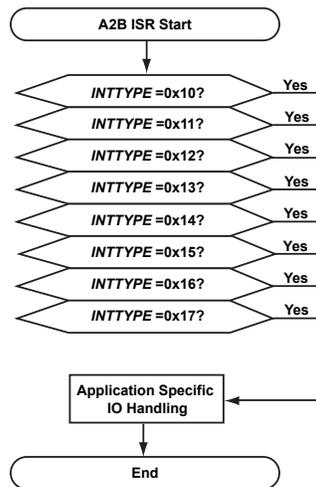


Figure 11-5: General Purpose IO Pin Interrupts

IO0= 0x10 (Sub node only)
 IO1= 0x11
 IO2= 0x12
 IO3= 0x13
 IO4= 0x14
 IO5= 0x15
 IO6= 0x16
 IO7= 0x17

Action: read the [A2B_INTTYPE](#) register and take action that is specific to the application.